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*European Standard (Telecommunications series)*

## **Digital Video Broadcasting (DVB); Interaction channel for satellite distribution systems;**

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European Broadcasting Union



Union Européenne de Radio-Télévision



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DEN/JTC-DVB-96

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## Foreword

This European Standard (Telecommunications series) has been produced by Joint Technical Committee (JTC) of the European Broadcasting Union (EBU), Comité Européen de Normalisation Electrotechnique (CENELEC) and the European Telecommunications Standards Institute (ETSI).

NOTE: The EBU/ETSI JTC Broadcast was established in 1990 to co-ordinate the drafting of standards in the specific field of broadcasting and related fields. Since 1995 the JTC Broadcast became a tripartite body by including in the Memorandum of Understanding also CENELEC, which is responsible for the standardization of radio and television receivers. The EBU is a professional association of broadcasting organizations whose work includes the co-ordination of its members' activities in the technical, legal, programme-making and programme-exchange domains. The EBU has active members in about 60 countries in the European broadcasting area; its headquarters is in Geneva.

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Founded in September 1993, the DVB Project is a market-led consortium of public and private sector organizations in the television industry. Its aim is to establish the framework for the introduction of MPEG-2 based digital television services. Now comprising over 200 organizations from more than 25 countries around the world, DVB fosters market-led systems, which meet the real needs, and economic circumstances, of the consumer electronics and the broadcast industry.

<b>National transposition dates</b>	
Date of adoption of this EN:	24 November 2000
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# 1 Scope

The present document is the baseline specification for the provision of the interaction channel for GEO satellite interactive networks with fixed return channel satellite terminals (RCST). The present document facilitates the use of RCSTs for individual or collective installation (e.g. SMATV) in a domestic environment. It also supports the connection of such terminals with in-house data networks. The present document may be applied to all frequency bands allocated to GEO satellite services.

The solutions provided in the present document for interaction channel for satellite interactive networks are a part of a wider set of alternatives to implement interactive services for Digital Video Broadcasting (DVB) systems.

---

# 2 References

The following documents contain provisions which, through reference in this text, constitute provisions of the present document.

- References are either specific (identified by date of publication, edition number, version number, etc.) or non-specific.
- For a specific reference, subsequent revisions do not apply.
- For a non-specific reference, the latest version applies.
- A non-specific reference to an ETS shall also be taken to refer to later versions published as an EN with the same number.

- [1] ETSI EN 300 421: "Digital Video Broadcasting (DVB); Framing structure, channel coding and modulation for 11/12 GHz satellite services".
- [2] ETSI TR 101 202: "Digital Video Broadcasting (DVB); Implementation guidelines for Data Broadcasting".
- [3] ETSI ETS 300 802: "Digital Video Broadcasting (DVB); Network-independent protocols for DVB interactive services".
- [4] ETSI EN 300 468: "Digital Video Broadcasting (DVB); Specification for Service Information (SI) in DVB systems".
- [5] ETSI EN 301 192: "Digital Video Broadcasting (DVB); DVB specification for data broadcasting".
- [6] ETSI EN 301 459: "Satellite Earth Stations and Systems (SES); Harmonized EN for Satellite Interactive Terminals (SIT) and Satellite User Terminals (SUT) transmitting towards satellites in geostationary orbit in the 29,5 GHz to 30 GHz frequency bands covering essential requirements under article 3.2 of the R&TTE Directive".
- [7] IETF RFC 1483 (1993): "Multiprotocol Encapsulation over ATM Adaptation Layer 5".
- [8] ETSI TR 100 815: "Digital Video Broadcasting (DVB); Guidelines for the handling of Asynchronous Transfer Mode (ATM) signals in DVB systems".
- [9] ISO/IEC 13818-1 (1996): "Information technology - Generic coding of moving pictures and associated audio information; Part 1: Systems".
- [10] ETSI ETR 154: "Digital Video Broadcasting (DVB); Implementation guidelines for the use of MPEG-2 Systems, Video and Audio in satellite, cable and terrestrial broadcasting applications".
- [11] ITU-T Recommendation Q.2931 (1995): "Broadband Integrated Services Digital Network (B-ISDN) - Digital Subscriber Signalling System No. 2 (DSS 2) - User-Network Interface (UNI) - Layer 3 specification for basic call/connection control".



- [12] IEEE 802.3 (1996): "Information technology - Telecommunications and information exchange between systems - Local and metropolitan area networks - Specific requirements - Part 3: Carrier sense multiple access with collision detection (CSMA/CD) access method and physical layer specifications".
- [13] ITU-T Recommendation I.432 (1999): "B-ISDN user-network interface - Physical layer specification".
- [14] ETSI ES 200 800: "Digital Video Broadcasting (DVB); DVB interaction channel for Cable TV distribution System (CATV); DVB document A023 Rev. 1, June 1999.
- [15] IETF RFC 2104 (1997): "HMAC: Keyed-Hashing for Message Authentication".
- [16] ANSI/IEEE 754 (1985): "IEEE Standard for Binary Floating-Point Arithmetic".
- [17] ISO/IEC 13818-6 (1998): "Information technology - Generic coding of moving pictures and associated audio information - Part 6: Extensions for DSM-CC".
- [18] ITU-T Recommendation I.363-5 (1996): "B-ISDN ATM Adaptation Layer specification: Type 5 AAL".

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## 3 Definitions, symbols and abbreviations

### 3.1 Definitions

For the purposes of the present document, the following term and definition applies:

**reserved:** term "reserved", when used in the clauses defining the coded bit stream, indicates that the value may be used for future extensions. Unless otherwise specified within the present document, all reserved bits shall be set to '1'.

### 3.2 Symbols

For the purposes of the present document, the following symbols apply:

*	multiplication
^	power
~	concatenation
mod	modulo division
(unsigned char)x	ANSI C cast operator: converts value x to unsigned char
""	empty string (zero length)
nonce1	random string (NCC)
nonce2	random string (RCST)
$N_{atm}$	Number of ATM cells in an ATM TRF burst (1, 2 or 4).
$N_{mpeg}$	Number of MPEG packets in an optional MPEG2-TS TRF burst (1 and $2*n$ for $n=1$ to 12).
$N_{p,atm}$	Number of bytes of the prefix used on ATM TRF bursts (0, 2 or 4).
$N_{p,sync}$	Number of bytes of the optional prefix used on SYNC bursts: 0, 2...31 for concatenated code, 0, 12 or 16 for the Turbo code (see subclause 6.4).

### 3.3 Abbreviations

For the purposes of the present document, the following abbreviations apply:

AAL	ATM Adaptation Layer
ACQ	Acquisition burst
ATM	Asynchronous Transfer Mode
AVBDC	Absolute Volume-Based Dynamic Capacity
BCD	Binary Coded Decimal
BTP	Burst Time Plan
CBC	Cipher Block Chaining

CMF	Control & Monitoring Functions
CMT	Correction Message Table
CR	Capacity Requests
CRA	Constant-Rate Assignment
CRC	Cyclic Redundancy Check
CRSC	Circular Recursive Systematic Convolutional
CSC	Common Signalling Channel
CTRL/MNGM	Control/Management virtual channel used in DULM
DES	Data Encryption Standard
DSM-CC	Digital Storage Medium - Command and Control
DULM	Data Unit Labelling Method
DVB	Digital Video Broadcast
DVB-S	Digital Video Broadcast by Satellite
EKE	Explicit Key Exchange
EN	European Norm
ETS	European Telecommunications Standard
ETSI	European Telecommunications Standards Institute
FCA	Free Capacity Assignment
FCT	Frame Composition Table
FLS	Forward Link Signalling
GEO	Geostationary Earth Orbit
GFC	Generic Flow Control
HMAC	Hash-based Message Authentication Code
HPA	High Power Amplifier
I	In-phase
ID	Identifier
IDU	Indoor unit
IE	Information Element
IEC	International Electrotechnical Commission
IEEE	Institute of Electrical and Electronics Engineers
IF	Intermediate Frequency
IP	Internet Protocol
ISDN	Integrated Services Digital Network
ISO	International Standards Organization
ITU	International Telecommunication Union
IV	Initialization Vector
LFSR	Linear Feedback Shift Register
LLC	Logical Link Control
LSB	Least Significant Bit
M&C	Monitoring and Control
MAC	Medium Access Control
MF-TDMA	Multiple-Frequency Time-Division Multiple Access
MIB	Management Information Base
MKE	Main Key Exchange
MPEG	Moving Picture Experts Group
MSB	Most Significant bit
NCC	Network Control Centre
NCR	Network Clock Reference
NIT	Network Information Table
NIU	Network Interface Unit
OBO	Output Back-Off
ODU	Outdoor unit
OSI	Open Systems Interconnection
PAT	Program Association Table
PC	Personal Computer
PCR	Program Clock Reference
PID	Packet IDentifier
PMT	Program Map Table
ppm	parts per million
PRBS	Pseudo Random Binary Sequence
PRNG	Pseudo-Random Number Generator
PSI	Program Specific Information

PSTN	Public Switched Telephone Network
PVC	Permanent Virtual Circuit
Q	Quadrature
QKE	Quick Key Exchange
QPSK	Quadrature Phase-Shift Keying
RBDC	Rate-Based Dynamic Capacity
RCST	Return Channel Satellite Terminal
RF	Radio Frequency
RMT	RCS Map Table
RS	Reed-Solomon
SAC	Satellite Access Control
SAR	Segmentation And Re-assembly
SCT	Superframe Composition Table
SDT	Service Description Table
SI	Service Information
SIT	Satellite Interactive Terminal
SMATV	Satellite Master Antenna Television
SNAP	Sub Network Access Protocol
SNMP	Simple Network Management Protocol
SPT	Satellite Position Table
SUT	Satellite User Terminal
SVC	Switched Virtual Circuit
SYNC	Synchronization burst type
TBTP	Terminal Burst Time Plan
TCT	Time-slot Composition Table
TDMA	Time-Division Multiple Access
TG	Traffic Gateway
TIM	Terminal Information Message
TRF	Traffic (burst type)
TS	Transport Stream
Tx	Transmitter
UNI	User Network Interface
VBDC	Volume-Based Dynamic Capacity
VCI	Virtual Circuit Identifier
VPI	Virtual Path Identifier

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## 4 Reference Models for Satellite Interactive Networks in DVB

### 4.1 Protocol Stack Model

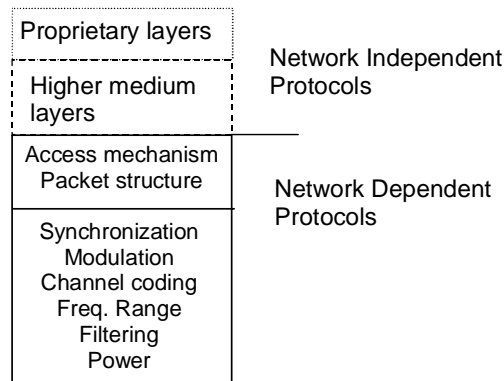
For interactive services supporting broadcast to the end user with return channel, a simple communications model consists of the following layers:

**physical layer:** where all the physical (electrical) transmission parameters are defined.

**transport layer:** defines all the relevant data structures and communication protocols like data containers, etc.

**application layer:** is the interactive application software and runtime environment (e.g. home shopping application, script interpreter, etc.).

A simplified model of the OSI layers was adopted to facilitate the production of specifications for these layers. Figure 1 points out the lower layers of the simplified model and identifies some of the key parameters for the lower two layers.



**Figure 1: Layer structure for generic system reference model**

The present document addresses the satellite interactive network dependent aspects only.

## 4.2 System Model

Figure 2 shows the system model which is to be used within DVB for interactive services.

In the system model, two channels are established between the service provider and the user:

- **Broadcast Channel:** a unidirectional broadband Broadcast Channel including video, audio and data is established from the service provider to the users. It may include the Forward Interaction Path.
- **Interaction Channel:** a bi-directional Interaction Channel is established between the service provider and the user for interaction purposes. It is formed by:
  - **Return Interaction Path (Return Channel):** from the user to the service provider. It is used to make requests to the service provider, to answer questions or to transfer data.
  - **Forward Interaction Path:** from the service provider to the user. It is used to provide information from the service provider to the user and any other required communication for the interactive service provision. It may be embedded into the Broadcast Channel. It is possible that this channel is not required in some simple implementations which make use of the Broadcast Channel for the carriage of data to the user.

The RCST is formed by the Network Interface Unit (consisting of the Broadcast Interface Module and the Interactive Interface Module) and the Set Top Unit. The RCST provides interface for both Broadcast and Interaction Channels. The interface between the RCST and the interaction network is via the Interactive Interface Module.

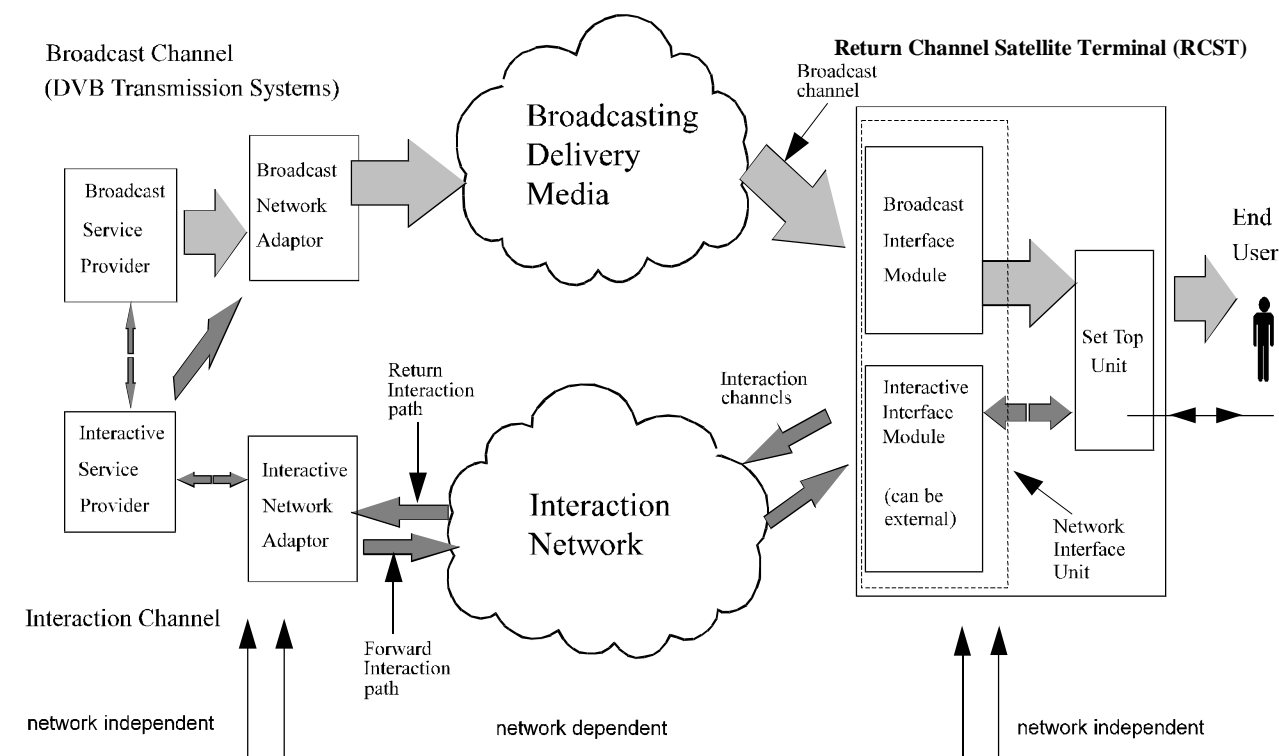


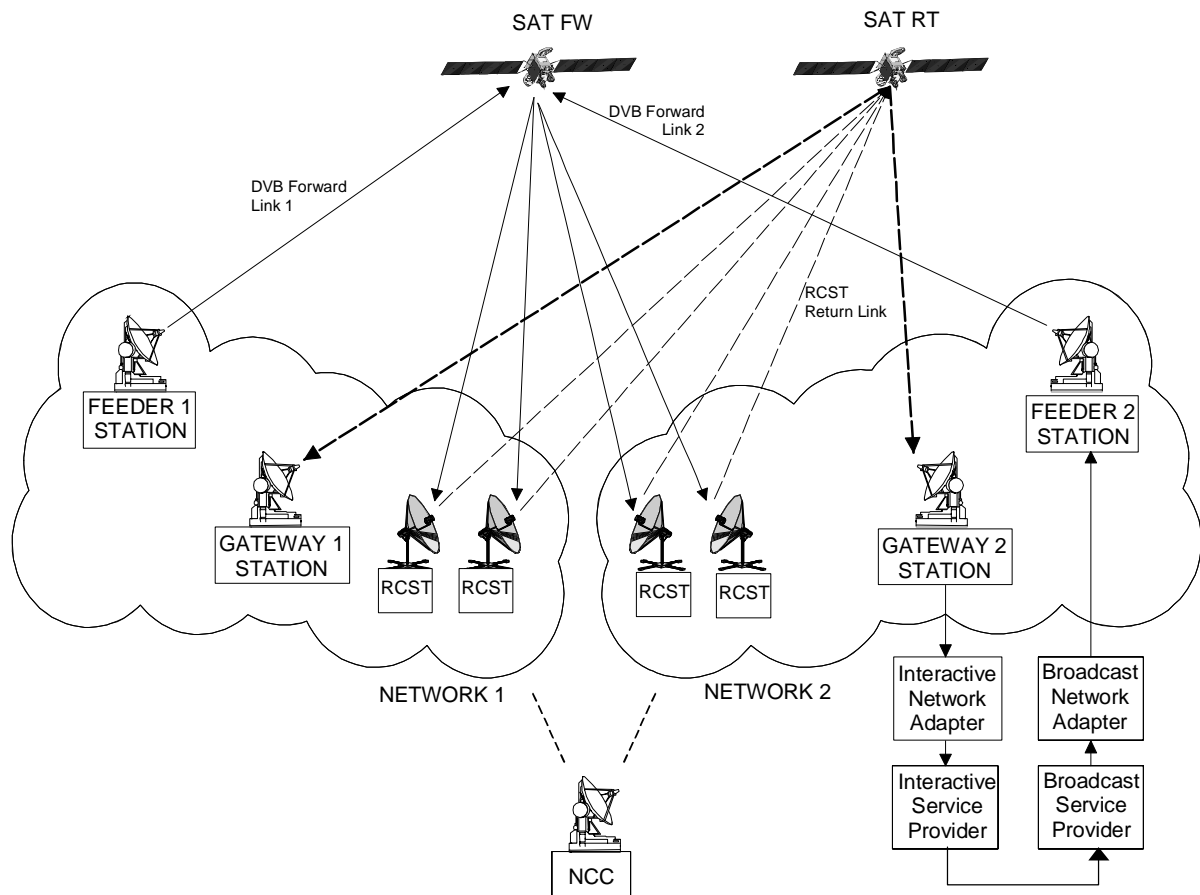
Figure 2: A generic system Reference Model for Interactive Systems

### 4.3 Reference Model of the Satellite Interactive Network

An overall Satellite Interactive Network, within which a large number of Return Channel Satellite Terminal (RCST) will operate, will comprise the following functional blocks, as shown in figure 3:

- **Network Control Centre:** a NCC provides monitoring & control functions. It generates control and timing signals for the operation of the Satellite Interactive Network to be transmitted by one or several Feeder Stations.
- **Traffic Gateway:** a TG receives the RCST return signals, provides accounting functions, interactive services and/or connections to external public, proprietary and private service providers (data bases, pay-per-view TV or video sources, software download, tele-shopping, tele-banking, financial services, stock market access, interactive games etc.) and networks (Internet, ISDN, PSTN etc...).
- **Feeder:** a Feeder transmits the forward link signal, which is a standard satellite digital video broadcast (DVB-S) uplink, onto which are multiplexed the user data and/or the control and timing signals needed for the operation of the Satellite Interactive Network.

An RCST is e.g. a SIT or a SUT as described in [6].



**Figure 3: Reference Model for the Satellite Interactive Network**

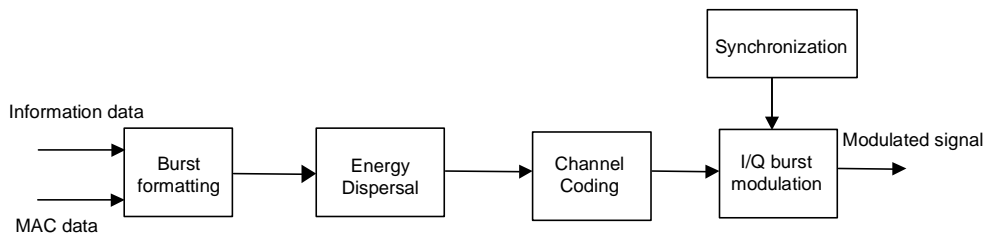
The forward link carries signalling from the NCC and user traffic to RCSTs. The signalling from the NCC to RCSTs that is necessary to operate the return link system is called "Forward Link Signalling" in the following. Both the user traffic and forward link signalling can be carried over different forward link signals. Several RCST configurations are possible depending on the number of forward link receivers present on the RCST.

## 5 Forward Link

The RCST shall be able to receive digital signals conforming to EN 300 421 [1], TR 101 202 [2], ETS 300 802 [3], EN 300 468 [4], EN 301 192 [5] and ETR 154 [10].

## 6 Return Link Base-band Physical Layer Specification and Multiple Access Definition

Specifications for the base-band physical layer are given in this chapter. Figure 4 represents the generic digital signal processing to be performed at the RCST transmitter side, from the burst formatting of the serial information bit-stream, to the modulation representing the digital to analogue conversion. The signal processing to be performed by each subset is described in the following sections.



**Figure 4: Block diagram of the RCST return link baseband signal processing**

## 6.1 RCST synchronization

### 6.1.1 Timing Control

The synchronization of the RCST is an important feature of the satellite interactive network. Constraints are imposed on the RCSTs to obtain an efficient TDMA system with minimum interference between users and maximum throughput, although they can be minimized if the NCC performs tasks such as satellite frequency translation error and common-mode Doppler compensation for RCST carrier frequency. For this reason, the synchronization scheme is based on information contained within the Forward Link Signalling as follows:

- NCR (Network Clock Reference);
- signalling in DVB/MPEG2-TS private sections.

The NCR is distributed with a specific PID within the MPEG2 Transport Stream that carries the Forward Link Signalling. The NCR distribution follows the PCR distribution mechanism as defined in ISO/IEC 13818-1 [9], which is usually derived from an MPEG video encoder, whereas here the NCR is derived from the NCC reference clock. The NCC reference clock will have an accuracy of 5 ppm or better.

### 6.1.2 Carrier synchronization

The MPEG2-TS that carries the Forward Link Signalling contains a NCR information which provides a 27 MHz reference of the NCC reference clock to the RCSTs. The RCST reconstructs the reference clock from the received NCR information as implemented in MPEG decoders for MPEG2 Transport Streams (MPEG2-TS). The RCST then performs a comparison to determine the offset between the local reference clock which controls the RCST upconverter local oscillator and the reference clock recovered from the received NCR. It then compensates the carrier frequency according to this offset. This local carrier synchronization provides a way of adjusting the transmit frequency of all RCSTs on the network to almost the same frequency.

Normalized carrier frequency accuracy shall be better than  $10^{-8}$  (root mean square).

### 6.1.3 Burst synchronization

The RCSTs retrieve the centre frequency, the start time and the duration of their transmit bursts by examining the forward link signalling (more precisely the SCT, FCT and TCT tables described in subclause 8.3.1).

The contention between RCSTs on the return link is resolved as described in subclause 6.7.

The bursts are sent according to the Burst Time Plan (BTP) received in the Forward Link Signalling (see subclause 6.7.2). The BTP is expressed in terms of centre frequency and absolute start time (given in NCR-counter value) of superframes and associated frequency and time offsets of burst allocations along with a description of the time slot properties. A superframe always starts at a given value of the RCST local NCR counter, which serves as a reference for all burst allocations within the superframe. For the purpose of synchronizing to the network the RCST reconstructs, in addition to the reference clock, the absolute value of the NCC reference clock. The RCST compares the reconstructed value with the NCR value given by the BTP. The time reference for counting timeslots occurs when the values are equal.

Burst synchronization accuracy shall be within 50% of a symbol period. The resolution shall be 1 NCR count interval. The burst synchronization accuracy is the worst case deviation of the scheduled start of burst time and the actual start of burst time at the transmitter output. The scheduled start of burst time is the point in time when the ideal reconstructed NCR equals the value written in the TBTP for that burst. The ideal reconstructed NCR is defined as observed at the output of an ideal delay-less DVB-S receiver. Compensation for the receiver delay, if required to achieve the specified accuracy, shall be done by the RCST.

### 6.1.4 Symbol clock synchronization

The symbol clock for the transmitter shall be locked to the NCR based clock, in order to avoid time drift with respect to the NCC reference clock. The RCST need not compensate for symbol clock Doppler.

Symbol clock accuracy shall be within 20 ppm from the nominal symbol\_rate value in the TCT (see subclause 8.5.5.4). The symbol clock rate shall have a short-term stability that limits the time error of any symbol within a burst to 1/20 symbol duration.

## 6.2 Burst Format

There are four types of bursts: traffic (TRF), acquisition (ACQ), synchronization (SYNC) and common signalling channel (CSC). The burst formats are described in the following.

### 6.2.1 Traffic (TRF) burst formats

Traffic (TRF) bursts are used for carrying useful data from the RCST to the Gateway. Two types of traffic bursts carrying either ATM cells or MPEG2-TS packets are defined herebelow. Channel coding of these bursts is defined in subclause 6.4. A TRF is usually followed by a guard time to decrease transmitted power and compensate for time offset as described in subclause 6.5.4.

#### 6.2.1.1 ATM TRF burst

The payload of an ATM traffic burst is composed of  $N_{\text{atm}}$  concatenated ATM cells, each of length 53 bytes, plus an optional  $N_{\text{p,atm}}$  byte prefix, as described in subclause 6.6.1.2. ATM cells follow the structure of an ATM cell but do not necessarily support ATM classes of service. See figure 5 for a description of the ATM TRF burst.

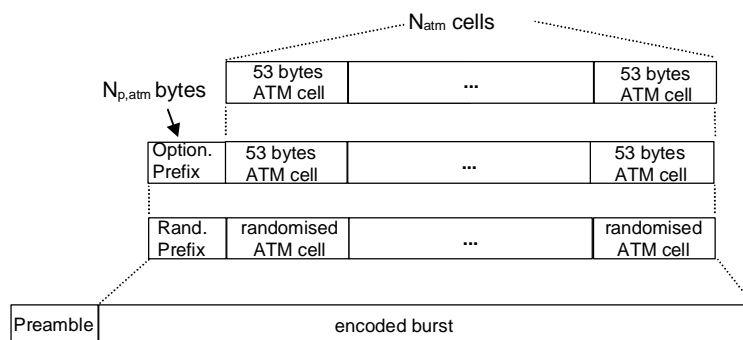


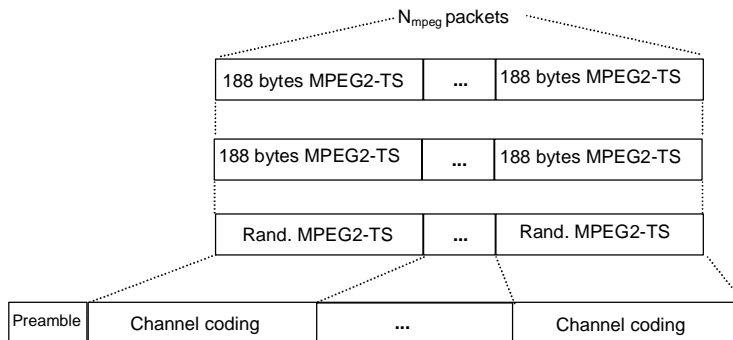
Figure 5: Composition of an ATM TRF burst

#### 6.2.1.2 Optional MPEG2-TS TRF burst

In the case that MPEG2-TS Packets are the basic containers a burst contains  $N_{\text{mpeg}}$  concatenated MPEG2-TS packets, each of length 188 bytes. The burst is composed of several channel coding blocks as described in subclause 6.4. See figure 6 for a description of the MPEG2-TS TRF burst.

RCSTs can deduce the number of MPEG2 packets in a TRF time slot from the time\_slot\_duration field of the TCT (see subclause 8.5.5.4), after subtracting the time duration of other fields. Transmission of MPEG2-TS TRF bursts is optional. The RCST will inform the NCC that it supports this mechanism in the CSC burst (subclause 6.2.3).





**Figure 6: Composition of the optional TRF burst carrying MPEG2-TS packets**

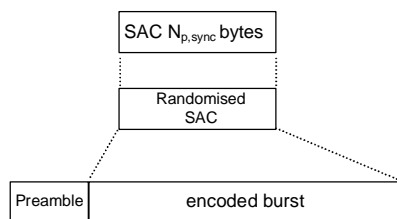
## 6.2.2 Synchronization and acquisition burst formats

Synchronization and Acquisition bursts are required to accurately position RCST burst transmissions during and after logon. Two separate burst types are defined for this purpose (SYNC and ACQ) as defined in the following paragraphs.

### 6.2.2.1 Synchronization (SYNC) burst format

A SYNC burst is used by an RCST for the purpose of maintaining synchronization and sending control information to the system. SYNC bursts are composed of a preamble for burst detection (configurable and indicated to the RCST through the TCT, as described in subclause 8.5.5.4), an optional  $N_{p, sync}$  byte Satellite Access Control (SAC) field as described in subclause 6.6.1.2, with the appropriate error control coding described in subclause 6.4. Like a TRF a SYNC is usually followed by a guard time to decrease transmitted power and compensate for time offset (see subclause 6.5.4). Figure 7 depicts the SYNC burst. The extent to which the SYNC burst is used depends on the capabilities of the NCC.

NOTE: SYNC bursts can be used in contention mode.

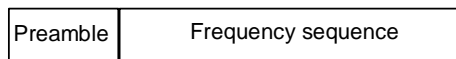


**Figure 7: Composition of a SYNC burst**

### 6.2.2.2 Acquisition (ACQ) burst

An ACQ burst can be used to achieve synchronization, prior to operational use of the network by the RCST.

Transmissions in an ACQ burst shall have the format shown on figure 8.

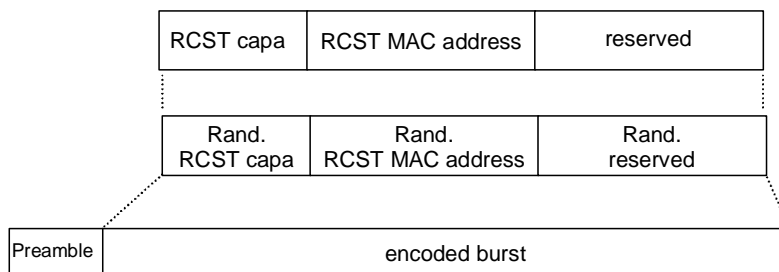


**Figure 8: Composition of the ACQ burst**

The preamble length and content (including the frequency sequence) are sent to the RCSTs via the TCT (see subclause 8.5.5.4). The ACQ is surrounded by a guard interval as described in subclause 6.5.4.

### 6.2.3 Common Signalling Channel (CSC) burst format

Common signalling channel (CSC) bursts are only used by an RCST to identify itself during logon. They are composed of a preamble for burst detection and start of burst detection, a field describing the RCST capabilities, the RCST MAC address and a reserved field (see figure 9). Coding of these bursts is defined in subclause 6.4. Table 1 gives the CSC burst content. The CSC is surrounded by a guard interval as described in subclause 6.5.4.



**Figure 9: Composition of a CSC burst**

**Table 1: CSC Burst Data Field Parameters**

Field Name	Size (bits)	Description/Content
Preamble	variable	Preamble for burst detection and start of burst detection. Definition by TCT. See subclause 8.5.5.4.
RCST Capability	24	See table 2
RCST MAC Address	48	RCST MAC address as per IEEE 802.3 [12]
Reserved	40	Reserved

The following table defines the different bit patterns within the RCST capability field. The 24-bit field is numbered from LSB to MSB using the notation b0 through b23.

**Table 2: RCST Capability**

Parameter	Bit Size	Description
Security mechanism	1 (b23)	'1' for RCST implementing security mechanism as described in subclause 9.4. '0' otherwise.
SNMP	1 (b22)	'1' for RCST supporting SNMP (see subclauses 8.4.2, 8.4.3 and 8.5.5.10.2). '0' otherwise.
ATM connectivity	1 (b21)	'1' for RCST capable of ATM connectivity (type B), '0' for not capable (Type A).
MPEG2-TS TRF	1 (b20)	'1' for RCST capable of MPEG2-TS TRF, '0' for not capable.
RCST boards	2 (b19-b18)	Number of RCST forward link receivers: '00' for 1 receiver, '01' for 2, '10' for more than 2, '11' reserved.
RCST ACQ	1 (b17)	'0' for RCST not requiring ACQ burst, '1' for ACQ required.
Multi_IDU	1 (b16)	'0' for single indoor unit/single outdoor unit configuration, '1' when two or more IDUs are connected to a single ODU.
S/W Version	8 (b15-b8)	System Dependent. Can be used to define the RCST software version.
Freq Hopping Range	2 (b7-b6)	Defines the RCST burst to burst frequency hopping range capability: '00' for 20MHz. Other patterns System Dependent
MF-TDMA	1 (b5)	'1' for RCST supporting dynamic MF-TDMA. '0' for RCST supporting fixed MF-TDMA (see subclause 6.7.1).
RCST Class	2 (b4-b3)	System Dependent
RCST Mode	3 (b2-b0)	'000' for Installation Mode (see subclause 8.5.5.10.5), '001' for Operational Mode, '010' for Reference RCST mode (can be used for measuring satellite frequency translation offset, downlink rain fade,...). Other patterns reserved.

## 6.2.4 Bit Numbering and Interpretation

For single-bit fields, a "1" shall represent "positive" information, such as "detected", "enabled", etc.

Multi-bit fields representing numerical values shall, as applicable, be unsigned or in two's complement representation.

The term "bit 0" shall refer to the least significant bit of a multi-bit field. The most significant bit of a k-bit unsigned field shall be designated "bit k-1". For a signed field, "bit k-1" shall be the sign bit and "bit k-2" the most significant magnitude-related bit.

## 6.2.5 Transmission Order

Fields in data structures shall be transmitted in the order in which they are defined.

Unsigned values shall be transmitted starting with the most significant bit and ending with the least significant bit.

Signed values shall be transmitted starting with the sign bit, followed by the most significant bit and ending with the least significant bit.

Bytes shall be processed MSB first.

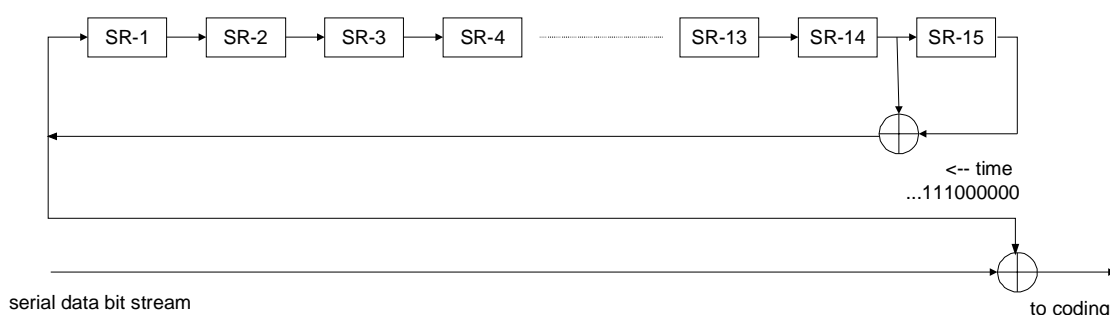
## 6.3 Randomization for energy dispersal

The return link data stream shall be organized in bursts as described in subclause 6.2. In order to comply with ITU Radio Regulations and to ensure adequate binary transitions, serial data bit stream in a burst shall be randomized. The polynomial of the Pseudo Random Binary Sequence (PRBS) shall be as the one of EN 300 421 [1] (see figure 10), i.e.  $1+x^{14}+x^{15}$ .

The data is randomized using the output of a 15 register linear feedback shift register (LFSR) randomized sequence (figure 10) to ensure a random distribution of ones and zeroes. The randomizer performs modulo-2 addition of the data with the pseudo-random sequence. The initial content of the SR-1 to SR-15 registers is given in table 3. The first bit of the pseudo-random sequence is to be added in the first bit of the serial data bit stream, i.e. the first bit after the burst preamble. The randomizer is reset to the initial content before processing the next burst.

**Table 3: Initial contents of the randomizer register**

Shift register	SR 1	SR 2	SR 3	SR 4	SR 5	SR 6	SR 7	SR 8	SR 9	SR 10	SR 11	SR 12	SR 13	SR 14	SR 15
Bit value	1	0	0	1	0	1	0	1	0	0	0	0	0	0	0



**Figure 10: Randomizer**

## 6.4 Coding

Coding for channel error protection is applied to traffic and control data, which are transmitted in the types of bursts described in subclause 6.2. Two coding schemes are described: Turbo (see subclause 6.4.4) and concatenated coding. RCST shall implement both schemes. Within a session (see clause 7), RCSTs are not requested to change the coding scheme (i.e. during a given session, an RCST will either use the Turbo or the concatenated code). In the case of the concatenated coding, the outer code is a by-passable Reed-Solomon (RS) code and the inner code is a by-passable non-systematic convolutional code (EN 300 421 [1]). For both coding schemes, a by-passable CRC can also be applied on CSC and SYNC bursts in order to allow error detection.

### 6.4.1 CRC error detection code

A CRC-16 can be applied on CSC and SYNC bursts in order to allow error detection. The CRC polynomial is  $x^{16}+x^{15}+x^2+1$ . The NCC indicates via the TCT (subclause 8.3.1.3) to the RCST if the CRC is to be applied. If used, the CRC is appended at the end of the burst before any other coding. CRC is applied on the randomized bit stream. The CRC is the remainder of the division of the burst payload by the polynomial. The CRC code is mandatory on turbo coded CSC bursts.

### 6.4.2 Reed Solomon outer coding

A Reed-Solomon RS(N-B, K-B, T) shortened code EN 300 421 [1] derived from the original RS(255, 239, 8) code, shortened by B bytes, can be applied for some burst formats. The code generator polynomial is  $g(x) = (x + \lambda^0)(x + \lambda^1)(x + \lambda^2) \dots (x + \lambda^{15})$ , where  $\lambda = 02_{\text{HEX}}$ . The generator polynomial is  $p(x) = x^8 + x^4 + x^3 + x^2 + 1$ . This code is similar to the one used in EN 300 421 [1].

For ATM traffic bursts the length of the encoded information word, K-B, is  $N_{\text{atm}} * 53 + N_{\text{p,atm}}$ . In the case that the basic container is an MPEG2-TS packet K - B = 188 applies.

The outer code can be bypassed. The outer code is always by-passed when using Turbo codes.

If both the CRC and RS codes are used, the burst payload CRC is first computed and the RS parity bytes are then added.

### 6.4.3 Convolutional inner coding

Processing of the convolutional encoder shall be in accordance with EN 300 421 [1], as summarized in the following.

The return link shall allow for a range of punctured convolutional codes, based on a rate 1/2 mother convolutional code with constraint length  $K = 7$  corresponding to 64 trellis states (see figure 11). The generator polynomials are  $G_0 = 171$  and  $G_1 = 133$  in octal representation. This choice will allow selection of the most appropriate level of error correction for a given service or data rate. Code rates of 1/2, 2/3, 3/4, 5/6 and 7/8 shall be supported. The inner code can be bypassed. In that case, the MSB bit is affected to the I channel, the next bit to the Q channel and so on. The convolutional inner code is always by-passed when using Turbo codes.

The encoder register shall be initialized to all zeroes before encoding the first data bit.

At the end of each data block, the encoder shall be flushed by 6 zero bits. This block is called the "Postamble". The output shall be continued until the encoder is in its all-zero state. If the inner code is bypassed, then the postamble is also omitted.

The puncturing pattern period counter shall be initialized before encoding the first data bit so that the first encoded (C2,C1) symbol always corresponds to an (Y1,X1) pair in table 4.

After encoding the last 0 bit of the postamble an incomplete symbol (##,C1) can remain if the message length is not divisible by the puncturing period. In that case the missing C2 is set to 0 and the burst is terminated.

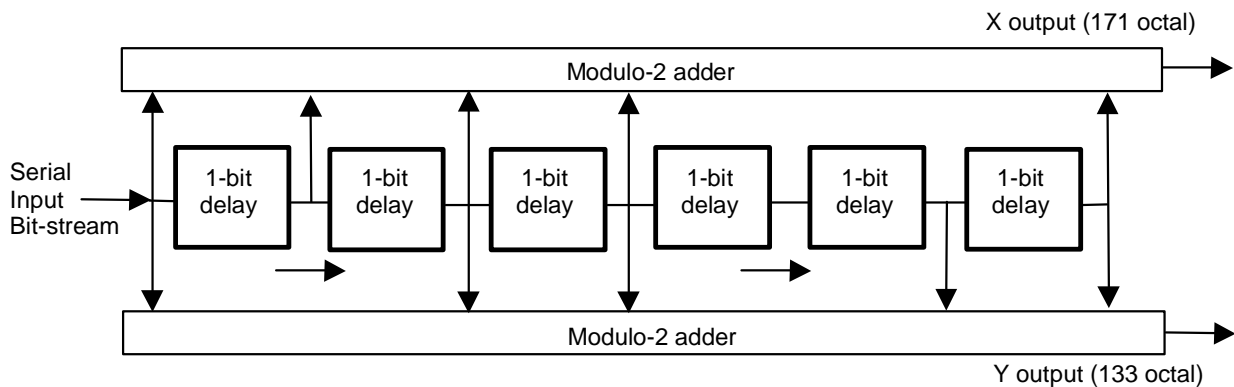


Figure 11: Convolutional code of rate 1/2

The punctured convolutional code shall be used as given in table 4, according to EN 300 421 [1].

Table 4: Punctured code definition

Original code			Code rates									
			1/2		2/3		3/4		5/6		7/8	
K	G <sub>1</sub> (X)	G <sub>2</sub> (Y)	P	d <sub>free</sub>	P	d <sub>free</sub>	P	d <sub>free</sub>	P	d <sub>free</sub>	P	d <sub>free</sub>
7	171 <sub>OC</sub> T	133 <sub>OC</sub> CT	X: 1 Y: 1	10	X: 1 0 Y: 1 1	6	X: 1 0 1 Y: 1 1 0	5	X: 1 0 1 0 1 Y: 1 1 0 1 0	4	X: 1 0 0 0 1 0 1 Y: 1 1 1 1 0 1 0	3
			C1=X <sub>1</sub> C2=Y <sub>1</sub>		C1= X <sub>1</sub> Y <sub>2</sub> Y <sub>3</sub> C2= Y <sub>1</sub> X <sub>3</sub> Y <sub>4</sub>		C1= X <sub>1</sub> Y <sub>2</sub> C2= Y <sub>1</sub> X <sub>3</sub>		C1= X <sub>1</sub> Y <sub>2</sub> Y <sub>4</sub> C2= Y <sub>1</sub> X <sub>3</sub> X <sub>5</sub>		C1= X <sub>1</sub> Y <sub>2</sub> Y <sub>4</sub> Y <sub>6</sub> C2= Y <sub>1</sub> Y <sub>3</sub> X <sub>5</sub> X <sub>7</sub>	

NOTE: 1 = transmitted bit.  
0 = non transmitted bit.

#### 6.4.4 Turbo code

The Turbo encoder is depicted in figure 12. It uses a double binary Circular Recursive Systematic Convolutional (CRSC) code. The MSB bit of the first byte after the burst preamble is assigned to A, the next bit to B and so on for the remaining of the burst content.

The encoder is fed by blocks of  $k$  bits or  $N$  couples ( $k = 2*N$  bits).  $N$  is a multiple of 4 ( $k$  is a multiple of 8).

The polynomials defining the connections are described in octal and symbolic notations as follows:

- for the feedback branch: 15 (in octal), equivalently  $1+D+D^3$  (in symbolic notation);
- for the Y parity bits: 13, equivalently  $1+D^2+D^3$ ;
- for the W parity bits: 11, equivalently  $1+D^3$ .

The input A bit is connected to tap "1" of the shift register and the input B bit is connected to the taps "1", D and D<sup>2</sup>.

First, the encoder (after initialization by the circulation state  $S_{C_1}$ , see subclause 6.4.4.2) is fed by the sequence in the natural order (switch on position 1) with incremental address  $i = 0, \dots, N-1$ . This first encoding is called C<sub>1</sub> encoding.

Then the encoder (after initialization by the circulation state  $S_{C_2}$ , see subclause 6.4.4.2) is fed by the interleaved sequence (switch in position 2) with incremental address  $j = 0, \dots, N-1$ . This second encoding is called C<sub>2</sub> encoding. The function  $\Pi(j)$  that gives the natural address  $i$  of the considered couple, when reading it at place  $j$  for the second encoding, is given in subclause 6.4.4.1.

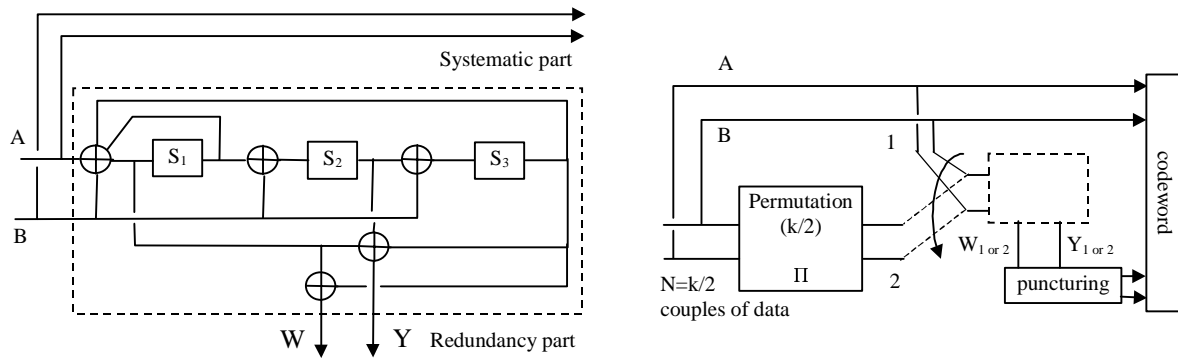


Figure 12: Encoder block diagram (turbo code)

#### 6.4.4.1 Description of the turbo code permutation

The permutation is done on two levels, the first one inside the couples (**level 1**), the second one between couples (**level 2**):

Set the permutation parameters  $P_0, P_1, P_2$  and  $P_3$

$$j = 0, \dots, N-1$$

##### level 1

if  $j \bmod 2 = 0$ , let  $(A,B) = (B,A)$  (invert the couple)

##### level 2

- if  $j \bmod 4 = 0$ , then  $P = 0$ ;
- if  $j \bmod 4 = 1$ , then  $P = N/2 + P_1$ ;
- if  $j \bmod 4 = 2$ , then  $P = P_2$ ;
- if  $j \bmod 4 = 3$ , then  $P = N/2 + P_3$ .

$$i = P_0 * j + P + 1 \bmod N$$

Table 5 provides the combinations of the default parameters to be used. Those parameters can be updated by the TCT (see subclause 8.5.5.4). The interleaving relations satisfy the odd/even rule (i.e. when  $j$  is even,  $i$  is odd and vice-versa) that enables the puncturing patterns to be identical for both encodings.

Table 5: Turbo code permutation parameters

Frame size in couples	$P_0$	$\{P_1, P_2, P_3\}$
$N = 48$ (12 bytes)	11	{24,0,24}
$N = 64$ (16 bytes)	7	{34,32,2}
$N = 212$ (53 bytes)	13	{106,108,2}
$N = 220$ (55 bytes)	23	{112,4,116}
$N = 228$ (57 bytes)	17	{116,72,188}
$N = 424$ (106 bytes)	11	{6,8,2}
$N = 432$ (108 bytes)	13	{0,4,8}
$N = 440$ (110 bytes)	13	{10,4,2}
$N = 848$ (212 bytes)	19	{2,16,6}
$N = 856$ (214 bytes)	19	{428,224,652}
$N = 864$ (216 bytes)	19	{2,16,6}
$N = 752$ (188 bytes)	19	{376,224,600}

### 6.4.4.2 Determination of the circulation states

The state of the encoder is denoted  $\mathbf{S}$  ( $0 \leq \mathbf{S} \leq 7$ ) with  $\mathbf{S} = 4.s_1 + 2.s_2 + s_3$  (see figure 12). The circulation states  $\mathbf{S}_{C1}$  and  $\mathbf{S}_{C2}$  are determined by the following operations:

1. initialize the encoder with state 0. Encode the sequence in the natural order for the determination of  $\mathbf{S}_{C1}$  or in the interleaved order for the determination of  $\mathbf{S}_{C2}$  (without producing redundancy). In both cases, the final state of the encoder is denoted  $\mathbf{S}_{N-1}^0$ ;
2. according to the length  $N$  of the sequence, use the following correspondence to find  $\mathbf{S}_{C1}$  or  $\mathbf{S}_{C2}$  (table 6).

**Table 6: Circulation state correspondence table**

$\mathbf{S}_{N-1}^0 \rightarrow$ $\downarrow N \bmod 7$	0	1	2	3	4	5	6	7
1	$\mathbf{S}_C=0$	$\mathbf{S}_C=6$	$\mathbf{S}_C=4$	$\mathbf{S}_C=2$	$\mathbf{S}_C=7$	$\mathbf{S}_C=1$	$\mathbf{S}_C=3$	$\mathbf{S}_C=5$
2	$\mathbf{S}_C=0$	$\mathbf{S}_C=3$	$\mathbf{S}_C=7$	$\mathbf{S}_C=4$	$\mathbf{S}_C=5$	$\mathbf{S}_C=6$	$\mathbf{S}_C=2$	$\mathbf{S}_C=1$
3	$\mathbf{S}_C=0$	$\mathbf{S}_C=5$	$\mathbf{S}_C=3$	$\mathbf{S}_C=6$	$\mathbf{S}_C=2$	$\mathbf{S}_C=7$	$\mathbf{S}_C=1$	$\mathbf{S}_C=4$
4	$\mathbf{S}_C=0$	$\mathbf{S}_C=4$	$\mathbf{S}_C=1$	$\mathbf{S}_C=5$	$\mathbf{S}_C=6$	$\mathbf{S}_C=2$	$\mathbf{S}_C=7$	$\mathbf{S}_C=3$
5	$\mathbf{S}_C=0$	$\mathbf{S}_C=2$	$\mathbf{S}_C=5$	$\mathbf{S}_C=7$	$\mathbf{S}_C=1$	$\mathbf{S}_C=3$	$\mathbf{S}_C=4$	$\mathbf{S}_C=6$
6	$\mathbf{S}_C=0$	$\mathbf{S}_C=7$	$\mathbf{S}_C=6$	$\mathbf{S}_C=1$	$\mathbf{S}_C=3$	$\mathbf{S}_C=4$	$\mathbf{S}_C=5$	$\mathbf{S}_C=2$

### 6.4.4.3 Rates and puncturing map

Seven code rates are defined for the Turbo mode:  $R = 1/3, 2/5, 1/2, 2/3, 3/4, 4/5, 6/7$ . This is achieved through selectively deleting the parity bits (puncturing). The puncturing patterns of table 7 are applied. These patterns are identical for both codes  $C_1$  and  $C_2$  (deleting is always done in couples). The puncturing rate is indicated to the RCSTs via the TCT (see subclause 8.5.5.4).

Rates  $1/3, 2/5, 1/2, 2/3$  and  $4/5$  are exact ones, independently of the block size. Rates  $3/4$  and  $6/7$  are exact ones only if  $N$  is a multiple of 3. In other cases, the actual rates are very slightly lower than the nominal ones.

Depending on the code rate, the length of the encoded block is:

- $2N + M$  for  $R < 1/2$ , with:
  - $M = N$  for  $R = 1/3$ ;
  - $M = N/2$  for  $R = 2/5$ .
- $N + M$  for  $R \geq 1/2$ , with:
  - $M = N$  for  $R = 1/2$ ;
  - $M = N/2$  for  $R = 2/3$ ;
  - for  $R = 3/4$ :
    - $M = N/3$  (if  $N \bmod 3 = 0$ ); or
    - $M = (N-4)/3+2$  (if  $N \bmod 3 = 1$ ); or
    - $M = (N-8)/3+3$  (if  $N \bmod 3 = 2$ ).
  - $M = N/4$  for  $R = 4/5$ ;
  - for  $R = 6/7$ :
    - $M = N/6$  (if  $N \bmod 3 = 0$ ); or
    - $M = (N-4)/6+1$  (if  $N \bmod 3 = 1$ ); or
    - $M = (N-8)/6+2$  (if  $N \bmod 3 = 2$ ).

**Table 7: Puncturing patterns for Turbo codes. "1"=keep**

$$\begin{array}{cc}
 \frac{1}{3} \begin{array}{l} Y \\ W \end{array} \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix} & \frac{2}{5} \begin{array}{l} Y \\ W \end{array} \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 \end{bmatrix} \\
 \\
 \frac{1}{2} \begin{array}{l} Y \\ W \end{array} \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} & \frac{2}{3} \begin{array}{l} Y \\ W \end{array} \begin{bmatrix} 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} \\
 \\
 \frac{3}{4} \begin{array}{l} Y \\ W \end{array} \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} & \frac{4}{5} \begin{array}{l} Y \\ W \end{array} \begin{bmatrix} 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} \\
 \\
 \frac{6}{7} \begin{array}{l} Y \\ W \end{array} \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}
 \end{array}$$

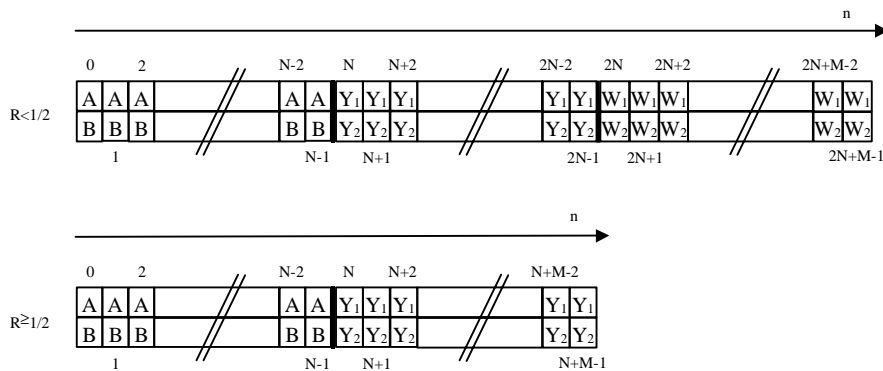
**6.4.4.4 Order of transmission and mapping to QPSK constellation**

Two orders of transmission are allowed:

- in the natural order, all couples (A,B) are transmitted first, followed by all couples (Y<sub>1</sub>,Y<sub>2</sub>) that remain after puncturing and then all couples (W<sub>1</sub>,W<sub>2</sub>) that remain after puncturing (see figure 13);
- in the reverse order, the couples (Y<sub>1</sub>,Y<sub>2</sub>) are transmitted first, in their natural order, followed by the couples (W<sub>1</sub>,W<sub>2</sub>), if any, and then finally by the couples (A,B).

Each couple is mapped to one QPSK constellation point as shown in figure 15. In figure 13, the row with the A symbols is mapped on the I channel (C<sub>1</sub> in figure 15).

The order of transmission is signalled by the NCC as an inner code parameter in the Time Slot Composition Table (subclause 8.5.5.4).



**Figure 13: Encoded blocks (natural order)**



## 6.5 Modulation

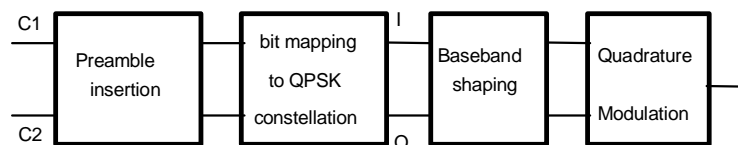
The signal shall be modulated using QPSK, with baseband shaping as described here below.

### 6.5.1 Bit mapping to QPSK constellation

Mapping into the QPSK constellation shall be as follows.

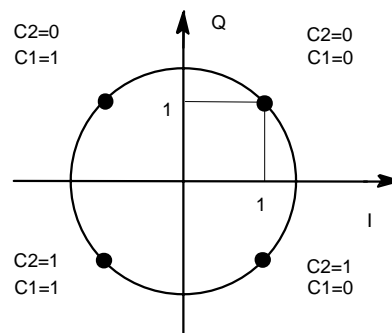
The preamble shall be configurable and indicated to the RCST through the TCT, as described in subclause 8.5.5.4.

Immediately after the preamble insertion, the outputs C1 and C2 of the encoder shall be sent without modification to the QPSK bit mapper (see figure 14).



**Figure 14: Processing after the encoder**

Conventional Gray-coded QPSK modulation with absolute mapping (no differential coding) shall be used. Bit mapping in the QPSK constellation shall follow figure 15. If the normalization factor  $1/\sqrt{2}$  is applied to the I and Q components, the corresponding average energy per symbol becomes equal to 1.



**Figure 15: Bit mapping into QPSK constellation**

The output C1 of the channel coding shall be mapped to the I channel of the modulation. The output C2 shall be mapped to the Q channel of the modulation.

Following the modulation process, the spectrum of the QPSK signal shall not be inverted.

### 6.5.2 Baseband shaping and quadrature modulation

Prior to modulation, the I and Q signals (mathematically represented by a succession of Dirac delta functions, multiplied by the amplitudes I and Q, spaced by the symbol duration  $TS = 1/RS$ ) shall be square root raised cosine filtered. The roll-off factor shall be 35 % as indicated by the TCT (see subclause 8.5.5.4). The baseband square root raised cosine filter shall have a theoretical function defined by the following expression:

$$\begin{aligned}
 H(f) &= 1 && \text{for } |f| < f_N(1 - \alpha) \\
 H(f) &= \sqrt{\frac{1}{2} + \frac{1}{2} \sin \frac{\pi(f_N - |f|)}{2\alpha f_N}} && \text{for } f_N(1 - \alpha) \leq |f| \leq f_N(1 + \alpha) \\
 H(f) &= 0 && \text{for } |f| > f_N(1 + \alpha)
 \end{aligned}$$

where  $f_N = \frac{1}{2T_s} = \frac{R_s}{2}$  is the Nyquist frequency and  $\alpha$  is the roll-off factor.

At the RCST antenna output (using a large output back-off), the allowed spectrum and group delay variation shall be according to the mask given in EN 300 421 [1] for every return link symbol rate supported by the terminal.

### 6.5.3 Output power control

In order to avoid degradation of the spectrum the RCST High Power Amplifier (HPA) may be operated with Output Power Backoff (OBO). The RCST shall have an output power accuracy of 0,5 dB in the linear range of the HPA with a power adjustment whose granularity is 0,5 dB (see subclauses 8.5.5.9 and 8.5.5.10.3).

### 6.5.4 Guard Time

Each burst is surrounded by a guard time which allows for RCST power switch-off transient and system timing errors. The guard time is allocated by the NCC as an implicit element of the TCT. The guard time is network dependant and is determined by overall system requirements. On optional MPEG2-TS TRF bursts, the guard time shall be shorter than half an MPEG2-TS packet.

## 6.6 MAC messages

All methods described below can be used by RCSTs for capacity requests and M&C messages. One or more of the methods may be employed in a Satellite Interactive Network. For the particular implementation, the RCSTs are configured at the time of logon by the logon initialize descriptor (subclause 8.5.5.10.4) that is transmitted in a TIM.

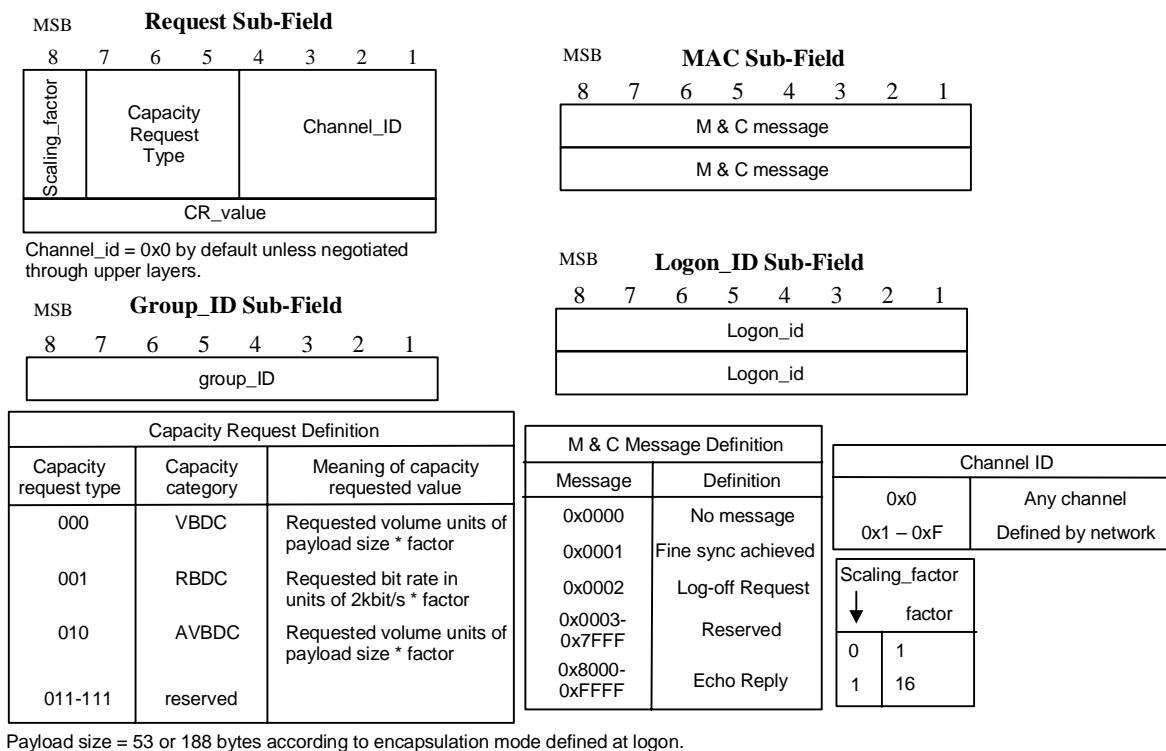
### 6.6.1 Prefix Method

#### 6.6.1.1 Prefix method mechanism

This mechanism is based on an optional  $N_{p,atm}$  bytes prefix attached to ATM traffic bursts. If used, the prefix carries control and management information from the RCSTs to the NCC, mainly composed of capacity requests. This mechanism is supported by the SAC request sub-field (see subclause 6.6.1.2) when appended to ATM Traffic bursts (see subclause 6.2.1.1).

#### 6.6.1.2 Satellite Access Control (SAC) field

The SYNC and the optional prefix attached to ATM TRF bursts contain the Satellite Access Control (SAC) field composed of signalling information added by the RCST for the purpose of requesting capacity on the session, or other additional MAC information. The SAC is composed of optional sub-fields that are defined in figure 16. The SAC field configuration for specific time slots is signalled by the NCC in the TCT as described in subclause 8.5.5.4.



**Figure 16: Composition of the SAC field**

Capacity requests for the different capacity categories are described in subclause 6.8. The Fine sync achieved message is used by the RCST to indicate that it has completed fine uplink synchronization on the return link and is ready to transmit traffic. The Logoff request message is used by the RCST to initiate logoff. The Echo Reply is a maintenance feature. The NCC can request the RCST echo back a predetermined sequence in this SAC for trouble-shooting purposes. The NCC would put this request in a TIM (see subclause 8.5.5.10.8). The group\_ID and logon\_ID sub-fields enable the NCC to identify the RCST that has transmitted a burst in a contention timeslot. When several of these SAC sub-fields are present, they appear in the following order: requests, M&C, group\_ID and logon\_ID (this order is identical to the order of flags that define the SAC in the TCT - subclause 8.5.5.4). If the SAC field length is larger than the sum of its sub-field lengths, then stuffing bytes shall follow after the last sub-field.

If the RCST does not have a message to send, it shall use the 'No message' in the M&C field. If the RCST does not have any capacity request to send, it shall send a VBDC request with an amount of 0.

## 6.6.2 Data unit labelling method (DULM)

The Data Unit Labelling method (DULM) is a "message-based" method that allows RCSTs to transmit control and/or management information to the NCC in the payload of the Data Units already assigned to them in TRF bursts. These Data Units are either ATM cells or MPEG TS packets, depending on the encapsulation mode of the RCST, as described in subclause 6.2.1.

The sequence of these "special" TRF bursts from an RCST to the NCC constitutes a dedicated virtual channel named CTRL/MNGM. This virtual channel can thus be utilized to support message-based delivery of control (e.g. bandwidth requests) and / or management information, possibly in combination with other methods described in this subclause.

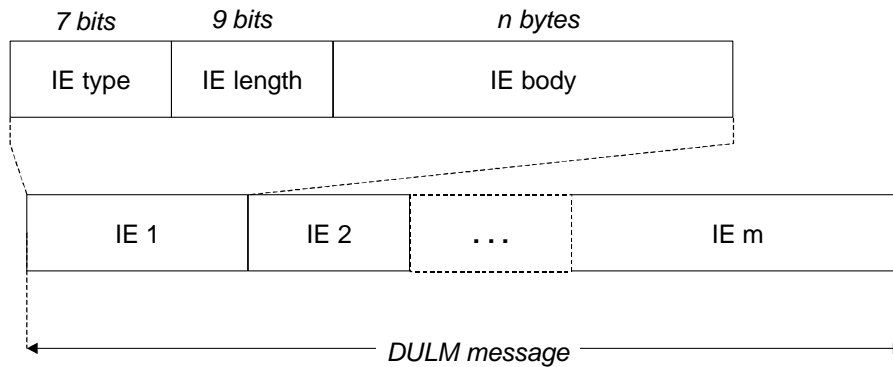
The use of the CTRL/MNGM virtual channel may be at the initiative of the NCC or of the RCST. For instance, an active RCST may insert a bandwidth request in an already assigned timeslot. The NCC may request status information from an RCST, and allocate timeslots to a terminal for the CTRL/MNGM channel.

### 6.6.2.1 DULM with ATM-formatting

For RCSTs, using ATM TRF bursts as defined in subclause 6.2.1.1, the CTRL/MNGM virtual channel shall be identified by a unique value in the Data Unit header, with the MSB of the PTI field set to 1 so as to allow discrimination from normal traffic information (see subclause 8.1.1).

A DULM message shall be composed of an integer number (between 1 and 64) of Information Elements (IE).

Each IE shall have the format described in figure 17, and be made of 2 bytes of header (IE type and IE length) plus a body of n bytes, with n between 1 and 512. A given DULM message may be composed of IEs of different lengths.



**Figure 17: DULM message with ATM TRF bursts**

The DULM message shall then be transmitted using standard AAL5 mechanisms, as specified in ITU-T Recommendation I.363-5 [18].

The Information Elements shall be as defined in table 8.

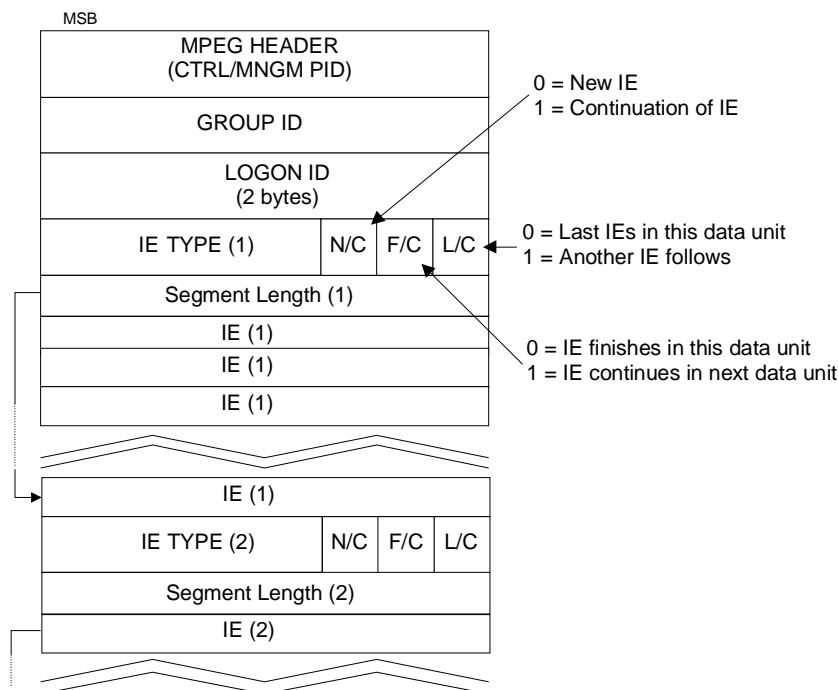
**Table 8: IEs for the ATM TRF case**

IE type	IE length*	IE body	
0x00	Capacity Request	2 bytes	As per subclause 6.6.1.2
0x01	M&C	2 bytes	As per subclause 6.6.1.2
0x02	Group_and_Logon_ID	3 bytes	As per subclause 6.6.1.2**
0x03 - 0x0C	Reserved		
0x0D	Security Sign-on Response	8 bytes	As per subclause 9.4.9.2
0x0E – 0x1E	Reserved		
0x1F	Wait		As per subclause 9.4.9.9
0x20 - 0x30	Reserved		
0x31	Main Key Exchange Response		As per subclause 9.4.9.4
0x32	Reserved		
0x33	Quick Key Exchange Response		As per subclause 9.4.9.6
0x34	Reserved		
0x35	Explicit Key Exchange Response		As per subclause 9.4.9.8
0x36 - 0x7F	Reserved		
NOTE 1: * "IE length": length (in bytes) of the IE body.			
NOTE 2: ** Group_and_Logon_ID: concatenation of the 1-byte Group_ID and the 2-byte Logon_ID fields, in this order.			

### 6.6.2.2 DULM with MPEG-formatting

For the RCSTs using the optional MPEG TRF bursts, a CTRL/MNGM PID shall be used in the header of CTRL/MNGM bursts. This PID is obtained by the RCST during the logon procedure (see subclause 8.5.5.10.4).

DULM messages shall use the transport mechanism described below, and illustrated in figure 18.



**Figure 18: DULM over optional MPEG packets**

The first 3 bytes of each Data Unit payload shall contain the RCST Group\_ID and Logon\_ID. Then one or more Information Elements (IE) follow with the structure described hereafter.

The first byte of an IE contains 4 fields:

- the [IETYPE] field is contained in the first 5 bits. It contains the IE type, as defined in table 9;
- the [N/C] is contained in the 6th bit. If [N/C]=0 the IE is starting in this Data Unit, if [N/C]=1 the IE is continued from the previous CTRL/MNGM Data Unit from the same RCST;
- the [F/C] is contained in the 7th bit. If [F/C]=0 the IE will finish in this CTRL/MNGM Data Unit, if [F/C]=1 the IE will continue in the next CTRL/MNGM Data Unit from the same RCST;
- the [L/C] is contained in the 8th bit. If [L/C]=0 the IE is the last of the CTRL/MNGM Data Unit, if [L/C]=1 another IE follows in the same CTRL/MNGM Data Unit.

The second byte contains the [SLENG] field. It indicates the segment length in number of bytes, defined as the part of the IE contained in the current Data Unit.

The following bytes contain the variable-length fragment.

**NOTE:** If the IE spans over several CTRL/MNGM Data Units, the IE header (IETYPE, [N/C], [F/C], [L/C], segment\_length) is duplicated on all Data Units.

Padding bytes set to all '0's shall be appended to the last IE of a CTRL/MNGM data unit.

Table 9: IEs for the MPEG TRF case

IE type (MPEG)		IE length	IE body
0x00	Capacity Request	2 bytes	As per subclause 6.6.1.2
0x01	M&C	2 bytes	As per subclause 6.6.1.2
0x02 - 0x0C	Reserved		
0x0D	Security Sign-on Response	8 bytes	As per subclause 9.4.9.2
0x0E - 0x10			
0x11	Main Key Exchange Response		As per subclause 9.4.9.4
0x12	Reserved		
0x13	Quick Key Exchange Response		As per subclause 9.4.9.6
0x14	Reserved		
0x15	Explicit Key Exchange Response		As per subclause 9.4.9.8
0x16 - 0x1E	Reserved		
0x1F	Wait		As per subclause 9.4.9.9

### 6.6.3 Mini-slot Method

This mechanism is based on a periodic assignment to logged-on RCSTs of bursts smaller than traffic timeslots. It carries control and management information from the RCSTs to the NCC and is used also for maintaining RCST synchronization. This mechanism is supported by the SAC request sub-field (see subclause 6.6.1.2) used in SYNC bursts (see subclause 6.2.2.1).

### 6.6.4 Contention based Mini-slot Method

As per the method described in the previous subclause, but the mini-slot can be accessed by a group of RCSTs on a contention basis. This mechanism is supported by the SAC request, group\_id and logon\_id sub-fields (see subclause 6.6.1.2) used in SYNC bursts (see subclause 6.2.2.1).

## 6.7 Multiple Access

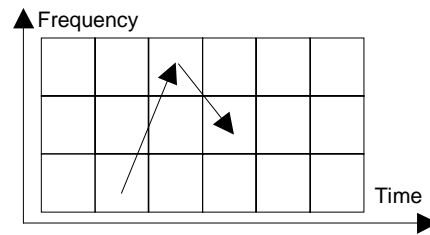
The multiple-access capability is either fixed or dynamic slot MF-TDMA. RCSTs shall indicate their capability by using the MF-TDMA field present on the CSC burst (see subclause 6.2.3).

### 6.7.1 MF-TDMA

The satellite access scheme is Multi-Frequency Time Division Multiple Access (MF-TDMA). MF-TDMA allows a group of RCSTs to communicate with a gateway using a set of carrier frequencies, each of which is divided into time-slots. The NCC will allocate to each active RCST a series of bursts, each defined by a frequency, a bandwidth, a start time and a duration.

#### 6.7.1.1 Fixed MF-TDMA

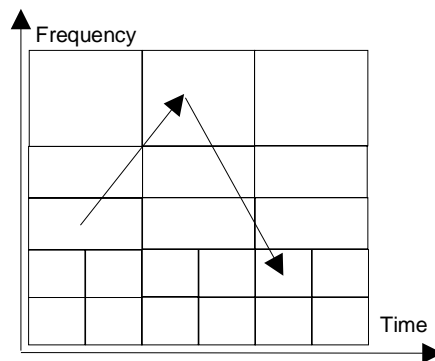
In Fixed-Slot MF-TDMA, the bandwidth and duration of successive traffic slots used by an RCST is fixed, as illustrated in figure 19 where the arrow indicates a typical sequence of slots assigned by the NCC to one RCST. In this case, TCT parameters (see subclause 8.5.5.4) defining the burst parameters (symbol\_rate, inner\_code\_type, inner\_code\_ordering, outer\_coding, inner\_code\_puncturing, modulation and baseband shaping) of a superframe are fixed. A fixed MF-TDMA RCST can send a mix of SYNC and single size TRF bursts provided that the burst parameters fulfil the previous requirement. If the NCC requests a change in these parameters, then they will apply to a new superframe with a delay as described in subclause 6.7.2.3.



**Figure 19: Fixed-slot MF-TDMA**

### 6.7.1.2 Dynamic MF-TDMA (Optional)

Dynamic-Slot MF-TDMA uses additional RCST flexibility to vary the bandwidth and duration of successive slots allocated to an RCST. In addition to changing carrier frequency and burst duration, the RCST may also change transmission rate and coding rate between successive bursts. The advantage of the more flexible RCST is more efficient adaptation to the widely varying transmission requirements typical of multimedia. The basic principle of the flexible RCST is illustrated in figure 20, where the arrows show an RCST using successive slots with different bandwidths and durations.



**Figure 20: Optional dynamic-Slot MF-TDMA, using a flexible RCST**

### 6.7.1.3 Frequency range

RCSTs have a specific frequency range for the frequency hopping from time-slot to time-slot. This frequency hopping range is communicated from an RCST to the NCC in a CSC burst during logon procedure (see subclause 6.2.3). The frequency hopping capability of individual RCSTs shall be at the manufacturer's discretion and shall be at least 20 MHz (i.e.  $\pm 10$  MHz around centre frequency).

The frequency agility is specified in terms of short term frequency hopping and long term frequency tuning.

Frequency hopping covers changes of frequency between adjacent slots in time. The settling time between bursts shall not exceed the guard interval defined in subclause 6.5.4. The settling time is the period required to achieve the frequency tolerance given in subclause 6.1.2.

Frequency tuning covers the change in centre frequency of hop ranges. The settling time shall not exceed 1 second.

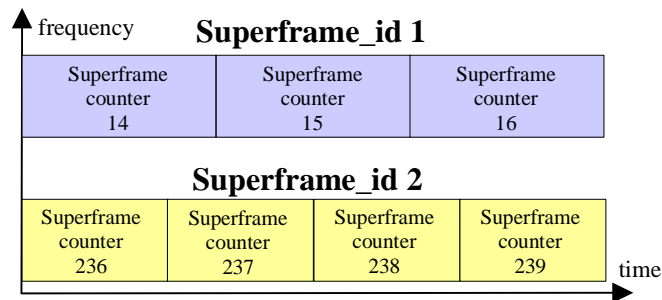
## 6.7.2 Segmentation of the return link capacity

In a Satellite Interactive Network, the timeslots of the return link are organized and numbered so that the network is able to allocate them to individual RCSTs.

### 6.7.2.1 Superframes

A superframe is a portion of time and frequency of the return link.

Within a Satellite Interactive Network, a Superframe\_ID identifies the return link resources accessed by a given set of RCSTs. Figure 21 shows a typical example whereby superframe\_IDs are indeed separate sets of carrier frequencies.



**Figure 21: Typical example of superframes of a Satellite Interactive Network**

In a Satellite Interactive Network, the global return link capacity may be segmented amongst sets of RCSTs, and the network will then separately manage several Superframe\_IDs. In the following, we only consider one Superframe\_ID.

As shown by figure 21, the consecutive superframes of a given Superframe\_ID are contiguous in time. Each occurrence of a superframe in time is labelled with a number called "superframe\_counter".

For each superframe (of a given Superframe\_ID), allocation of timeslots is communicated to the RCSTs via the TBTP table (see subclause 8.5.5.7). An RCST is allowed to transmit bursts only in timeslots which were allocated to it ("dedicated access"), or on random-access timeslots ("contention access").

**NOTE 1:** Some timeslots (like the ACQ and the SYNC bursts, see subclauses 6.2.2.1 and 6.2.2.2) can be assigned to RCSTs on the basis of a period which is much longer than the superframe via individual TIM messages (see subclauses 8.5.5.10.5 and 8.5.5.10.6). The period for these timeslots will be system dependent, but typically in the order of one second.

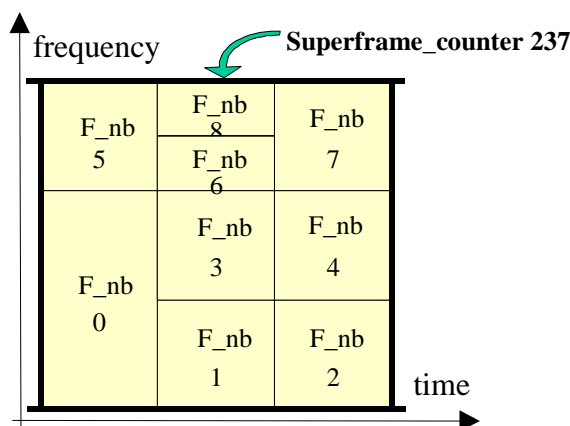
The superframe duration is therefore the elementary period of time for the assignment of resources to terminals.

**NOTE 2:** Some timeslots (like the SYNC bursts, see subclause 6.2.2.1) can be assigned to RCSTs on the basis of a period which is much longer than the superframe duration, in the order of one second (system-dependent).



### 6.7.2.2 Frames

A superframe is composed of frames, themselves composed of timeslots. The frame is at an intermediate level between the superframe and the timeslots. It is introduced for reasons of signalling efficiency (on the forward link signalling). The frame duration is not used as the basis of any timeslot allocation process.



**Figure 22: example of superframe composition**

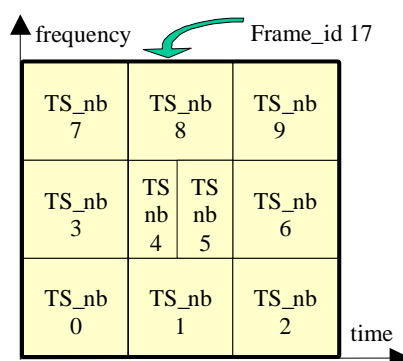
In a superframe, frames are numbered from 0 (lowest frequency, first in time) to N (highest frequency, last in time), ordered in time then in frequency as shown in figure 22. N shall be less than or equal to 31.

Frames of a superframe may not all have the same duration, bandwidth and timeslot composition.

Frames and superframes may all have the same duration, in which case frames can be seen as frequency sub-bands of the superframe. Anyway this property is not mandatory, and figure 22 shows an example of one superframe lasting 3 times more than each of its frames.

### 6.7.2.3 Timeslots

A frame is composed of timeslots (see the Frame Composition Table, subclause 8.5.5.3). A "frame\_id" identifies a particular arrangement of timeslots. For example, frame\_id = 1 could identify a sequence of 10 "user traffic" timeslots on the same carrier, and frame\_id = 2 a sequence of 4 "control" timeslots followed by 8 "user traffic" timeslots, all on the same carrier.



**Figure 23: example of frame composition**

### Figure 24: Void

A frame may span over several carrier frequencies. In a frame, timeslots are numbered from 0 (lowest frequency, first in time) to M (highest frequency, last in time), ordered in time then in frequency as shown in figure 23. M shall be less than or equal to 2047.

For the purpose of allocation, each timeslot is uniquely identified by its Superframe\_ID, Superframe\_counter, Frame\_number and Timeslot\_number.

The RCST shall process the TBTP message from the NCC for its allocation area, to extract the assignment count and timeslot allocations for its next uplink transmissions. The latency time from the arrival of the TBTP message at the RCST until the RCST is ready to transmit the bursts assigned by that TBTP shall not exceed 90 ms.

## 6.8 Capacity request categories

The timeslot allocation process shall support five capacity categories:

- Continuous Rate Assignment (CRA);
- Rate Based Dynamic Capacity (RBDC);
- Volume Based Dynamic Capacity (VBDC);
- Absolute Volume Based Dynamic Capacity (AVBDC);
- Free Capacity Assignment (FCA).

### 6.8.1 Continuous Rate Assignment (CRA)

CRA is rate capacity which shall be provided in full for each and every superframe while required. Such capacity shall be negotiated directly between the RCST and the NCC.

### 6.8.2 Rate Based Dynamic Capacity (RBDC)

RBDC is rate capacity which is requested dynamically by the RCST. RBDC capacity shall be provided in response to explicit requests from the RCST to the NCC, such requests being absolute (i.e. corresponding to the full rate currently being requested). Each request shall override all previous RBDC requests from the same RCST, and shall be subject to a maximum rate limit negotiated directly between the RCST and the NCC.

To prevent a terminal anomaly resulting in a hanging capacity assignment, the last RBDC request received by the NCC from a given terminal shall automatically expire after a time-out period whose default value is 2 superframes, such expiry resulting in the RBDC being set to zero rate. The time-out can be configured between 1 and 15 superframes (if set to 0 the time out mechanism is disabled) by the optional mechanism of subclause 8.4.2.

CRA and RBDC can be used in combination, with CRA providing a fixed minimum capacity per superframe and RBDC giving a dynamic variation component on top of the minimum.

### 6.8.3 Volume Based Dynamic Capacity (VBDC)

VBDC is volume capacity which is requested dynamically by the RCST. VBDC capacity shall be provided in response to explicit requests from the RCST to the NCC, such requests being cumulative (i.e. each request shall add to all previous requests from the same RCST). The cumulative total per RCST shall be reduced by the amount of this capacity category assigned in each superframe.

### 6.8.4 Absolute Volume Based Dynamic Capacity (AVBDC)

AVBDC is volume capacity which is requested dynamically by the RCST. This VBDC capacity shall be provided in response to explicit requests from the RCST to the NCC, such requests being absolute (i.e. this request replaces the previous ones from the same RCST). The AVBDC is used instead of VBDC when the RCST senses that the VBDC request might be lost (for example in the case of contention minislots).

### 6.8.5 Free Capacity Assignment (FCA)

FCA is volume capacity which shall be assigned to RCSTs from capacity which would be otherwise unused. Such capacity assignment shall be automatic and shall not involve any signalling from the RCST to the NCC. It shall be possible for the NCC to inhibit FCA for any RCST or RCSTs.

FCA should not be mapped to any traffic category, since availability is highly variable. Capacity assigned in this category is intended as bonus capacity which can be used to reduce delays on any traffic which can tolerate delay jitter.

## 7 Synchronization procedures

This clause defines the procedures to allow an RCST to logon to the satellite interactive network and also how the terminal can logoff from (or be logged off by) the network.

The period of time from the terminal logon to the terminal logoff is called a **session**.

### 7.1 Overall Events Sequencing

In order to be able to proceed to logon, the RCST shall be in the **Receive Synchronization state**, which is reached following the **Initial synchronization procedure** described in subclause 7.2.



The entry of an RCST into the system is then achieved through the following four phases:

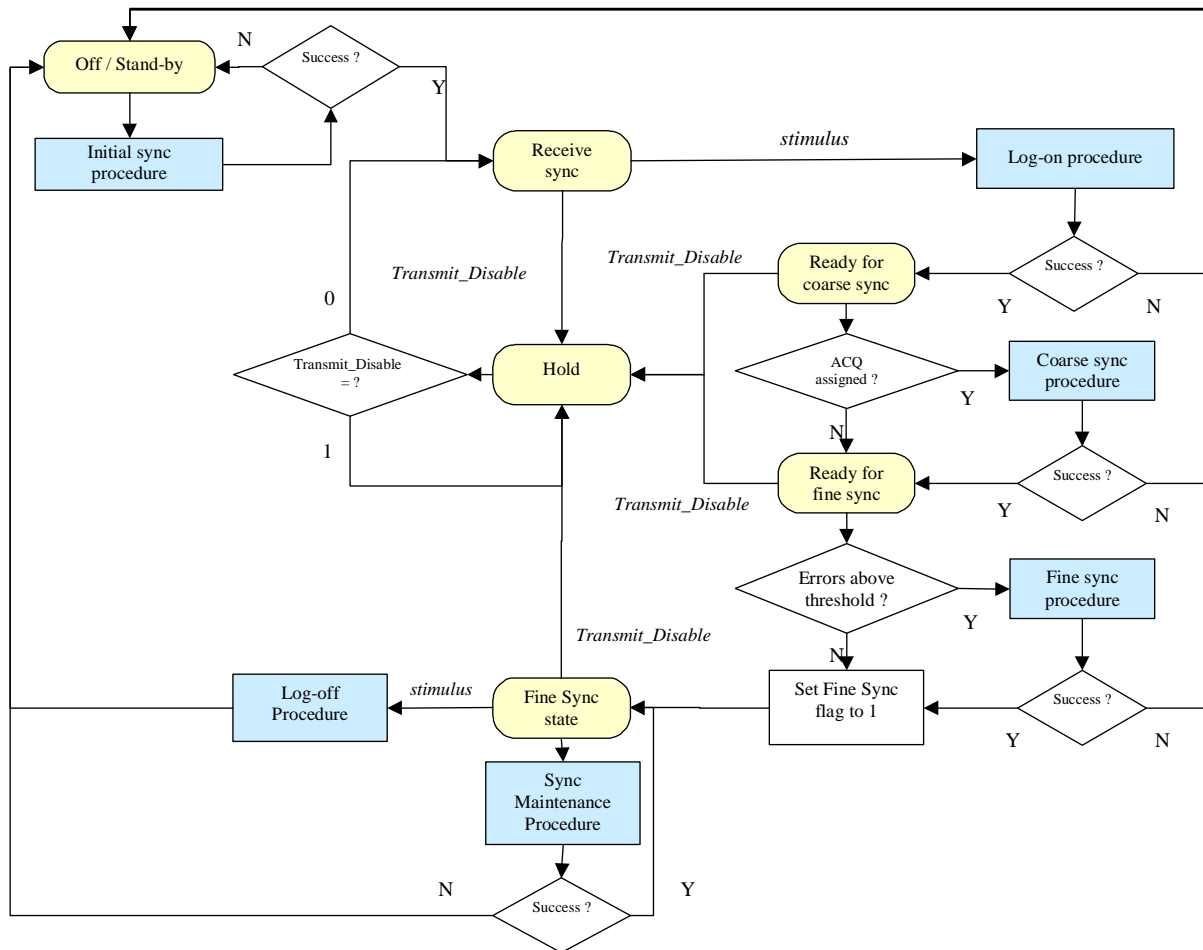
- **Logon procedure:** the RCST requests initial access to the network and gets initial logon information from the network (or alternatively the logon request may be rejected by the network). See subclause 7.3;
- **Acquisition Coarse synchronization procedure** (optional): the RCST improves its physical synchronization (frequency, time, and power adjustments). See subclause 7.4;
- **Fine synchronization procedure** (optional): the RCST completes its physical synchronization. See subclause 7.5;
- **Synchronization maintenance procedure:** the RCST maintains its physical synchronization during the entire session. See subclause 7.6.

Corresponding to the *procedures*, the RCST can be in one of the following *states*:

- **Hold:** the RCST is in hold mode. The RCST Status field in the TIM contains a flag called Transmit\_Disable (see subclause 8.5.5.8). A terminal that receives a TIM with this flag set to 1 shall cease transmission and release all assigned logon session parameters (i.e. logon\_id, group\_id, timeslot allocations) and enter the Hold state. This can happen when the terminal is in the Receive sync state, Ready for coarse sync state, Ready for fine sync state or the Fine Sync state. A terminal that is in the Hold state shall remain there after a power off or reset. A terminal goes from the Hold state to the Receive sync state only when it receives a TIM with the Transmit\_Disable flag set to 0;
- **Inactive Off / Stand-by** - the RCST is not powered or on a stand-by mode or has lost synchronization;
- **Receive sync** - the RCST has acquired the forward link;
- **Ready for coarse sync** – the RCST has been detected by the NCC, and may initiate a coarse synchronization procedure;
- **Ready for fine sync** – the RCST has been detected by the NCC, and may initiate a fine synchronization procedure;
- **Fine sync**– the RCST is synchronized and can send traffic.

The logoff procedure described in subclause 7.7 allows the RCST to leave the network.

The following figures give an overview of the sequence, where  represent states and  represent procedures. Figure 25 shows the RCST synchronization state diagram.



**Figure 25: RCST synchronization state diagram**

All the states, events, conditions and procedures are further described in this section.

The RCST is not allowed to transmit TRF bursts until it has reached the "Fine Sync" state.

### Signalling Messages

The exchange of signalling messages between RCST and NCC, including optional messages, is illustrated in figure 26.

This illustrates the normal flow of events in the case that the optional coarse and fine synchronization procedures are used.

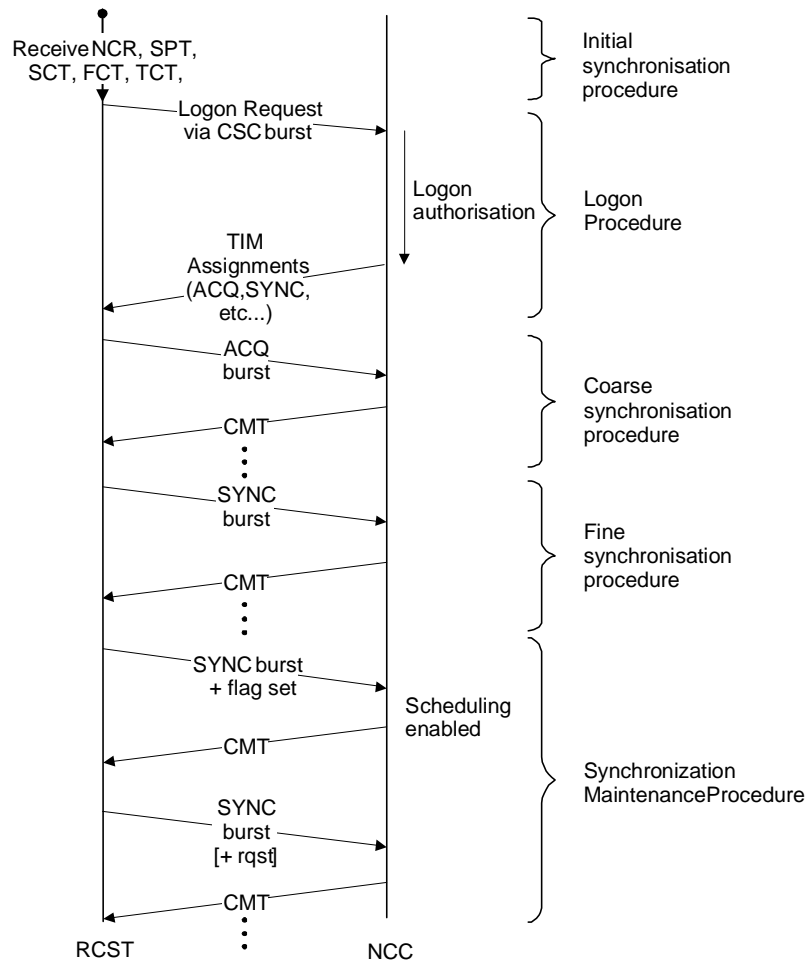


Figure 26: Example of RCST Network Entry Signalling Flow

## 7.2 Initial synchronization procedure

Following the power-up, the RCST shall proceed as detailed below:

- the RCST shall first follow the procedures described in subclause 8.5.5.11 to find all necessary control information related to the operation of the RCS network. This includes NCR synchronization, through which the RCST initiates its internal clock, by tracking the NCR which is transmitted by the NCC on the forward link;
- the RCST shall then calculate the satellite ranges for both forward and return links using the satellite ephemeris data contained within the Satellite Position Table (SPT) plus a knowledge of its own location (latitude, longitude and height above sea level). It shall use these ranges to calculate the corresponding satellite to RCST and RCST to satellite propagation delays. The nominal satellite position, which can be found in the NIT, shall be used in the case that the NCC does not transmit a SPT;
- the RCST shall continue to receive the NCR throughout the session. In the event that NCR synchronization is lost, the RCST ceases transmission (see subclause 7.7.3) and shall re-start the initial synchronization procedure. Similarly, any failure of the RCST during one of the later-described procedures takes the RCST back to the initial synchronization procedure;
- the RCST shall receive the burst time plan transmitted by the NCC at regular intervals. The BTP is contained in the Forward link Signalling, and is made of the Superframe, Frame and Timeslot Composition Tables. All these tables are described in clause 8. The RCST shall also acquire the broadcast TIM.

After following these steps, the RCST shall enter the Receive sync state.

## 7.3 Logon procedure

After the RCST has received all SI tables related to the structure of the satellite interactive network, it is ready to initiate a network logon, in order to be admitted into the system and be ready to handle traffic. The RCST can decide to move from the "Receive sync" state e.g. because it is booting up or because it wants to transmit data and is no longer logged-on after a long period of inactivity. Alternatively, the network may trigger the logon procedure by sending a "Wake up" signal to the RCST in an unicast TIM as described in subclause 8.5.5.8.

The logon procedure is illustrated in figure 27 and detailed below.

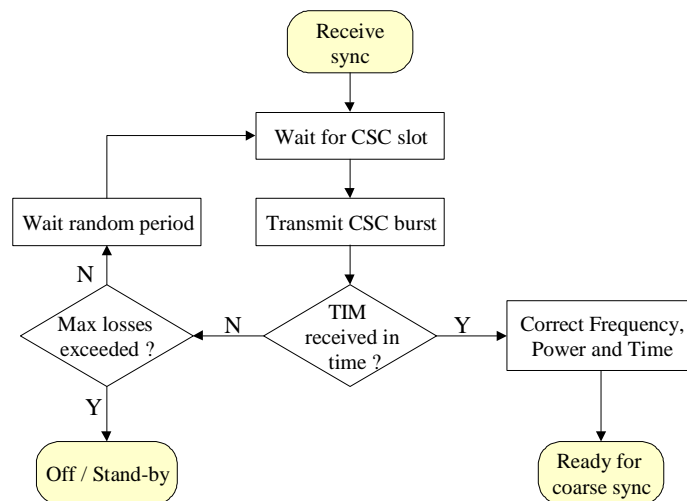
The RCST sends a logon request in a CSC timeslot using Slotted-Aloha random access. This request contains the RCST MAC address and a field indicating the capabilities of the terminal (see subclause 6.2.3). If it is received correctly, the NCC will proceed with the next step. In the absence of a reply from the NCC in due time, the RCST shall assume that there was a collision between multiple simultaneous requests and shall retry after a maximum, randomly-selected interval. If applicable, upper limits for repetition rate and duration of the logon request are specified in [6]. Parameters for this retransmission scheme are retrieved by the terminal from the Contention Control descriptor (see subclause 8.5.5.10.14), acquired during the Initial synchronization procedure.

The NCC verifies that transmission resources are available (ACQ and SYNC bursts) and checks if the administrative aspects are satisfied (e.g., account is valid, account is paid, etc.). If all conditions are met, the NCC proceeds with the next step.

The NCC sends a TIM message to the RCST as an acknowledgement. This "logon" TIM shall contain the information detailed in subclauses 8.5.5.8 and 8.5.5.10.1.

Subsequently, an ATM VPI/VCI shall be used by RCSTs sending ATM TRF bursts. For RCSTs sending optional MPEG packets in TRF bursts, a PID is assigned. VPI/VCI or PID values are assigned in the TIM Logon Initialize descriptor (subclause 8.5.5.10.4).

**NOTE:** Under exceptional conditions such as a severe network failure, access to the CSC bursts may be dedicated rather than random.



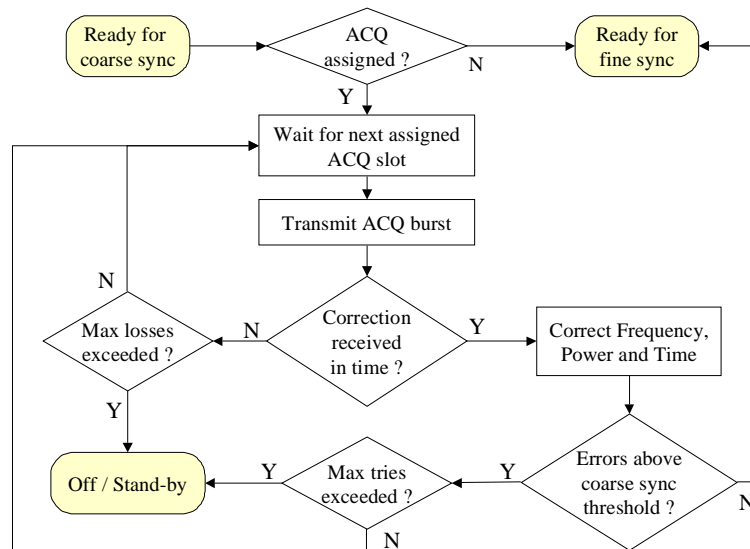
**Figure 27: Logon Procedure**

## 7.4 Coarse synchronization procedure (Optional)

In a network where all RCSTs are locked to the NCR on the forward link, the NCC can correct all frequency and timing errors other than differential Doppler between the RCST and the NCC. Initial burst time errors can be low when the satellite and terminal position are known. Provided the NCC/gateway receivers can handle these residual errors, which are small for a satellite maintained in a tight 'box', there is no requirement for the RCST to perform the ranging process of the coarse synchronization procedure. In this case the TIM will not contain an ACQ assign descriptor (subclause 8.5.5.10.5), and the RCST shall enter directly the "Ready for fine synchronization" state. Otherwise, if the descriptor is contained then the RCST shall perform the coarse synchronization procedure given here.

After the RCST has logged on and been given the authorization to proceed, it shall commence the acquisition phase to achieve timing and frequency synchronization and power adjustment. The procedure is illustrated in figure 28, and detailed below.

The RCST is assigned ACQ bursts via the TIM. The RCST sends an ACQ burst at its reserved time slot. The NCC measures the timing, frequency and power error of the ACQ burst, relative to the system reference, and sends this information back in the Correction Message Table (CMT, see subclause 8.5.5.9). The RCST in turn adjusts its transmission parameters, and retries. This process continues until the accuracy is within the "coarse sync thresholds" indicated to the RCSTs in the ACQ Assign Descriptor (see subclause 8.5.5.10.5). The same descriptor provides means to limit the number of loops in this procedure ("Max tries exceeded"). Parameters for this procedure are contained in the correction\_control\_descriptor (see subclause 8.5.5.10.15).

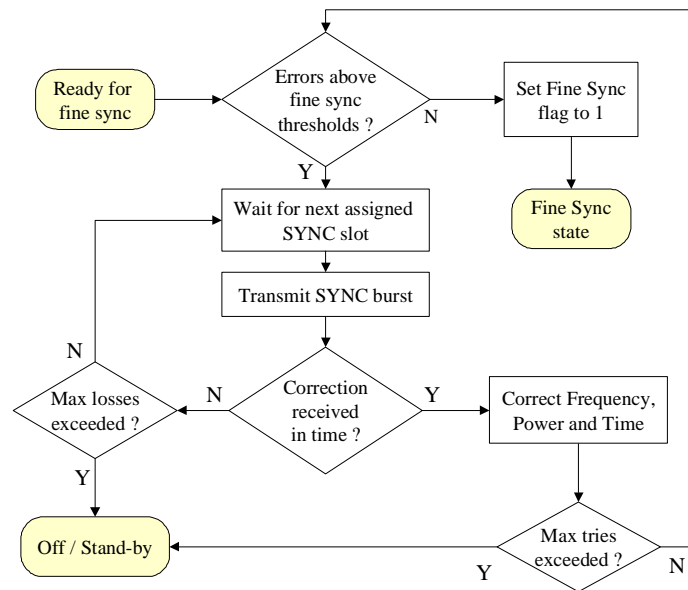


**Figure 28: Coarse Synchronization Procedure (optional)**

## 7.5 Fine Synchronization Procedure (Optional)

This procedure is quite similar in principle to the coarse synchronization procedure described in subclause 7.4, but it uses dedicated SYNC slots instead of ACQ bursts. This procedure is only performed if the errors indicated in the latest correction message (which was either included in the "logon" TIM received right after CSC, or in a CMT during the optional ACQ procedure) are larger than the "fine sync thresholds" indicated in the SYNC Assign descriptor of the "logon" TIM (subclause 8.5.5.10.6).

The fine synchronization procedure is shown in figure 29. Parameters for this procedure are contained in the correction\_control\_descriptor (see subclause 8.5.5.10.15).



**Figure 29: Fine synchronization procedure (optional)**

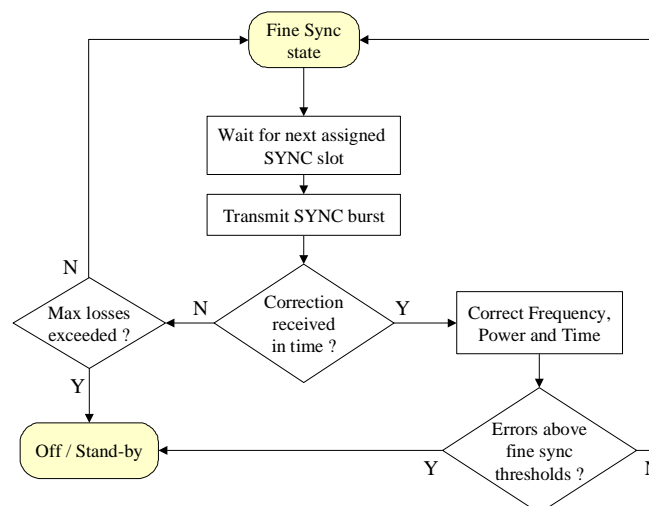
NOTE: It may be necessary to average a number of error measurements in order to achieve sufficient accuracy.

## 7.6 Synchronization Maintenance procedure

Upon achieving fine synchronization, the RCST is allowed to transmit TRF bursts. It shall in parallel proceed to maintain synchronization. This procedure is carried out continuously for the duration of the session as shown in figure 30. However, in case the NCC has indicated to the RCST that a security sign-on was required by the network (see the Logon Initialize descriptor subclause 8.5.5.10.4), the RCST shall wait for the security handshake to occur (see subclause 9.4). Once it has returned a "Security Sign-on Response" message, the RCST is allowed to transmit user traffic in its allocated TRF timeslots.

Loop counts, timeouts and thresholds are indicated to the RCST through the SYNC\_Assign descriptor embedded in a TIM (see subclause 8.5.5.10.6).

Parameters for this procedure are contained in the correction\_control\_descriptor (see subclause 8.5.5.10.15).



**Figure 30: Synchronization Maintenance Procedure**



## 7.7 Logoff Procedure

### 7.7.1 General

The RCST logoff procedure described in this subclause applies when the RCST is in the Fine Sync state.

When the RCST logs-off, it shall cease all transmission and its logical address and SYNC timeslot shall be removed from the active list and made available to other RCSTs that may join the network. The logoff procedure is initiated as a result of a session termination (normal) or of a failure (abnormal) as described in the following subclauses.

### 7.7.2 Normal

A normal logoff can be initiated either automatically or manually by the user at the end of a session. For an RCST initiated logoff, a logoff request message shall be transmitted in the M&C section. For an NCC initiated logoff, the logoff message shall be carried by a Terminal Information Message (TIM) addressed to the RCST (see subclause 8.5.5.8).

### 7.7.3 Abnormal

An RCST shall logoff in the following conditions:

- loss of receive synchronization, i.e.:
  - NCR not received for several consecutive seconds;
  - CMT burst correction not received for several consecutive SYNCs.

---

## 8 Control and management

This clause defines the messages to allow an RCST to logon to the satellite interactive network. These will be used to co-ordinate an identification of the calling RCST, a process to adjust the power of the RCST, and a logon procedure which gives an identification to the RCST that can be used to transmit meta-signalling to request traffic connections. The following subclauses detail the protocol stacks used in the forward and return link and each one of the messages exchanged between the NCC/Gateway and RCST and vice-versa.

As a minimum set of requirements the RCST shall comply with the Control & Monitoring Functions (CMF) specified in [6] if applicable. Among others, the present document requires that the RCST is only allowed to transmit, when it receives its control correctly.

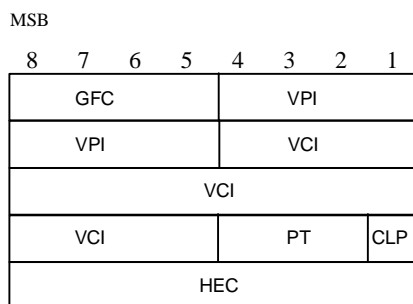
### 8.1 Protocol stack

The protocol stack is based on the DVB/MPEG2-TS standard in the forward link and is based on ATM cells or optional MPEG2-TS Packets mapped onto TDMA bursts in the return link. The forward link shall use DVB multiprotocol encapsulation according to TR 101 202 [2] to send IP datagrams over MPEG2 Transport Streams.

#### 8.1.1 RCST Type A (IP)

The RCST Type A shall be able to support IP services only.

ATM cells are used on the return link to benefit from the AAL5 segmentation and re-assembly (SAR) function. The mechanism to carry IP over ATM AAL5 shall be as specified in RFC 1483 [7] (section 4.1 Payload Format for Routed IP PDUs). From the two multiplexing methods defined there the one called "VC based multiplexing of Routed Protocols", shall be applied. The ATM cells shall be mapped into MF-TDMA traffic bursts as defined in subclause 6.2.1.1. The 53 bytes of the ATM cell are made of 5 bytes of header and 48 bytes of payload. ATM cells are used either for user traffic or for control and/or management traffic (handled by upper layers). For type A RCSTs, the 5-byte header of an ATM cell shall follow the ATM UNI cell format, as represented on figure 31.



**Figure 31: ATM cell header format**

The use of the different fields shall be as follow:

- GFC: reserved. This field has another meaning when the optional security mechanism is used (see subclause 9.4.6.3);
- VPI: set to the VPI value transmitted by the network to the terminal after logon (see subclause 8.5.5.10.4), or a VPI value obtained through upper layer signalling;
- VCI: set to the VCI value transmitted by the network to the terminal after logon (see subclause 8.5.5.10.4), or a VCI value obtained through upper layer signalling.

A given VPI/VCI pair shall not be allocated by the network to more than one RCST at a given time.

- PT: This field shall be used as follows:
  - bit 2 of octet 4 shall be used for AAL5 Segmentation and Reassembly process (AUU bit), as defined in ITU-T Recommendation I.363 ;
  - bit 3 of octet 4 is normally used by ATM signalling to indicate if the cell has experienced congestion. This bit shall be set to 0;
  - bit 4 of octet 4 (normally used by ATM signalling to indicate a traffic cell). This bit shall be set to 0 for traffic cells and to 1 for CTRL/MNGM bursts (see subclause 6.6.2).
- CLP: Cell Loss Priority. This bit is used to indicate the cells which would be first deleted in case of buffer congestion in one of ATM network nodes. This bit shall be set to 0;
- HEC: this field shall be generated as described in ITU-T Recommendation I.432 [13].

#### **Optional MPEG mode:**

Alternatively the IP SAR function can be provided by the DVB multiprotocol encapsulation method according to TR 101 202 [2] using MPEG2 Transport Stream packets as the container. The packets shall be mapped into MF-TDMA traffic bursts as defined in subclause 6.2.1.2.

The PID(s) for user traffic and, if applicable, CTRL / MNGM messages (see subclause 6.6.2) to be used in the return link MPEG TS packets are set by the logon initialize descriptor (see subclause 8.5.5.10.4) or obtained through upper layer signalling.

User traffic is sent between the User Device and a Gateway, whereas signalling is sent only between the NIU and the NCC. Figure 32 shows an example of a protocol stack for user traffic. The RCST and User Device are connected by 10 baseT. IP is used for user traffic on the network layer. Figure 33 shows the protocol stack for signalling.

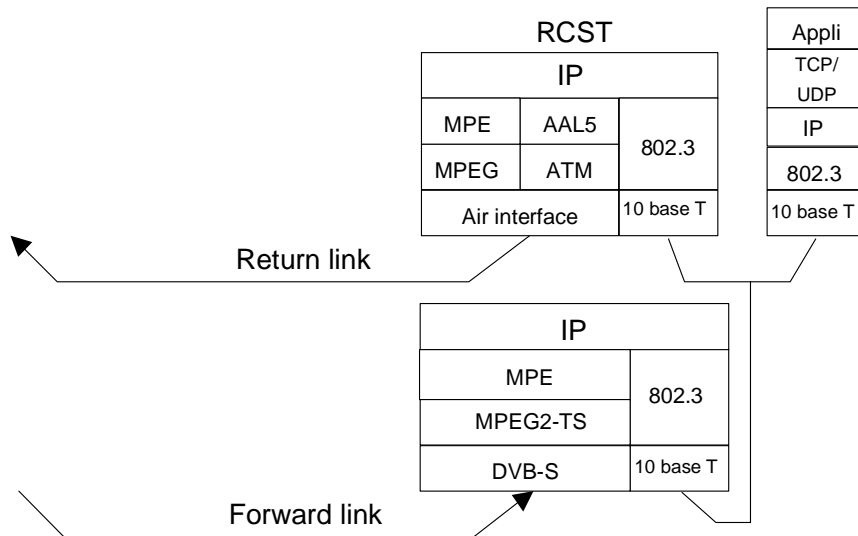


Figure 32: Example of protocol stack for user traffic with Type A RCST

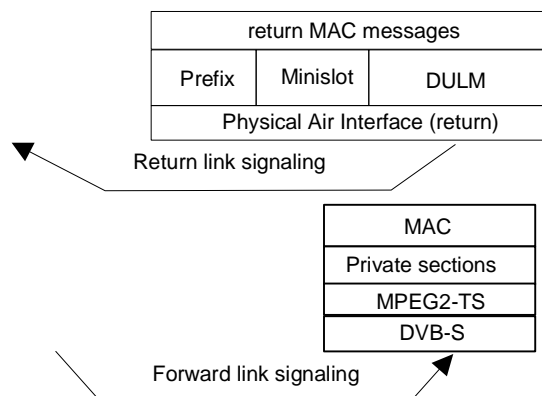


Figure 33: Protocol stack for signalling

### 8.1.2 Optional RCST Type B (Native ATM)

The RCST Type B shall be able to operate as RCST Type A and shall also be able to support native ATM protocols by encapsulating ATM cells within an MPEG2 Transport Stream on the forward link, as defined in TR 100 815 [8]. On the return link ATM cells shall be mapped into MF-TDMA traffic bursts as defined in subclause 6.2.1.1.

The RCST can support Permanent Virtual Circuits (PVC's) and Switched Virtual Circuits (SVC's) on the forward- and return link as an UNI. Standard compliant signalling according to ITU-T Recommendation Q.2931 [11] shall be used. In difference to a normal ATM environment there is a shared medium between the terminal and the gateway. During the Logon procedure a VPI/VCI is assigned to the RCST by the Logon initialize descriptor that is carried in the TIM, see subclauses 8.5.5.10.4 and 8.5.5.8. The assigned VPI/VCI shall overwrite the values 0/5 that are normally used for ITU-T Recommendation Q.2931 [11] signalling.

ATM connectivity in RCSTs is optional.

## 8.2 RCST Addressing

On the Forward Link, RCSTs shall be uniquely identified by a physical MAC address and a logical address.

The *MAC address* is a physical address stored in non-volatile memory. It corresponds to a unique RCST physical identifier. It shall follow the IEEE 802.3 [12] standard and shall consist of 48 bits. The value 0xFFFFFFFF shall be reserved for broadcasting to all RCSTs. The MAC address shall be used inside a CSC burst and in DSM-CC private data sections that carry TIMs. It will also be used to encapsulate IP datagrams into MPEG2-TS frames TR 101 202 [2].

The logical address is composed of two fields, the *Group\_ID* and *Logon\_ID* which are assigned to each RCST during logon. They are used for addressing individual RCSTs until logoff.

The *Group\_ID* corresponds to a group of logged-on RCSTs. It shall consist of 8 bits. The value 0xFF shall be reserved for system use.

The *Logon\_ID* uniquely identifies the RCST within a group identified with the *Group\_ID*. The *Logon\_ID* shall consist of 16 bits. The value 0xFFFF shall be reserved for system use (contention mode on the return link).

For a Type A RCST any data (user traffic) that is destined to a specific RCST shall be transmitted using the RCST MAC address. Any data (user traffic) that is destined to all Type A RCSTs shall be transmitted using the broadcast MAC address (0xFFFFFFFF).

Independently, each protocol used at higher layers may impose its own addressing scheme, e.g. IP addresses, etc.

## 8.3 Forward Link Signalling

DVB defines a set of tables built upon the MPEG PSI tables to provide detailed information regarding the broadcast network. Such DVB tables are referred as the Service Information (SI) tables. In a two-way Satellite Interactive Network, consisting of a forward and return link via satellite, medium access control information and other signalling are communicated through the forward link and shall be transmitted in a DVB compliant manner. Thus, the specifications for Service Information (SI) in DVB systems EN 300 468 [4] shall apply. The forward link signalling consists of general SI tables, carrying information about the structure of the satellite interactive network, and RCST specific messages sent to individual RCSTs, private data fields defined for standard DVB-SI tables, special Transport Stream packets (PCR Insertion) and descriptors, including private descriptors for standard DVB-SI tables.

### 8.3.1 General SI Tables

General SI data describing the Satellite Interactive Network organization are structured as six types of table. These tables are broadcast in private sections (c.f. ISO/IEC 13818-1 [9]). Where applicable the use of descriptors allows a flexible approach to the organization of the tables and allows for future compatible extensions. The precise definition of the table content is given in subclause 8.5.5.

#### 8.3.1.1 Superframe Composition Table (SCT)

This table describes the sub-division of the entire satellite interactive network into superframes and frames. The table contains for each superframe, a superframe identification, a centre frequency, an absolute start time expressed as an NCR value and a superframe count. Each superframe is further divided into frames. Each type of frame is identified by a *frame\_id*. The frame position within a superframe is given by a frame number used for timeslot assignments in the TBTP. The frame numbering convention is described in subclause 6.7.2. The frames are positioned relative to the centre frequency and start time of the associated superframe.

#### 8.3.1.2 Frame Composition Table (FCT)

This table describes the partitioning of the frames into time-slots. The table contains for each *frame\_id* (i.e. for each frame type) a frame duration, the total number of timeslots contained in the frame, the start times and frequency offsets for the timeslots. The transmission parameters (such as symbol rate, code rate, preamble, etc.) for each timeslot are referred by a time-slot identifier (*timeslot\_id*) to a description conveyed in the TCT.

### 8.3.1.3 Time-Slot Composition Table (TCT)

This table defines the transmission parameters for each time-slot type identified by the time-slot identifier. It provides information about the timeslot properties such as symbol rate, code rate, preamble, payload content (TRF with ATM cells, TRF with MPEG2 TS packets, CSC, ACQ, SYNC) and others.

### 8.3.1.4 Satellite Position Table (SPT)

This table contains the satellite ephemeris data required to update the burst position at regular intervals. The table shall contain this data for those satellites that form an active part of a particular Satellite Interactive Network.

### 8.3.1.5 Correction Message Table (CMT)

The NCC sends the Correction Message Table to groups of RCSTs. The purpose of the CMT is to advise the logged-on RCSTs what corrections shall be made to their transmitted bursts. The CMT provides correction values for burst frequency, timing and amplitude. The CMT contains the corrections for the RCSTs with the most recently measured ACQ and SYNC bursts.

### 8.3.1.6 Terminal Burst Time Plan (TBTP)

This message is sent by the NCC to a group of terminals. The group is addressed by a logical Group\_ID, while each individual terminal is addressed by a logical Logon\_ID. Both Group\_ID and Logon\_ID are notified to the terminal at logon time. It contains one or more entries for each RCST, with each entry defining an assignment of a contiguous block of timeslots. Each traffic assignment is described by the number of the start timeslot in the block and a repetition factor giving the number of consecutive timeslot allocations.

The TBTP allows timeslots to be allocated once or continuously. In the latter case a mechanism is provided to add or remove time slot allocations from the terminal burst time plan.

## 8.3.2 Terminal Information Message (TIM)

This message is sent by the NCC either to an individual RCST addressed by its MAC address (uni-cast message) or broadcast to all RCSTs using a reserved broadcast MAC address and contains static or quasi static information about the forward link such as configuration.

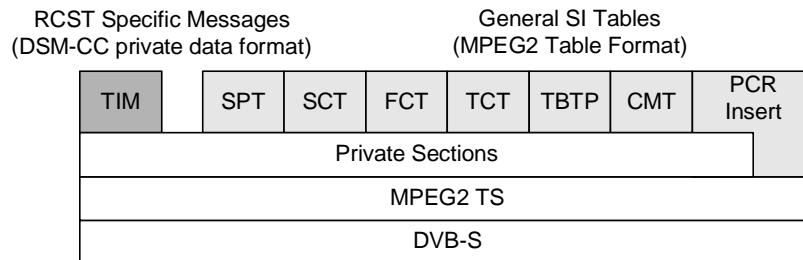
This message may also be used to facilitate the handing over of an RCST to a different group or network group or network or to switch a group of RCSTs to a different forward link signalling service on another MPEG2-TS for example. This message is sent in a DSM-CC private data section (c.f. ISO/IEC 13818-6 [17]).

## 8.3.3 PCR Insertion TS Packet

The PCR Insertion TS Packet shall be used for inserting the NCR value used for return link synchronization. The PID shall be assigned via the PMT of the Forward Link Signalling service (see subclause 8.5.5.11). The optional payload may contain information about the delay between all relevant satellite to NCC and satellite to Gateway combinations.

## 8.3.4 Summary

Figure 34 gives an overview of the protocol stack for forward link signalling. This protocol stack corresponds to the protocol shown in figure 33.



**Figure 34: Protocol Stack for Forward Signalling**

### 8.3.5 Repetition Rates

All sections of the SCT, FCT, TCT, SPT and broadcast TIM shall be transmitted at least every 10 seconds to allow newly activated RCSTs to acquire the necessary start-up state. In addition, the TIM shall be updated as required to reflect system status changes requiring immediate notification of the RCSTs.

The TBTP shall be updated every superframe.

The update rate of the NCR value in the PCR Insertion TS Packet shall be between 200 times per second and 10 times per second. If used, the optional PCR Insertion TS packet payload section shall be transmitted at least once a second.

The uni-cast TIM shall be updated as needed to reflect changes affecting a given RCST.

## 8.4 Return Link Signalling

### 8.4.1 RCST synchronization and Identification messages

The NCC manages RCST synchronization. The synchronization process is a set of signalling messages exchanged between the NCC and the RCST on reserved timeslots, which allow fine tuning of all synchronization parameters, timing, frequency, and power (see clause 7).

The messages described below are used for the synchronization of the RCST. The ACQ and SYNC bursts are special bursts containing symbols dedicated for the NCC to be able to measure frequency and timing offset.

The NCC requests an RCST to transmit a sequence of ACQ bursts via the TIM by assigning a BTP ACQ time slot for a limited number of repeats. In systems which can tolerate the required signalling bandwidth, the NCC requests an RCST to transmit an ACQ bursts as and when required by an ACQ burst assignment signalled via the TBTP.

In systems using periodic SYNC transmissions, the NCC also requests an RCST to periodically transmit a SYNC burst via the TIM by assigning a SYNC time slot once every N superframes. In systems using ad hoc (non periodic) SYNC transmissions, the NCC requests an RCST to transmit a SYNC burst as and when required by a SYNC timeslot assignment signalled via the TBTP.

The NCC determines power, frequency and burst time error of the RCST. Corrections for frequency and burst time are sent in a CMT for correction by the RCST.

**Forward :** Message/DSM-CC and SI Section/MPEG2-TS/DVB-S

**Return:** Special bursts/Air Interface

**Messages used:** TIM (forward) – [DSM-CC] or TBTP [SI]

CMT (forward) – [SI]

CSC (return)

ACQ (return)

SYNC (return)

## 8.4.2 Configuration parameters between RCST and NCC (optional)

The NCC uses the configuration parameters exchanged between RCST and NCC to identify the functional capability of the RCST and therefore what transmission characteristics it can demand from that particular RCST. In addition the NCC can obtain information such as manufacturer identification, RCST version (number of forward link receivers, RCST type (amplifier power), user identification, hardware version, software version, RCST position (latitude, longitude, altitude), Outdoor Unit characteristics (power, antenna size and antenna gain), type of RCST connection (SMATV, single user, multiple user), installer identification, postal code of the area, date & time of installation and more). A private Management Information Base (MIB) in the RCST stores the configuration parameter values in variables. The NCC uses SNMP commands to obtain the current configuration parameter values from the RCST MIB. The NCC sets a flag in the RCST MIB to acknowledge receipt of the configuration parameter values. An SNMP agent in the RCST responds to commands from an SNMP client in the NCC. This exchange of configuration parameters is optional, the RCST indicates in the CSC burst (see subclause 6.2.3) if it implements SNMP.

**Forward:** SNMP/UDP/IP/DSM-CC/MPEG2-TS/DVB-S

**Return:** SNMP/UDP/IP/Traffic bursts/Air Interface

**Messages:** **get-request [MIB variable] (forward)**  
**get-next-request [MIB variable] (forward)**  
**get-response [MIB variable, value] (return)**  
**set-request [acknowledgement flag] (forward)**

NOTE: Private MIB variables are out of the scope of the present document.

## 8.4.3 Other Messages for Network Management (optional)

The NCC and RCST send other SNMP messages whenever they are needed for network management. Such messages implement installation procedures, software upgrades, transmit authorization or prohibition, individual/group control and traffic forward link assignment, RCST status enquiries, and RCST or NCC requests to leave the network. The NCC queries MIB values to determine RCST status and stores values in the MIB variables to trigger RCST actions. The RCST sends *trap* messages to notify the NCC when it has accomplished triggered actions or to issue requests. For instance, the NCC sends a *set-request* message with a MIB variable value to authorize or prohibit transmission. Similarly, the RCST sends a *trap* message to the NCC to request to leave the network. This exchange of messages is optional, the RCST indicates in the CSC burst (see subclause 6.2.3) if it implements SNMP.

**Forward:** SNMP/UDP/IP/DSM-CC/MPEG2-TS/DVB-S

**Return:** SNMP/UDP/IP/Traffic bursts/Air Interface

**Messages:** **get-request [MIB variable] (forward)**  
**get-next-request [MIB variable] (forward)**  
**get-response [MIB variable] (return)**  
**set-request [MIB variable, value] (forward)**  
**trap [MIB variable value, value] (return)**

NOTE: Private MIB variables are out of the scope of the present document.

## 8.4.4 Burst Time Plan Exchange

The Burst Time Plan (BTP) is sent to all the RCSTs affected by using the Terminal Burst Time Plan message (TBTP). This information is the basis for the RCSTs to know when to transmit their bursts. Capacity Requests (CR) messages are sent by the RCSTs to the NCC for all volume based connections depending on the data present in the queue. For constant bit rate connections, the TDMA scheduler at the NCC will automatically update the BTP according to the set-up parameters.

<b>Forward:</b>	<b>Message/SI Table/MPEG2-TS/DVB-S</b>
<b>Return:</b>	<b>Capacity requests (CR)/Air Interface</b>
<b>Messages:</b>	<b>TBTP (forward)</b>
	<b>CR (return)</b>

## 8.5 Coding of SI for Forward Link Signalling

### 8.5.1 Introduction

Forward Link Signalling for Satellite Interactive Network is divided into five parts:

- general SI Tables (SCT, FCT, TCT, SPT, TBTP and CMT);
- RCST Specific Messages (TIM);
- Private Data fields defined for standard DVB-SI tables;
- Special Transport Stream packets (PCR Insertion);
- descriptors, including private descriptors for standard DVB-SI tables.

The following subclauses describe the coding and definition of these tables, messages and descriptors.

### 8.5.2 SI Table Mechanism

SI Tables for Satellite Interactive Network are transmitted in private sections as defined in ISO/IEC 13818-1 [9], Systems.

The syntax and semantics of SI tables defined in EN 300 468 [4] for the segmentation of tables in one or more sections, the section length, the section identification, the mapping of sections into Transport Streams, the repetition rates and the random access shall also be applicable to SI tables for Satellite Interactive Network defined in the present document.

If required, the tables defined in the present document may be scrambled to prohibit traffic analysis. If the tables are scrambled, the same mechanisms as defined in EN 300 468 [4] for scrambled SI tables shall apply.

### 8.5.3 DSM-CC Section Mechanism

RCST Specific Messages for Satellite Interactive Network are transmitted in DSM-CC private data sections as defined in ISO/IEC 13818-6 [17] in general and in EN 300 468 [4] for DVB.

If required, RCST Specific Messages defined in the present document may be scrambled to prohibit traffic analysis. They may be scrambled on an individual basis.



## 8.5.4 Coding of PID and table\_id fields

**Table 10: PID and table\_id allocation SI Tables**

Table and private data sections defined in the present document	PID	Table_id
RMT	Assigned*	0x41
SCT	Assigned	0xA0
FCT	Assigned	0xA1
TCT	Assigned	0xA2
SPT	Assigned	0xA3
CMT	Assigned	0xA4
TBTP	Assigned	0xA5
PCR packet payload	Assigned	0xA6
Reserved for future use		0xA7 to 0xAF
TIM	Assigned	0xB0
Reserved for future use		0xB1 to 0xBF
NOTE: *This PID shall be defined to be a given value across all interactive networks in a given system.		

Table 10 lists the PID and table\_id values which shall be used for the TS packets which carry SI tables and RCST Specific Messages defined in the present document.

## 8.5.5 Table definitions

The following sub-clauses describe the syntax and semantics of the different types of table.

NOTE: The symbols and abbreviations, and the method of describing syntax used in the present document are the same as those defined in subclauses 2.2 and 2.3 of ISO/IEC 13818-1 [9].

The mnemonics defined in subclause 2.2.6 of ISO/IEC 13818-1 [9] are used in the tables of the present document. For convenience, the present document defines a further three mnemonics to cover frequently used parameter formats, as follows:

- spfmsbf = single precision floating point value, which is a 32 bit value formatted in accordance with ANSI/IEEE Standard 754 (1985) [16]. The most significant bit (i.e. the most significant bit of the exponent) is first;
- upcrmsf = unsigned PCR count value. The coding of this type of parameter shall be identical to the coding of the program\_clock\_reference (PCR) described in ISO/IEC 13818-1 [9], comprising a base field of up to 33 bits and a 9 bit extension field. Where the number of bits is less than the full 42 bit PCR format, then the least significant 9 bits shall correspond to the extension field and the remaining bits shall correspond to the least significant bits of the base field. The most significant bit is first;
- flagsmf = multi-bit boolean flags field (e.g. status byte / word), with the most significant bit first. Each flag is asserted on a logic '1'. In an extension to the ISO/IEC 13818-1 [9] syntax, an individual flag is referenced using standard object reference "dot" notation, i.e. the "foo" flag in parameter "status" is referenced as "status.foo" and takes the value TRUE if the foo flag bit is '1'.

### 8.5.5.1 Standard Section Headers

The following standard headers have been tailored for forward link signalling use, and represent a specific subset of the more general formats specified in the DVB and ISO standards.

## 8.5.5.1.1 SI Section Header

Table 11: Standard SI Section Header

Syntax	No. of bits	Mnemonic
table_id	8	uimsbf
section_syntax_indicator	1	bslbf
reserved_for_future_use	1	bslbf
reserved	2	bslbf
section_length	12	uimsbf
interactive_network_id	16	uimsbf
reserved	2	bslbf
version_number	5	uimsbf
current_next_indicator	1	bslbf
section_number	8	uimsbf
last_section_number	8	uimsbf

The standard header for a SI section occupies a total of 64 bits, and shall be as defined in table 11.

Semantics for the standard SI section header:

- table\_id: This 8 bit field identifies the table. See table 10 for the table\_id values;
- section\_syntax\_indicator: The section\_syntax\_indicator is a 1-bit field which shall be set to "1";
- section\_length: This is a 12-bit field, the first two bits of which shall be "00". It specifies the number of bytes of the section, starting immediately following the section\_length field and including the CRC. The section\_length shall not exceed 1 021 so that the entire section has a maximum length of 1 024 bytes;
- interactive\_network\_id: This is a 16-bit field, which serves as a label to identify the Satellite Interactive Network, to which the table shall apply;
- version\_number: This 5-bit field is the version number of the sub\_table. The version\_number shall be incremented by 1 when a change in the information carried within the sub\_table occurs. When it reaches value 31, it wraps around to 0. When the current\_next\_indicator is set to "1", then the version\_number shall be that of the currently applicable sub\_table defined by the table\_id and interactive\_network\_id\_mask. When the current\_next\_indicator is set to "0", then the version\_number shall be that of the next applicable sub\_table defined by the table\_id and interactive\_network\_id;
- current\_next\_indicator: This 1-bit indicator, when set to "1" indicates that the sub\_table is the currently applicable sub\_table. When the bit is set to "0", it indicates that the sub\_table sent is not yet applicable and shall be the next sub\_table to be valid;
- section\_number: This 8-bit field gives the number of the section. The section\_number of the first section in the sub\_table shall be "0x00". The section\_number shall be incremented by 1 with each additional section with the same table\_id and interactive\_network\_id;
- last\_section\_number: This 8-bit field specifies the number of the last section (that is, the section with the highest section\_number) of the sub\_table of which this section is part.

## 8.5.5.1.2 DSM-CC Private Section Header

**Table 12: Standard DSM-CC Private Section Header**

Syntax	No. of bits	Mnemonic
table_id	8	uimsbf
section_syntax_indicator	1	bslbf
private_indicator	1	bslbf
reserved	2	bslbf
section_length	12	uimsbf
MAC_address_6	8	uimsbf
MAC_address_5	8	uimsbf
reserved	2	bslbf
payload_scrambling_control	2	bslbf
address_scrambling_control	2	bslbf
LLC_SNAP_flag	1	bslbf
current_next_indicator	1	bslbf
section_number	8	uimsbf
last_section_number	8	uimsbf
MAC_address_4	8	uimsbf
MAC_address_3	8	uimsbf
MAC_address_2	8	uimsbf
MAC_address_1	8	uimsbf

The standard header for a DSM-CC private section occupies a total of 96 bits, and shall be as defined in table 12.

Semantics for the standard DSM-CC private section header:

- table\_id: This 8 bit field identifies the table. See table 10 for the table\_id values;
- section\_syntax\_indicator: The section\_syntax\_indicator is a 1-bit field which shall be set to "1" to denote that a CRC32 check field is used at the end of the section;
- private\_indicator: The private\_indicator is a 1 bit field that shall be set to the complement of the section\_syntax\_indicator (i.e. to '0');
- section\_length: This is a 12-bit field, the first two bits of which shall be "00". It specifies the number of bytes of the section, starting immediately following the section\_length field and including the CRC. The section\_length shall not exceed 1 021 so that the entire section has a maximum length of 1 024 bytes;
- MAC\_address\_[1]: this 48 bit field contains the MAC address of the destination. The MAC address is fragmented in 6 fields of 8 bits, labelled MAC\_address\_1 to MAC\_address\_6. The MAC\_address\_1 field contains the most significant byte of the MAC address, while MAC\_address\_6 contains the least significant byte;

NOTE: The order of the bits in the byte is not reversed, and the MSB of each byte is still transmitted first.

- payload\_scrambling\_control: this 2 bit field defines the scrambling mode of the payload section (see table 13). This includes the payload starting after the MAC\_address\_1 but excludes the CRC32 field. The scrambling method applied is user private. If the optional security mechanism described in subclause 9.4 is active, this field shall be as defined in subclause 9.4.6.3;

**Table 13: Coding of the payload\_scrambling\_control field**

value	payload scrambling control
00	unscrambled
01	defined by service
10	defined by service
11	defined by service

- address\_scrambling\_control: this 2 bit field defines the scrambling mode of the MAC address section (see table 14). This field enables a dynamic change of MAC addresses. The scrambling method applied is user private;

**Table 14: Coding of the address\_scrambling\_control field**

value	address scrambling control
00	unscrambled
01	defined by service
10	defined by service
11	defined by service

- LLC\_SNAP\_flag: This 1 bit flag shall be set to '0' to indicate that the payload does not use LLC/SNAP encapsulation;
- current\_next\_indicator: This 1-bit field shall be set to "1".
- section\_number: This 8-bit field gives the number of the section. The section\_number of the first section in the message shall be "0x00". The section\_number shall be incremented by 1 with each additional section for the same message;
- last\_section\_number: This 8-bit field specifies the number of the last section (that is, the section with the highest section\_number) of the message of which this section is part.

### 8.5.5.2 Superframe Composition Table (SCT)

The SCT (see table 15) conveys information relating to the organization of the satellite interactive network, in particular the sub-division of the framing structure. The combination of the interactive\_network\_id and the superframe\_id allows each superframe to be uniquely identified. To each satellite interactive network is assigned an individual interactive\_network\_id which serves as a unique identification code. The interactive\_network\_id shall be unique throughout a given network.

The SCT shall be segmented into superframe composition sections using the syntax described in EN 300 468 [4]. Any sections forming part of an SCT shall be transmitted in TS packets with a PID value assigned in the PMT. Any sections of the SCT shall have the table\_id value as defined in table 10.

**Table 15: Syntax of Superframe Composition Section**

Syntax	No. of bits		Information Mnemonic
	Reserved	Information	
Superframe_composition_section(){			
SI_private_section_header		64	-
superframe_loop_count		8	uimsbf
for(i=0;i<=superframe_loop_count;i++){			
superframe_id		8	uimsbf
uplink_polarization	6	2	bslbf
superframe_start_time_base		33	uimsbf
superframe_start_time_ext	6	9	uimsbf
superframe_duration		32	upcrmsf
superframe_centre_frequency		32	uimsbf
superframe_counter		16	uimsbf
frame_loop_count	3	5	uimsbf
for(j=0;j<=frame_loop_count;j++){			
frame_id		8	uimsbf
frame_start_time		32	upcrmsf
frame_centre_frequency_offset		24	tcimsbf
}			
}			
CRC_32		32	rpchof
}			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. Unless otherwise stated, these bits shall be set to '0'.		

Semantics for the `superframe_composition_section`:

- `SI_private_section_header`: This is the standard SI private section header defined in table 11, and occupies a total of 64 bits;
- `superframe_loop_count`: This is a 8 bit field indicating one less than the number of iterations in the loop that follows. A zero count indicates one loop;
- `superframe_id`: This is an 8-bit field which serves as a label for identification of this superframe from any other superframe within the satellite interactive network;
- `uplink_polarization`: This is a 2-bit field specifying the polarization of the transmitted signal (see table 16);

**Table 16: Polarization definition**

Polarization	Value
linear – horizontal	00
linear – vertical	01
circular – left	10
circular - right	11

- `superframe_start_time_base` and `superframe_start_time_ext`: These two fields give the absolute time of the beginning of the superframe identified by the `superframe_id`. The coding of the fields shall be identical to the coding of the `program_clock_reference` (PCR) described in ISO/IEC 13818-1 [9], with the two fields corresponding to the base and extension parts of the PCR respectively.

NOTE: The separation of base and extension by 6 reserved bits matches the format widely used in other DVB tables for PCR values.

- `superframe_duration`: This 32-bit field gives the duration of the superframe identified by the `superframe_id`, in terms of PCR count intervals. The 32 bits correspond to a maximum duration of 93,2 seconds;
- `superframe_centre_frequency`: This 32-bit field gives the absolute centre frequency of the superframe identified by the `superframe_id`. The frequency is given in multiples of 100 Hz;
- `superframe_counter`: This 16 bit field gives the superframe count value, modulo 65536. This is used to avoid ambiguity in the processing of the TBTP message;
- `frame_loop_count`: This 5 bit field indicates one less than the number of iterations in the loop that follows. A zero count indicates one loop. The frame numbers follow the numbering convention defined in subclause 6.7.2;
- `frame_id` gives the frame type identifier for the  $j^{\text{th}}$  frame, corresponding to a frame type defined in the FCT;
- `frame_start_time`: This 32 bit field gives the start time of the  $j^{\text{th}}$  frame relative to the superframe start time, in terms of PCR count intervals. The 32 bits correspond to a maximum duration of 93,2 seconds;
- `frame_centre_frequency_offset`: This 24-bit field gives the signed offset of the centre frequency of the  $j^{\text{th}}$  frame relative to the `superframe_centre_frequency` parameter (SCT). The frequency is given in multiples of 100 Hz;
- `CRC_32`: This is a 32-bit field that contains the CRC value that gives a zero output of the registers in the decoder defined in annex B of EN 300 468 [4] after processing the entire section.

### 8.5.5.3 Frame Composition Table (FCT)

The FCT conveys information describing the different frame types. This table defines the structure in the frequency / time space for each frame type.

The FCT shall be as defined in table 17. It shall be segmented into frame composition sections using the syntax described in EN 300 468 [4]. Any sections forming part of an FCT shall be transmitted in TS packets with a PID value assigned in the PMT. Any sections of the FCT shall have the table\_id value as defined in table 10.

**Table 17: Syntax of frame composition section**

Syntax	No. of bits		Information Mnemonic
	Reserved	Information	
Frame_composition_section(){			
SI_private_section_header		64	-
frame_ID_loop_count		8	uimsbf
for(i=0;i<=frame_ID_loop_count;i++){			
frame_id		8	uimsbf
frame_duration		32	upcrmsf
total_timeslot_count	5	11	uimsbf
start_timeslot_number	5	11	uimsbf
timeslot_loop_count		8	uimsbf
for(j=0;j<=timeslot_loop_count;j++){			
timeslot_frequency_offset		24	tcimsbf
timeslot_time_offset		32	upcrmsf
timeslot_id		8	uimsbf
repeat_count		8	uimsbf
}			
}			
CRC_32		32	rpchof
}			
NOTE:	*Reserved bits are of type bs1bf, and shall precede the Information bits on the same line. Unless otherwise stated, these bits shall be set to '0'.		

Semantics for the frame\_composition\_section:

- SI\_private\_section\_header: This is the standard SI private section header defined in table 11, and occupies a total of 64 bits.
- frame\_ID\_loop\_count: This is a 8-bit field indicating one less than the number of iterations of the frame loop that follows. A zero count indicates one loop.
- frame\_id: This 8-bit field serves as a label for identification of the  $i^{\text{th}}$  frame type from any other frame type.
- frame\_duration: This 32-bit field gives the duration of the  $i^{\text{th}}$  frame identified by the frame\_id, in terms of PCR count intervals. The 32 bits correspond to a maximum duration of 93,2 seconds.
- total\_timeslot\_count: This 11 bit field defines the total number of timeslots in the  $i^{\text{th}}$  frame.
- start\_timeslot\_number: This 11 bit field defines the number of the first timeslot of the  $i^{\text{th}}$  frame defined in this section, following the numbering scheme defined in subclause 6.7.2. This simplifies the partitioning of a frame across sections, since the definition of a non-homogenous frame format can exceed the capacity of one section.
- timeslot\_loop\_count: This is a 8-bit field indicating one less than the number of iterations of the timeslot loop that follows. A zero count indicates one loop. The order follows the numbering scheme defined in subclause 6.7.2, starting with the start\_timeslot\_number.
- timeslot\_frequency\_offset: This 24-bit field gives the signed value of the offset of the centre frequency of the  $j^{\text{th}}$  timeslot group relative to the frame\_centre\_frequency parameter (FCT). The frequency is given in multiples of 100 Hz.
- timeslot\_time\_offset: this 32 bit field gives the time offset of the  $j^{\text{th}}$  timeslot group relative to the frame start time, in terms of PCR count intervals. The 32 bits correspond to a maximum duration of 93,2 seconds.

- timeslot\_id: This 8 bit field identifies the type of timeslot for the  $j^{\text{th}}$  timeslot group, and corresponds to a timeslot\_id defined in the TCT.
- repeat\_count: This 8 bit field value is the number of repeats of the preceding timeslot\_id. The value is one less than the total number of successive timeslots of the same type. E.g. a value of 0 indicates no repeats (1 occurrence only), while a value of 2 indicates 2 further repeats for a total of 3. All repeats shall have the same timeslot\_frequency\_offset value, but the timeslot\_time\_offset value shall be increased from the prior timeslot by the timeslot\_duration value in the TCT for that timeslot\_id.
- CRC\_32: This is a 32-bit field that contains the CRC value that gives a zero output of the registers in the decoder defined in annex B of EN 300 468 [4] after processing the entire section.

#### 8.5.5.4 Timeslot Composition Table (TCT)

The TCT (see table 18) conveys information relating to the properties of the types of timeslot, such as TRF with ATM cells, TRF with MPEG2-TS packets, CSC, SYNC and ACQ. The timeslot\_id allows each kind of timeslot to be uniquely identified. Only timeslot properties are described within this table, the time and frequency co-ordinates of a particular timeslot are contained in the SCT/FCT.

**Table 18: Timeslot Composition Section**

Syntax	No. of bits		Information Mnemonic
	Reserved	Information	
Timeslot_composition_section(){			
SI_private_section_header		64	-
timeslot_loop_count		8	uimsbf
For(i=0; i <= timeslot_loop_count; i++){			
timeslot_id		8	uimsbf
symbol_rate		24	uimsbf
timeslot_duration		24	upcrmsf
burst_start_offset		16	upcrmsf
inner_code_type		1	bslbf
inner_code_ordering		1	bslbf
outer_coding		2	bslbf
inner_code_puncturing		4	bslbf
modulation		5	bslbf
baseband_shaping		3	bslbf
timeslot_payload_type		8	uimsbf
SAC_length	3	5	bslbf
request_flag		1	bslbf
M_and_C_flag		1	bslbf
group_ID_flag		1	bslbf
logon_ID_flag		1	bslbf
capacity_requests_number		3	bslbf
New_permutation		1	bslbf
If( (Inner_code_type == 1) && (New_permutation == 1) ) {			
P0	3	5	uimsbf
P1	6	10	uimsbf
P2	6	10	uimsbf
P3	6	10	uimsbf
}			
preamble_length		8	uimsbf
for (j=0; j < preamble_length; j++){			
preamble_symbol		2	bslbf
}			
while (!bytealigned){			
stuffing_bit }		1	bslbf
}			
CRC_32		32	rpchof
}			
NOTE: *Reserved bits are of type bslbf, and shall precede the Information bits on the same line. Unless otherwise stated, these bits shall be set to '0'.			

The TCT shall be segmented into timeslot composition sections using the syntax described in EN 300 468 [4]. Any sections forming part of a TCT shall be transmitted in TS packets with a PID value assigned in the PMT. Any sections of the TCT shall have the table\_id as defined in table 10.

Semantics for the timeslot\_composition\_section:

- SI\_private\_section\_header: This is the standard SI private section header defined in table 11, and occupies a total of 64 bits;
- timeslot\_loop\_count: This is a 8-bit field indicating one less than the number of iterations of the timeslot loop that follows;
- timeslot\_id: This is an 8-bit field which serves as a label for identification of the  $i^{\text{th}}$  timeslot including all properties of the slot;
- symbol\_rate: This 24 bit field gives the symbol rate in symbol/s;
- timeslot\_duration: This 24 bit field defines the duration of the timeslot, in terms of PCR count intervals. The 24 bits corresponds to a maximum duration of 364 ms;
- burst\_start\_offset: This 16 bit field defines the burst start offset from the slot start, in terms of PCR count intervals. The 16 bits corresponds to a maximum offset of 1,42 ms;
- inner\_code\_type: This bit specifies the inner coding type to use. A "0" means the K=7 convolutional code (see Section 0) and a "1" means the Turbo Code (see Section 0);
- inner\_code\_ordering: This bit specifies the order of transmission of the encoded block for Turbo Code. A "0" means the natural order, a "1" the reverse order (see Section 0). This bit is reserved for the K=7 convolutional code;
- outer\_coding: This is a 2-bit field specifying the outer coding mode to be used. The value "01" indicates that the RS code, as specified in Section 0, is switched on. The value "10" indicates that the CRC error detection code, as described in Section 0 is used. The value "00" indicates that both the RS and CRC codes are used. The value "11" indicates that neither the CRC nor the RS are used;
- inner\_code\_puncturing: This is a 4-bit field (table 19) specifying the inner coding rates to be used in combination with the selected inner coding;

**Table 19: Inner coding**

Code Rate	Value	K=7	Turbo
1 / 2	0000	x	x
2 / 3	0001	x	x
3 / 4	0010	x	x
5 / 6	0011	x	Not used
7 / 8	0100	x	Not used
1 / 3	0101	Not used	x
2 / 5	0110	Not used	x
4 / 5	0111	Not used	x
6 / 7	1000	Not used	x
Reserved for future use	1001 to 1110		
Inner code is omitted	1111	x	x
NOTE: x means supported.			

- modulation: This is a 5-bit field which serves as an identifier of the modulation scheme. The value "0 0000" is not defined. The value "0 0001" indicates QPSK. The values "0 0010" to "1 1111" are reserved for future use;
- baseband\_shaping: This is a 3-bit field which serves as an identifier of the used baseband shaping. The value 000 indicates a root raised cosine filtering with roll off factor 0.35. Other values are reserved for future use;



**Table 20: Timeslot Types**

Timeslot Type	Value
Reserved	0x00
TRF with one ATM cell	0x01
TRF with two ATM cells	0x02
Reserved	0x03
TRF with four ATM cells	0x04
TRF with MPEG2-TS packet(s)	0x05
CSC	0x06
ACQ	0x07
SYNC	0x08
Reserved for future use	0x09 to 0xFF

- timeslot\_payload\_type: This is an 8-bit field which serves as a label for identification of the type of timeslot. The value are assigned in table 20;
- SAC\_length: This is a 5-bit field specifying the length of the SAC field in bytes;
- Request\_flag: The value '1' indicates that the SAC field of this time slot contains at least one Request sub-field as described in subclause 6.6.1.2. In this case, the number of requests is indicated with the capacity\_requests\_number field described herebelow. The value '0' indicates that there is no Request sub-field;
- M\_and\_C\_flag: The value '1' indicates that the SAC field of this time slot contains an M&C sub-field as described in subclause 6.6.1.2. The value '0' indicates that there is no M&C sub-field;
- group\_ID\_flag: The value '1' indicates that the SAC field of this time slot contains a group\_ID sub-field as described in subclause 6.6.1.2. The value '0' indicates that there is no group\_ID sub-field;
- logon\_ID\_flag: The value '1' indicates that the SAC field of this time slot contains a logon\_ID sub-field as described in subclause 6.6.1.2. The value '0' indicates that there is no logon\_ID sub-field;
- capacity\_requests\_number: this 3-bit field indicates one less than the number of capacity\_requests allowed in the SAC as described in subclause 6.6.1.2. In networks using the TRF-prefix signalling method (capacity requests appended to TRF bursts), this field would normally have a value of 0 ('000'). This field is reserved if the request flag is equal to '0';
- New\_permutation: when set to 0, this bit specifies that the turbo code permutation to be used is the default one as defined in table 5. When set to 1, this bit indicates that the permutation parameters are defined in the sequel of the TCT. This bit is reserved if the turbo code is not used (inner\_code\_type==0).
- P0, P1, P2 and P3: define the set of permutation parameters to be used for the specific burst type.
- preamble\_length: This is an 8-bit field specifying the preamble length in symbols. The loop which follows is aligned however on byte boundaries. This means that for example if the preamble length is 14 Symbols, the loop over the preamble bytes will consist of 4 bytes with the last 4 bits being merely stuffing bits;
- preamble\_symbol: This is a 2-bit field specifying a QPSK symbol of the preamble. Symbols are given in the order they shall be transmitted. The first bit specifies the I bit and the second bit the Q bit of a QPSK symbol. As the preamble is not scrambled, a proper sequence shall be selected in order to comply with ETSI requirements concerning off-axis EIRP;
- stuffing\_bit: Since the preamble description is byte aligned, stuffing bits are present until the next byte boundary. The stuffing bits may take any value and shall be discarded by the RCST;
- CRC\_32: This is a 32-bit field that contains the CRC value that gives a zero output of the registers in the decoder defined in annex B of EN 300 468 [4] after processing the entire section.

### 8.5.5.5 Satellite Position Table (SPT)

The SPT conveys information about the positions of all the satellites used for the forward and return links. In the case where the NCC computes the nominal RCST-to-satellite range in 27 MHz clock periods at RCST registration, the SPT table needs not be transmitted.

The SPT shall be segmented into satellite position sections using the syntax described in EN 300 468 [4]. Any sections forming part of an SPT shall be transmitted in TS packets with a PID value assigned in the PMT. Any sections of the SPT shall have the table\_id value as defined in table 10. The SPT is defined in table 21.

**Table 21: Syntax of satellite position table section**

Syntax	No. of bits		Mnemonic
	Reserved	Information	
satellite_position_section(){			
SI_private_section_header		64	-
<b>satellite_loop_count</b>		8	uimsbf
for(i=0;i<=satellite_loop_count;i++){			
satellite_id		8	uimsbf
x_coordinate		32	spfmsbf
y_coordinate		32	spfmsbf
z_coordinate		32	spfmsbf
}			
CRC_32		32	rpchof
}			
NOTE: *Reserved bits are of type bslbf, and shall precede the Information bits on the same line. Unless otherwise stated, these bits shall be set to '0'.			

Semantics for the satellite\_position\_section:

- SI\_private\_section\_header: This is the standard SI private section header defined in table 11, and occupies a total of 64 bits;
- satellite\_loop\_count: This 8 bit field defines one less than the number of satellites belonging to the system. A zero count indicates one loop;
- satellite\_id: This 8 bit field defines a system assigned satellite identifier;
- x\_coordinate: This 32 bit field defines the x co-ordinate of the satellite ephemeris in meters;
- y\_coordinate: This 32 bit field defines the y co-ordinate of the satellite ephemeris in meters;
- z\_coordinate: This 32 bit field defines the z co-ordinate of the satellite ephemeris in meters.

NOTE: The position of the satellites will be expressed as Cartesian coordinates x, y, z in the geodetic reference frame ITRF96 (IERS Terrestrial Reference Frame). This system coincides with the WGS84 (World Geodetic System 84) reference system at the one meter level.

- CRC\_32: This is a 32 bit field that contains the CRC value that gives a zero output of the registers in the decoder defined in annex B of EN 300 468 [4] after processing the entire section.

### 8.5.5.6 PCR Insertion Transport Stream packet

The PCR Insertion transport stream (TS) packet is used to insert the NCR count for return link synchronization purpose.

#### 8.5.5.6.1 TS packet format

The format of the TS packet shall be as defined in ISO/IEC 13818-1 [9] subclause 2.4.3. The value of the fields in the TS packet (ISO/IEC 13818-1 [9] table 2.3) shall be as follows:

- sync\_byte: as defined by the standard;
- transport\_error\_indicator: set to '0';

- payload\_unit\_start\_indicator: If the payload of the TS packet contains the PCR Insertion TS packet payload section, the payload\_unit\_start\_indicator value shall be '1', indicating that the first byte of the payload of this Transport Stream packet carries a pointer field. This pointer field is an 8-bit field whose value shall be the number of bytes, immediately following the pointer field until the first byte of the section that is present in the payload of the Transport Stream packet (so a value of 0x00 in the pointer field indicates that the section starts immediately after the pointer field);
- transport\_priority: set to '0';
- PID: set to the assigned PID value;
- transport\_scrambling\_control: as defined by the standard, depending on whether or not transport stream scrambling is used;
- adaptation\_field\_control field set to '10' whenever a PCR value only is being sent (no payload) and set to '11' to indicate that the packet contains an adaptation field followed by payload to indicate that an update to one or more of the forward or return link combinations data is present;
- adaptation\_field(): present – see adaptation field definition below (subclause 8.5.5.6.2),
- data\_byte [loop]: present if adaptation\_field\_control field set to '11' - see payload definition below (subclause 8.5.5.6.3).

#### 8.5.5.6.2 Adaptation field

The adaptation field shall be coded as per ISO/IEC 13818-1 [9] subclause 2.4.3.4, with field values as follows:

- adaptation\_field\_length: set to 7;
- discontinuity\_indicator: set to '0';
- random\_access\_indicator: set to '0';
- elementary\_stream\_priority\_indicator: set to '0';
- PCR\_flag: set to '1';
- OPCR\_flag: set to '0';
- splicing\_point\_flag: set to '0';
- transport\_private\_data\_flag: set to '0';
- adaptation\_field\_extension\_flag: set to '0';
- program\_clock\_reference\_base; program\_clock\_reference\_extension: set to the inserted NCR value, as defined in the standard.

The remaining fields of the adaptation field are not applicable.

#### 8.5.5.6.3 Optional Payload field

The optional payload of the TS packet may contain the delays between NCCs and satellites as well as between Traffic Gateways and satellites. The corresponding syntax shall be conforming to the one defined in the standard ISO/IEC 13818-1 [9] for the definition of private sections. The section shall be entirely contained in the payload of the TS packet containing the NCR.

If the payload is used to transmit delay information, the syntax in table 22 shall be used.

**Table 22: Syntax of the optional PCR Insertion TS packet payload section**

Syntax	No of bits		Information mnemonic
	reserved	information	
PCR Insertion TS packet payload section () {			
SI_private_section_header		64	-
forward_link_combinations		8	uimsbf
for (i=0; i<forward_link_combinations; i++){			
satellite_id		8	uimsbf
NCC_id		8	uimsbf
propagation_delay }		32	upcrmsf
return_link_combinations		8	uimsbf
for (i=0; i<return_link_combinations; i++) {			
satellite_id		8	uimsbf
gateway_id		8	uimsbf
propagation_delay}		32	upcrmsf
descriptor_length		8	uimsbf
for(i=0; i<N; i++) {			
descriptor() }			
CRC_32		32	rpchof
}			

- SI\_private\_section\_header: This is the standard SI private section header defined in table 11, and occupies a total of 64 bits.
- forward\_link\_combinations : This 8 bit field defines one less than the number of all applicable NCC to satellite combinations on the forward link. A zero count indicates one loop.
- satellite\_id: This 8 bit field defines the identifier of the satellite end of one combination. The satellite\_id is assigned by the system administrator.
- NCC\_id: This 8 bit field defines the identifier of the NCC end of one combination. The NCC\_id is assigned by the system administrator.
- propagation\_delay: This 32 bit field defines the propagation\_delay between NCC and satellite as a PCR count. For the forward link it is the delay from NCC to satellite while for the return link it is the delay from satellite to Gateway. The 32 bits corresponds to a maximum delay of 93.2 seconds. RCSTs may use this information to compute delays.
- return\_link\_combinations: This 8 bit field defines one less than the number of all applicable satellite to Gateway combinations on the return link. A zero count indicates one loop.
- Gateway\_id : This 8 bit field defines the identifier of the Gateway end of one combination. The Gateway\_id is assigned by the system administrator.
- descriptor\_length: This 8 bit field gives the total length in bytes of future descriptors.
- CRC\_32: This is a 32 bit field that contains the CRC value that gives a zero output of the registers in the decoder defined in annex B of EN 300 468 [4] after processing the entire section.

### 8.5.5.7 Terminal Burst Time Plan (TBTP)

The TBTP defines the dynamic timeslot assignments generated every superframe.

The TBTP shall be as defined in table 23. It shall be segmented into TBTP sections using the syntax described in EN 300 468 [4]. Any sections forming part of an TBTP shall be transmitted in TS packets with a PID value assigned in the PMT. Any sections of the TBTP shall have the table\_id as defined in table 10.

Table 23: Terminal Burst Time Plan Section

Syntax	No. of bits		Mnemonic
	Reserved	Information	
Terminal_burst_time_plan(){			
SI_private_section_header		64	-
Group_ID		8	uimsbf
Superframe_count		16	uimsbf
frame_loop_count	3	5	uimsbf
for (i=0;i<=frame_loop_count;i++) {			
frame_number	3	5	uimsbf
BTP_loop_count	5	11	uimsbf
for (j=0;j<= BTP_loop_count;j++){			
logon_id		16	uimsbf
Multiple_channels_flag		1	bslbf
Assignment_type		2	bslbf
Start_slot	2	11	uimsbf
If (Multiple_channels_flag == 1)			
Channel_ID	4	4	uimsbf
Assignment_count}		8	uimsbf
}			
CRC_32		32	rpchof
}			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. Unless otherwise stated, these bits shall be set to '0'.		

Semantics for the terminal\_burst\_time\_plan:

- SI\_private\_section\_header: This is the standard SI private section header defined in table 11, and occupies a total of 64 bits;
- Group\_ID: This is an 8 bit field that identifies the service group for which the message applies. See subclause 8.2;
- Superframe\_count: This 16 bit field identifies the modulo 65,536 superframe count to which the BTP applies;
- Frame\_loop\_count: This 5 bit field specifies one less than the number of superframe frame loops that follow. A zero count indicates one loop. Each entry in the loop corresponds to the definition of one frame in the superframe;
- frame\_number: This 5 bit field specifies the frame number within the superframe, using the frame numbering as defined in subclause 6.7.2;
- BTP\_loop\_count: This 11-bit field specifies one less than the number of the frame BTP entry loops that follow. A zero count indicates one loop. Each entry in the loop provides a definition of a block of timeslots allocated. Where a given terminal is assigned multiple blocks, there will be one entry per block. The entries for a given terminal may not be consecutive in the section;
- Logon\_ID: This 16 bit field gives the identifier assigned to the terminal at logon time. This number is limited in scope to within the Group\_ID (i.e. the same logon\_ID can be given to a different terminal within a different Group\_ID). See subclause 8.2;
- Multiple\_channels\_flag: this 1-bit flag indicates the presence of a Channel\_ID field for the current assignment. A value of 1 corresponds to the Channel\_ID field being present. If the field is equal to 0, channel\_id=0 shall be assumed by the RCST;
- Assignment\_Type: This 2 bit field defines the repetitive nature of the assignment, as defined in table 24;

**Table 24: Assign Type**

Assign type	Value
one time assignment	00
repeating assignment	01
assignment release	10
Reserved	11

- Start\_slot: This 11 bit field gives the number of the first timeslot in the block, relative to the frame (numbering as defined in subclause 6.7.2);
- Channel\_ID: This 4-bit field indicates the channel to which timeslots are being assigned. Values are managed by the NCC. This field is present if the multiple\_channels\_flag is set to '1';
- Assignment\_count: This 8 bit field gives one less than the number of timeslots assigned in the block. These timeslots will be consecutively numbered in ascending order starting with the start\_slot;
- CRC\_32: This is a 32 bit field that contains the CRC value that gives a zero output of the registers in the decoder defined in annex B of EN 300 468 [4] after processing the entire section.

### 8.5.5.8 Terminal Information Message (TIM)

The TIM is used to transmit static or quasi-static information about the forward link, such as configuration parameters. The format supports two variants of this message:

- a terminal specific (uni-cast) message containing information for a specific terminal. This variant may be encrypted for security reasons. It is sent during logon initialization, and when a parameter changes;
- a broadcast message giving general information applicable to all terminals. This variant is unlikely to be encrypted. It will need to be transmitted sufficiently often that newly powered terminals can acquire any necessary information within a reasonable time. This matches a similar requirement for the SCT, FCT and TCT signals, suggesting the same repeat interval (see subclause 8.3.5).

**Table 25: Terminal Information Message Section**

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
terminal_information_message_section() {			
DSM-CC_private_section_header		96	
If MAC_address == RCST MAC (not broadcast) {			
RCST_Status }		8	flagmsf
else {			
Network_status }		8	flagmsf
descriptor_loop_count		8	uimsbf
for (i= 0; i<= descriptor_loop_count; i++) {			
descriptor() }			
If MAC_address == RCST MAC (not broadcast) {			
Pad_bytes }		see text	bslbf
CRC_32		32	rpchof
}			
NOTE: *Reserved bits are of type bslbf, and shall precede the Information bits on the same line. They shall be ignored by the RCST. The value of these bits is undefined. For an encrypted uni-cast TIM, the bit values shall be varied in a random manner to avoid encryption spoofing.			

The TIM shall be as defined in table 25. It shall be segmented into terminal information message sections using the syntax described in EN 300 468 [4]. Any sections forming part of a TIM shall be transmitted in TS packets with a PID value assigned in the PMT. Any sections of the TIM shall have the table\_id as described in table 10.

Semantics for the terminal\_information\_message\_section:

- DSM-CC\_private\_section\_header: This is the standard DSM-CC private section header defined in table 12, and occupies a total of 96 bits;
- RCST\_Status: This 8 bit field gives status flags defining the network state of the RCST (see table 26);

**Table 26: RCST Status**

Bit	Identifier
(MSB) 7	ID_encrypt
6	Logon_fail_(busy)
5	Logon_denied
4	Log_off
3	Transmit_Disable
2	Rain_Fade_release
1	Rain_Fade_detect
(LSB) 0	Wake_up

Semantics for the flag bits are as follows, where a logic '1' asserts the condition defined:

- ID\_encrypt: Indicates that the terminal shall use TBTP logon ID encryption;
- Logon\_fail\_(busy): Indicates that the terminal cannot enter the network because of lack of resources;
- Logon\_denied: Indicates that the terminal is not authorized to enter the network;
- Log\_off: Indicates that the network has enforced a terminal logoff. The terminal shall immediately cease transmission and enter the logged off state;
- Transmit\_Disable: Indicates that the terminal shall immediately cease transmission and enter the hold state, and shall not resume until it receives a uni-cast TIM with this bit reset to '0';
- Rain\_Fade\_release: Indicates that the NCC is performing a reconfiguration procedure to restore settings following cessation of a rain fade event;
- Rain\_Fade\_detect: Indicates that the NCC has detected a rain fade event and is performing a reconfiguration procedure to establish rain fade settings;
- Wake\_up: indicates that the NCC wants to wake up the RCST;
- Network\_Status: This 8 bit field gives status flags defining the network state for RCSTs within the scope of the broadcast MAC address. The flag bits shall be as defined in table 27.

**Table 27: Network Status**

Bit	Identifier
(MSB) 7	ID_encrypt
6	Reserved
5	Reserved
4	CSC_link_failure_rec
3	Link_failure_recovery
2	Return_link_failure
1	NCC_Receive_Failure
(LSB) 0	Scheduler_Failure

Semantics for the flag bits are as follows, where a logic '1' asserts the condition defined:

- ID\_encrypt: Indicates that terminals shall use TBTP logon ID encryption;
- CSC\_link\_failure\_rec: Indicates that the system is recovering from a failure of a forward or return link. The NCC affects CSC bursts in the TBTP. Terminals shall wait for their reserved timeslot or use one of the CSC contention timeslots;

- **Link\_failure\_recovery**: Indicates that the system is recovering from a failure of a forward or return link. Terminals shall follow a pre-defined procedure for large outage recovery until a new broadcast TIM is received with this bit reset to '0';
- **Return\_link\_failure**: Indicates that the NCC has detected a failure of the return link. All RCST's shall cease transmission until a new broadcast TIM is received with this bit reset to '0';
- **NCC\_Receive\_Failure**: Indicates that the NCC has a receive channel failure. Affected terminals shall suspend transmission until a new broadcast TIM is received indicating that the failure is no longer present;
- **Scheduler\_Failure**: Indicates that the NCC has a scheduler failure. Affected terminals shall suspend transmission of all traffic, but not SYNC bursts, until a new broadcast TIM is received indicating that the failure is no longer present. SYNC bursts transmitted while the fault is present shall have the request field of the SAC set to '0';
- **descriptor\_loop\_count**: This 8 bit field defines one less than the number of descriptors in the following loop. A zero count indicates one loop;
- **descriptor()**: The descriptors that may be inserted into the TIM are defined in subclause 8.5.5.10. It should be noted that the allowed descriptors are different between the unicast and broadcast versions of the TIM (as defined by the scope of the MAC address);
- **Pad\_bytes**: Each Pad\_byte shall be an 8 bit field. Sufficient Pad\_bytes are inserted so as to pad the encrypted portion of the message to the encryption algorithm boundary. The content of these fields is undefined and shall be randomized to prevent code spoofing. The encrypted portion of the message starts immediately following the MAC Address fields and finishes immediately prior to the CRC\_32 field;
- **CRC\_32**: This is a 32 bit field that contains the CRC value that gives a zero output of the registers in the decoder defined in annex B of EN 300 468 [4] after processing the entire section.

#### 8.5.5.9 Correction Message Table (CMT)

The CMT shall be as shown in table 28. It shall be segmented into terminal information message sections using the syntax described in EN 300 468 [4]. Any sections forming part of a CMT shall be transmitted in TS packets with a PID value assigned in the PMT. Any sections of the CMT shall have the table\_id value as defined in table 10. This message provides closed loop feedback to a number of terminals to allow them to adjust the transmit power level, frequency and burst timing as required to maintain link integrity, as well as other error condition feedback messages.



Table 28: Correction Message Table Section

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
correction_message_table_section(){			
SI_private_section_header		64	-
Entry_loop_count		8	uimsbf
for (i = 0; i <= Entry_loop_count; i++) {			
Group_ID		8	uimsbf
Logon_ID		16	uimsbf
Time_correction_flag		1	bslbf
Power_correction_flag		1	bslbf
Frequency_correction_flag		1	bslbf
Slot_Type		2	bslbf
Burst_time_scaling		3	uimsbf
if (Time_correction_flag == 1)			
Burst_time_correction		8	tcimsbf
if (Power_correction_flag == 1)			
Power_control_flag		1	bslbf
if(Power_control_flag==1)			
Power_correction		7	tcimsbf
else			
EbN0		7	tcimsbf
if (Frequency_correction_flag == 1)			
Frequency_correction}		16	tcimsbf
CRC_32		32	rpchof
}			
NOTE: *Reserved bits are of type bslbf, and shall precede the Information bits on the same line. Unless otherwise stated, these bits shall be set to '0'.			

Semantics for the correction\_message\_table\_section:

- SI\_private\_section\_header: This is the standard SI private section header defined in table 11, and occupies a total of 64 bits;
- Entry\_loop\_count: This field specifies one less than the number of correction message loops that follow. A zero count indicates one loop;
- Group\_ID: This 8 bit field defines which Group ID the RCST is assigned to, as identified by the Terminal Information Message (TIM). This matches the Group\_ID used in the TBTP sections (see subclause 8.5.5.7);
- Logon\_ID: This 16 bit field identifies the assigned terminal logon identifier, as identified by the TIM. Nominally, this is the same as used in the TBTP sections. When scrambling is used, the Logon ID used in the CMT will be the unscrambled version, to avoid a possible compromise to the TBTP security;
- Time\_correction\_flag; Power\_correction\_flag; Frequency\_correction\_flag; Slot\_type; Burst\_time\_scaling; Burst\_time\_correction; Power\_control\_flag; Power\_correction; EbN0; Frequency\_correction: These fields are identical to the matching fields of the Correction\_message\_descriptor, and define one measurement correction set. See subclause 8.5.5.10.3 for a definition of the descriptor;

NOTE: The CMT does not use the descriptor directly, since the added descriptor\_tag and descriptor\_length parameters would represent a substantial, and unnecessary, overhead.

- CRC\_32: This is a 32 bit field that contains the CRC value that gives a zero output of the registers in the decoder defined in annex B of EN 300 468 [4] after processing the entire section.

## 8.5.5.10 Descriptor coding

### 8.5.5.10.1 Descriptor Identification and location

Table 29 lists the descriptors defined within the present document, giving the descriptor-tag values and the intended placement within the tables.

**Table 29: Descriptor Tags and Locations**

Descriptor	Tag value	PMT	RMT	TIM*	
				Broadcast	Unicast
Network_layer_info_descriptor	0xA0				X
Correction_message_descriptor	0xA1				X
Logon_initialize_descriptor	0xA2				X
ACQ_assign_descriptor	0xA3				X
SYNC_assign_descriptor	0xA4				X
Encrypted_logon_ID_descriptor	0xA5				X
Echo_value_descriptor	0xA6			X	X
Linkage_descriptor (private data)**	0x4A		X		
RCS_content_descriptor	0xA7	X			
Satellite_forward_link_descriptor	0xA8		X		X
Satellite_return_link_descriptor	0xA9		X		X
table_update_descriptor	0xAA			X	
Contention_control_descriptor	0xAB			X	
Correction_control_descriptor	0xAC			X	
NOTE 1: * For the TIM message sections, the definition of intended placement differs between the unicast and broadcast versions of that message.					
NOTE 2: ** Private extension to existing DVB descriptor.					

After logon, the unicast TIM contains at least the following descriptors: Correction\_message\_descriptor, Logon\_initialize\_descriptor, SYNC\_assign\_descriptor and Satellite\_return\_link\_descriptor.

### 8.5.5.10.2 Network Layer Info Descriptor (optional)

The Network Layer Info descriptor provides a mechanism by which network level information, such as the meta-signalling VPI/VCI and base IP address, can be passed to the Management Plane of the RCST during, or prior to, the start-up configuration phase of logon. As such, the message content is passed transparently through the lower layers covered by the present document and is not defined here. The descriptor is defined in table 30.

**Table 30: Network Layer Info Descriptor**

Syntax	No. of bits	Mnemonic
Network_layer_info_descriptor(){		
descriptor_tag	8	uimsbf
descriptor_length	8	uimsbf
Message_body	see text	
}		

Semantics for the Network\_layer\_info\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- Message\_body: This variable length field shall contain a datagram for passing to the target application. The length of the message body shall not exceed 255 bytes, and should preferably be limited such that the section containing this descriptor fits within a single TS packet. This datagram will take the form of an SNMP message. The messages that can be passed by this method are beyond the scope of the present document. Basic functionalities related with the network are provided with the logon initialize descriptor (subclause 8.5.5.10.4). This descriptor is optional, the RCST indicates in the CSC burst (see subclause 6.2.3) if it implements SNMP.

### 8.5.5.10.3 Correction Message descriptor

The Correction Message descriptor defines a transmit parameter correction set for one terminal measurement. It shall be as defined in table 31.

**Table 31: Correction Message Descriptor**

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
Correction_message_descriptor(){			
descriptor_tag		8	uimsbf
descriptor_length		8	uimsbf
Time_correction_flag		1	bslbf
Power_correction_flag		1	bslbf
Frequency_correction_flag		1	bslbf
Slot_Type		2	bslbf
Burst_time_scaling		3	uimsbf
If (Time_correction_flag == 1)			
Burst_time_correction		8	tcimsbf
If (Power_correction_flag == 1)			
Power_control_flag		1	bslbf
If(Power_control_flag==1)			
Power_correction		7	tcimsbf
else			
EbN0		7	tcimsbf
If (Frequency_correction_flag == 1)			
Frequency_correction		16	tcimsbf
}			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. For non-encrypted message types, these bits shall be set to '0'. For encrypted message types, the value of these bits is undefined.		

Semantics for the Correction\_message\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- Time\_correction\_flag, Power\_correction\_flag, Frequency\_correction\_flag: these three bits are used to indicate the presence of time, power and frequency correction fields, respectively, in the remainder of the descriptor;
- Slot\_type: This 2 bit field identifies the type of burst being measured, as defined in table 32;

**Table 32: Slot Type**

Value	Identifier
00	TRF
01	CSC
10	ACQ
11	SYNC

- Burst\_time\_scaling: This 3 bit field gives the power-of-2 scaling to apply to the Burst\_time\_correction parameter, i.e. a value of 2 indicates a scaling factor of 4 (= shift left 2 bits). In case there is no time correction in this descriptor, i.e. the Time\_correction\_flag is equal to 0, the Burst\_time\_scaling field shall be set to 000;
- Burst\_time\_correction: This 8 bit field gives the required correction to burst timing as a two's complement binary PCR clock count (i.e. in counts of the 27 MHz PCR clock) that shall be scaled according to the Burst\_time\_scaling field above. To minimize truncation errors, the N LSB bits of the scaled value shall be set to an approximate mid-range value of '1' followed by '0's, with N being the value of the Burst\_time\_scaling field. For example, with N=2, the resulting clock count value is "dd dddd dd10";

- Power\_control\_flag: This 1 bit field defines how uplink power control for the return link is carried out. The value '1' indicates that the NCC transmits a power correction value for the RCST and the value '0' indicates that the NCC transmits a measured  $E_b/N_0$  value instead;
- Power\_correction: This 7 field gives the required correction to uplink power on the return link in 0.5 dB steps as a two's complement integer value;
- EbN0: This 7 bit field gives the measured  $E_b/N_0$  value on the return link in 0.5 dB steps as two's complement integer value. The RCST can use the value to monitor the return link. If available, RCSTs should use this field for information only;
- Frequency\_correction: This 16 bit field gives the required correction to frequency in 10 Hz steps, as a two's complement integer value. A negative value indicates that the terminal shall reduce frequency. For systems not implementing frequency correction, this field shall be set to all 0's.

#### 8.5.5.10.4 Logon Initialize Descriptor

This descriptor provides parameters needed for initial logon (see table 33).

**Table 33: Logon Initialize Descriptor**

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
Logon_initialize_descriptor() {			
descriptor_tag		8	uimsbf
descriptor_length		8	uimsbf
Group_ID		8	uimsbf
Logon_ID		16	uimsbf
Security_handshake_required	3	1	bslbf
Prefix_flag		1	bslbf
Data_unit_labelling_flag		1	bslbf
Mini_slot_flag		1	bslbf
Contention_based_mini_slot_flag		1	bslbf
Traffic_burst_type	2	1	bslbf
If (Traffic_burst_type == 0) {			
Connectivity		1	bslbf
If (Connectivity == 0) {			
Return_VPI	4	8	uimsbf
Return_VCI		16	uimsbf
}			
Else {			
Return_signalling_VPI	4	8	uimsbf
Return_signalling_VCI		16	uimsbf
Forward_signalling_VPI	8	8	uimsbf
Forward_signalling_VCI		16	uimsbf
}			
}			
Else {			
Return_TRF_PID		13	uimsbf
If (Data_unit_labelling_flag == 1) {			
Return_CTRL_MNGM_PID	3	13	uimsbf
}			
}			
}			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. They shall be ignored by the RCST. The value of these bits is undefined. For an encrypted uni-cast TIM, the bit values shall be varied in a random manner to avoid encryption spoofing.		

Semantics for the Logon\_initialize\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- Group\_ID: This 8 bit field defines which Group ID the terminal is assigned to. This matches the Group\_ID used in the TBTP and CMT sections;
- Logon\_ID: This 16 bit field identifies the assigned terminal logon identifier, which is used in the TBTP and CMT sections;
- Security\_handshake\_required: the value '1' indicates that the security handshake as described in subclause 9.4 is to be used. Otherwise it is set to '0';
- Prefix\_flag: The value '1' indicates that the Prefix Method according to subclause 6.6.1 is implemented in the network and the value '0' indicates that it is not implemented;
- Data\_unit\_labelling\_flag: The value '1' indicates that the Data Unit Labelling Method according to subclause 6.6.2 is implemented in the network and the value '0' indicates that it is not implemented;
- Mini\_slot\_flag: The value '1' indicates that the Mini-slot Method according to subclause 6.6.3 is implemented in the network and the value '0' indicates that it is not implemented;
- Contention\_based\_mini\_slot\_flag: The value '1' indicates that the Contention-based Mini-slot Method according to subclause 6.6.4 is implemented in the network and the value '0' indicates that it is not implemented;
- Traffic\_burst\_type: This 1 bit field defines the traffic burst type to be used on the return link. The value '0' indicates ATM TRF according to subclause 6.2.1.1 and the value '1' indicates optional MPEG2-TS TRF according to subclause 6.2.1.2;
- Connectivity: This 1 bit field defines the connectivity to be used. The value '0' indicates IP connectivity according to subclause 8.1.1 (Type A RCST) and the value '1' indicates optional ATM connectivity according to subclause 8.1.2 (Type B RCST). In the case of MPEG2-TS TRF the connectivity is always IP and therefore not signalled;
- Return\_VPI, Return\_VCI: These fields define the VPI/VCI that the RCST shall use in ATM cells on the return link;
- Return\_signalling\_VPI, Return\_signalling\_VCI: These fields define the VPI/VCI that is used on the return link for ITU-T Recommendation Q.2931 [11] signalling instead of the normal value 0/5. The signalling is used to set up connections for traffic. These parameters can be the same for all RCSTs;
- Forward\_signalling\_VPI, Forward\_signalling\_VCI: These fields define the VPI/VCI that is used on the forward link for ITU-T Recommendation Q.2931 [11] signalling instead of the normal values 0/5. The signalling is used to set up connections for traffic;
- Return\_TRF\_PID: This 13 bit field defines the PID that the RCST shall use in optional MPEG2 TS packets on the return link for traffic information. This parameter can be the same for all RCSTs;
- Return\_CTRL\_MNGM\_PID: This 13 bit field defines the PID that the RCST shall use in optional MPEG2 TS packets on the return link for CTRL/MNGM information. This parameter can be the same for all RCSTs.

### 8.5.5.10.5 ACQ Assign Descriptor

This descriptor (table 34) defines the assignment of Acquisition bursts for the optional coarse synchronization procedure at logon. If used, it will only be present following detection of a valid CSC from the terminal. Some systems may choose to assign acquisition timeslots via the TBTP instead of using this descriptor.

**Table 34: ACQ Assign Descriptor**

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
ACQ_assign_descriptor() {			
descriptor_tag		8	uimsbf
descriptor_length		8	uimsbf
ACQ_achieved_time_threshold		8	uimsbf
ACQ_achieved_frequency_threshold		16	uimsbf
ACQ_repeats	2	6	uimsbf
ACQ_start_superframe		16	uimsbf
ACQ_frame_number	3	5	uimsbf
ACQ_repeat_period	2	6	uimsbf
ACQ_slot_number	5	11	uimsbf
}			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. They shall be ignored by the RCST. The value of these bits is undefined. For an encrypted uni-cast TIM, the bit values shall be varied in a random manner to avoid encryption spoofing.		

Semantics for the ACQ\_assign\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- ACQ\_achieved\_time\_threshold: This 8 bit field gives the magnitude of the burst timing error threshold value to be used for the transition from the coarse synchronization procedure to the "ready for fine sync" state, see subclause 7.1. The value is scaled by the same amount as the Burst\_time\_correction value of the Correction Message Table and Correction Message Descriptor, allowing a direct magnitude comparison with that value. Transition to the "ready for fine sync" state occurs when the magnitude of the measured Burst\_time\_correction value is less than, or equal to, the ACQ\_achieved\_time\_threshold;
- ACQ\_achieved\_frequency\_threshold: This 16 bit field gives the magnitude of the carrier frequency error threshold value to be used for the transition from the coarse synchronization procedure to "ready for fine sync" state, see subclause 7.1. Transition to the "ready for fine sync" state occurs when the magnitude of the measured Frequency\_correction value is less than, or equal to, the ACQ\_achieved\_frequency\_threshold;

NOTE: If both ACQ\_achieved\_time\_threshold and ACQ\_achieved\_frequency\_threshold have non-zero values, transition to the "ready for fine sync" state occurs when both criterion are fulfilled.

- ACQ\_repeats: This 6 bit field defines the maximum number of times that the ACQ shall be repeated during the Coarse Synchronization procedure. A special value of '00 0000' is reserved to indicate that the terminal is being given a sustained acquisition burst. This special value is reserved for use during terminal alignment procedures, typically at the time of installation;
- ACQ\_start\_superframe: This 16 bit field gives the modulo 65536 superframe number at which the ACQ assignment starts (see SCT Superframe\_counter parameter in subclause 8.5.5.2);
- ACQ\_frame\_number: This 5 bit field identifies which frame number in the superframe contains the ACQ burst. This number matches the frame numbering defined in subclause 6.7.2;
- ACQ\_repeat\_period: This 6 bit field gives the number of superframes between ACQ repeats;
- ACQ\_slot\_number: This 11 bit field identifies the timeslot number to use for ACQ. This number matches the numbering defined in subclause 6.7.2.

### 8.5.5.10.6 SYNC Assign Descriptor

This descriptor (table 35) provides the static BTP assignment of SYNC bursts. It is typically sent following the detection of a valid CSC from the terminal, but may also be sent when it is necessary to modify the SYNC assignment.

**Table 35: SYNC Assign Descriptor**

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
SYNC_assign_descriptor() {			
descriptor_tag		8	uimsbf
descriptor_length		8	uimsbf
SYNC_achieved_time_threshold		8	uimsbf
max_SYNC_tries		8	uimsbf
SYNC_achieved_frequency_threshold		16	uimsbf
SYNC_start_superframe		16	uimsbf
SYNC_frame_number	3	5	uimsbf
SYNC_repeat_period		16	uimsbf
SYNC_slot_number	5	11	uimsbf
}			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. They shall be ignored by the RCST. The value of these bits is undefined. For an encrypted uni-cast TIM, the bit values shall be varied in a random manner to avoid encryption spoofing.		

Semantics for the SYNC\_assign\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- SYNC\_achieved\_time\_threshold: This 8 bit field gives the magnitude of the burst timing error threshold value to be used for the transition from the fine synchronization procedure to the "fine sync" state, see subclause 7.1. The value is scaled by the same amount as the Burst\_time\_correction value of the Correction Message Table and Correction Message Descriptor, allowing a direct magnitude comparison with that value. Transition to the "fine sync" state occurs when the magnitude of the measured Burst\_time\_correction value is less than, or equal to, the SYNC\_achieved\_time\_threshold;
- max\_SYNC\_tries: This 8 bit field gives the maximum number of bursts that the RCST can send to achieve fine synchronization, from the start of the "ready for fine sync" state (see subclause 7.1);
- SYNC\_achieved\_frequency\_threshold: This 16 bit field gives the magnitude of the carrier frequency error threshold value to be used for the transition from the fine synchronization procedure to the "fine sync" state, see subclause 7.1. Transition to the "fine sync" state occurs when the magnitude of the measured Frequency\_correction value is less than, or equal to, the SYNC\_achieved\_frequency\_threshold;

NOTE: If both SYNC\_achieved\_time\_threshold and SYNC\_achieved\_frequency\_threshold have non-zero values, transition to the "fine sync" state occurs when both criterion are fulfilled.

- SYNC\_start\_superframe: This 16 bit field gives the modulo 65536 superframe number at which the SYNC assignment starts (see SCT Superframe\_counter parameter in subclause 8.5.5.2);
- SYNC\_frame\_number: This 5 bit field identifies which frame number in the superframe contains the SYNC burst. This number matches the frame numbering defined in subclause 6.7.2;
- SYNC\_repeat\_period: This 16 bit field gives the number of superframes between SYNC repeats, for example, SYNC\_repeat\_period=0 means that the SYNC slot is assigned on each superframe, SYNC\_repeat\_period=1 means that two superframes containing the SYNC slot assignment are separated by 1 superframe that does not have the SYNC slot assigned, and so on;
- SYNC\_slot\_number: This 11 bit field identifies the timeslot number to use for SYNC. This number matches the numbering defined in subclause 6.7.2.

### 8.5.5.10.7 Encrypted Logon ID Descriptor

This descriptor (table 36) is provided as one possible method of TBTP key management, when the TBTP is encrypted. It requires that the TIM message itself be encrypted with an RCST specific key.

**Table 36: Encrypted Logon ID Descriptor**

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
Encrypted_Logon_ID_descriptor() {			
descriptor_tag		8	uimbsf
descriptor_length		8	uimbsf
ID_start_time		16	uimbsf
ID_update_period		16	uimbsf
ID_loop_count		8	uimbsf
for (i = 0; i <= ID_loop_count; i++) {			
Encrypted_Logon_ID }		16	uimbsf
}			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. They shall be ignored by the RCST. The value of these bits is undefined and shall be varied in a random manner to avoid encryption spoofing.		

Semantics for the Encrypted\_Logon\_ID\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- ID\_Start\_time: This 16 bit field gives the superframe count value at which the first Encrypted\_Logon\_ID parameter in the following loop takes effect. This corresponds to the superframe counter defined in the SCT and TBTP;
- ID\_update\_period: This 16 bit field gives the period, in terms of superframe counts, at which successive Encrypted\_Logon\_ID parameters take effect following the first such parameter;
- ID\_loop\_count: This 8 bit field defines one less than the number of encrypted logon ID's that follow. A zero count indicates one loop. The ID's follow in sequential time order, earliest first;
- Encrypted\_Logon\_ID: This 16 bit value gives the value for a future encrypted logon ID.

### 8.5.5.10.8 Echo Value Descriptor

This descriptor (table 37) supports a simple loop-back RCST diagnostic test.

**Table 37: Echo Value Descriptor**

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
Echo_value_descriptor() {			
descriptor_tag		8	uimbsf
descriptor_length		8	uimbsf
Echo_Value		16	'1' + bslbf
}			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. They shall be ignored by the RCST. The value of these bits is undefined. For an encrypted uni-cast TIM, the bit values shall be varied in a random manner to avoid encryption spoofing.		



Semantics for the Echo\_value\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- Echo\_value: This 16 bit field defines the value to be echoed back (for example in the 2 M&C bytes of the return linkSAC field). The MSB shall be set to '1'.

#### 8.5.5.10.9 Linkage descriptor (private data)

This is an extension to the standard DVB Linkage descriptor (see EN 300 468 [4] section 6.2.11), occupying the private data bytes provision at the end of that descriptor. It allows the RCST to identify which interactive network services its population group. The Linkage descriptor shall be as EN 300 468 Table 39, with the changes as highlighted in bold font in table 38.

**Table 38: Linkage Descriptor - private data**

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
linkage_descriptor(){			
descriptor_tag		8	uimsbf
descriptor_length		8	uimsbf
transport_stream_id		16	uimsbf
original_network_id		16	uimsbf
service_id		16	uimsbf
linkage_type		8	uimsbf
<b>Interactive_Network_ID</b>		16	uimsbf
<b>Population_ID_loop_count</b>		8	uimsbf
<b>for (i=0; i&lt;=Population_ID_loop_count; i++) {</b>			
<b>population_ID_base</b>		16	uimsbf
<b>population_ID_mask</b>		16	uimsbf
<b>}</b>			
<b>for (i=0; i&lt;N; i++) {</b>			
private_data_byte		8	bslbf
<b>}</b>			
<b>}</b>			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. They shall be ignored by the RCST. The value of these bits is undefined. For an encrypted uni-cast TIM, the bit values shall be varied in a random manner to avoid encryption spoofing.		

Semantics for the private data part of the linkage\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is defined in EN 300 468 [4] and is recalled in table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- transport\_stream\_id: This is a 16-bit field which identifies the TS containing the information service indicated;
- original\_network\_id: This 16-bit field gives the label identifying the network\_id of the originating delivery system of the information service indicated;
- service\_id: This is a 16-bit field which uniquely identifies an information service within a TS. The service\_id is the same as the program\_number in the corresponding program\_map section;
- linkage\_type: This is an 8-bit field specifying the type of linkage. Its value is 0x81 for 'RCS FLS';
- Interactive\_Network\_ID: This 16 bit field gives the label identifying the network\_ID for the interactive network that services the population\_ID's following;

- Population\_ID\_loop\_count: This 8 bit field indicates one less than the number of population\_ID ranges in the following list;
- population\_ID\_base and population\_ID\_mask: These two 16 bit values, in combination, define a range of population\_ID's associated with this linkage descriptor / interactive network. The population\_ID is the identifier for the population group that the RCST belongs to. This is a value assigned by the network operator, and is a configuration parameter known to the RCST prior to forward link acquisition. The population\_ID\_base parameter defines the fixed bit pattern part of the population\_ID range, while the population\_ID\_mask parameter defines those bit positions of the population ID that are 'don't care'. A '1' value in a bit of the mask indicates that the corresponding bit of the RCST population\_ID may be '1' or '0'. A '0' value in a bit of the mask indicates that the corresponding bit of the RCST population\_ID shall match the value of that bit in the base parameter;
- private\_data\_byte: This is an 8 bit field, the value of which is privately defined. It retains the functionality of the original linkage\_descriptor for further extensions.

#### 8.5.5.10.10 RCS content descriptor

The RCS content descriptor provides the definition of the PID assignments to the RCS specific tables, and is shown in table 39. This descriptor is used in the second loop of the Program Map table (PMT), defined in ISO/IEC 13818-1 [9], subclause 2.4.4.8 / table 2-28. Each descriptor defines the General RCS SI Tables and RCS Specific Messages associated with one elementary\_PID.

NOTE: The PCR Insertion packet PID is signalled earlier in the PMT section as the PCR\_PID parameter.

**Table 39: RCS Content Descriptor**

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
RCS_content_descriptor() {			
descriptor_tag		8	uimsbf
descriptor_length		8	uimsbf
for (i=0; i<N; i++) {			
Table_id		8	uimsbf
}			
}			
NOTE:	*Reserved bits are of type bsbf, and shall precede the Information bits on the same line. They shall be ignored by the RCST. The value of these bits is undefined. For an encrypted uni-cast TIM, the bit values shall be varied in a random manner to avoid encryption spoofing.		

Semantics for the RCS\_content\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- table\_id: This 8 bit field gives the table\_id value for a General RCS SI Table or RCS Specific Message, as defined in table 10.

### 8.5.5.10.11 Satellite forward link descriptor

The satellite forward link descriptor defines the forward link, and is used in place of the Satellite delivery system descriptor of EN 300 468 [4] for RCS systems. It is shown in table 40.

**Table 40: Satellite Forward Link Descriptor**

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
Satellite_forward_link_descriptor() {			
descriptor_tag		8	uimsbf
descriptor_length		8	uimsbf
satellite_ID		8	uimsbf
beam_ID		16	uimsbf
NCC_ID		8	uimsbf
link_usage		3	bslbf
local_link_ID		5	uimsbf
frequency		32	uimsbf
orbital_position		16	bslbf
west_east_flag		1	bslbf
polarization		2	bslbf
modulation		5	bslbf
symbol_rate		24	uimsbf
FEC_inner		4	bslbf
reserved	4		bslbf
for (i=0; i<N; i++) {			
private_data_byte		8	bslbf
}			
}			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. They shall be ignored by the RCST. The value of these bits is undefined. For an encrypted uni-cast TIM, the bit values shall be varied in a random manner to avoid encryption spoofing.		

Semantics for the Satellite\_forward\_link\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- satellite\_ID: This 8 bit field identifies which satellite is carrying the forward link, and corresponds to the satellite\_ID field in the SPT / PCR Insertion packet payload;
- beam\_ID: This 16 bit field identifies the beam number of the satellite carrying the forward link;
- NCC\_ID: This 8 bit field identifies which NCC is transmitting the forward link, and corresponds to the NCC\_ID field in the PCR Insertion packet payload (subclause 8.5.5.6);
- link\_usage: This 3 bit field allows link discrimination for RCSTs that can operate on multiple forward links simultaneously. The usage codes are shown in table 41.

**Table 41: Forward Link usage codes**

Usage code	Value
000	Combined signalling / data link
001	Signalling link only
010	Data link only
111	Release data link
011 to 110	Reserved for future use

The first two codes are mutually exclusive in any Transport Stream ID (TS\_id) entry of the RMT (there can be only one signalling link for a given population\_ID). The signalling link only value can only be used where all RCSTs covered by that TS\_id are capable of receiving multiple forward links simultaneously.

The release data link value allows a previously assigned data link to be released, without logging off the terminal.

For RCSTs that can operate on multiple forward links simultaneously, there are two possible modes of defining links during logon:

- the RMT defines only the signalling link, and the TIM defines any data only links. This is the most flexible method;
- the RMT uses multiple descriptors, one for each forward link defined. Only one of these may be a signalling link.

NOTE 1: In DVB terms, this mode "stretches" the definition of a transport stream to actually cover a group of inter-related transport streams.

For both modes, the TIM can command changes to the initial configuration defined via the RMT.

NOTE 2: Changing a signalling link may cause service interruption and is not recommended.

RCSTs capable only of single carrier operation shall use only the combined signalling / data link descriptor and shall ignore descriptors for other link\_usage codes.

- local\_link\_ID: This 5 bit field is used to simplify changes to the definition of a link for RCSTs that can operate on multiple forward links simultaneously, and is a RCS local value defined by the interactive network operator. It allows the NCC to indicate which of the forward links is being created, changed or released. The last two operations can only be performed via a TIM message.
- frequency: This 32-bit field gives the frequency value. The frequency is given in multiples of 100 Hz.
- orbital\_position: The orbital\_position is a 16 bit field giving the 4-bit BCD values specifying 4 characters of the orbital position in degrees where the decimal point is after the third character (e.g. 019.2 degrees).
- west\_east\_flag: The west\_east\_flag is a 1 bit field indicating if the satellite position is in the western or eastern part of the orbit. A value '0' indicates the western position and a value '1' indicates the eastern position.
- polarization: The polarization is a 2 bit field specifying the polarization of the transmitted signal (see table 42).

**Table 42: Forward Link Polarization**

Polarization	Value
linear - horizontal	00
linear - vertical	01
circular - left	10
circular - right	11

- modulation: This is a 5-bit field specifying the modulation scheme used on the satellite forward link according to table 43.

**Table 43: Forward Link Modulation**

Modulation	Value
Not defined	0 0000
QPSK	0 0001
Reserved for future use.	0 0010 to 1 1111

- symbol\_rate: The symbol\_rate is a 24 bit field giving the symbol rate in multiples of 100 symbols/s.
- FEC\_inner: The FEC\_inner is a 4 bit field specifying the inner FEC scheme used, as per table 44.

Table 44: Inner FEC scheme

Code Rate	Value
$\frac{1}{2}$	0000
$\frac{2}{3}$	0001
$\frac{3}{4}$	0010
$\frac{5}{6}$	0011
$\frac{7}{8}$	0100
Inner code is omitted	1111
Reserved for future use	0101 to 1110

- `private_data_byte`: This is an 8 bit field, the value of which is privately defined. It can be used, for example, to indicate system specific NCC information.

#### 8.5.5.10.12 Satellite return link descriptor

Table 45: Satellite Return Link Descriptor

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
Satellite_return_link_descriptor() {			
descriptor_tag		8	uimsbf
descriptor_length		8	uimsbf
satellite_ID		8	uimsbf
beam_ID		16	uimsbf
gateway_ID		8	uimsbf
reserved	8		bslbf
orbital_position		16	bslbf
west_east_flag	7	1	bslbf
Superframe_ID		8	uimsbf
Tx_frequency_offset		24	tcimsbf
for (i=0; i<N; i++) {			
private_data_byte		8	bslbf
}			
}			
NOTE: *Reserved bits are of type bslbf, and shall precede the Information bits on the same line. They shall be ignored by the RCST. The value of these bits is undefined. For an encrypted uni-cast TIM, the bit values shall be varied in a random manner to avoid encryption spoofing.			

The satellite return link descriptor defines the characteristics of the return link and is shown in table 45.

Semantics for the `Satellite_return_link_descriptor`:

- `descriptor_tag`: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- `descriptor_length`: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the `descriptor_length` field;
- `satellite_ID`: This 8 bit field identifies which satellite is carrying the return link, and corresponds to the `satellite_ID` field in the SPT / PCR Insertion packet payload;
- `beam_ID`: This 16 bit field identifies the beam number of the satellite carrying the return link.;
- `gateway_ID`: This 8 bit field identifies which Gateway is receiving the return link, and corresponds to the `gateway_ID` field in the PCR Insertion packet payload;
- `orbital_position`: The `orbital_position` is a 16 bit field giving the 4-bit BCD values specifying 4 characters of the orbital position in degrees where the decimal point is after the third character (e.g. 019.2 degrees);
- `west_east_flag`: The `west_east_flag` is a 1 bit field indicating if the satellite position is in the western or eastern part of the orbit. A value '0' indicates the western position and a value '1' indicates the eastern position;

- Superframe\_ID: This 8 bit field identifies which superframe the terminal is to use. This identifier matches an entry in the SCT;
- Tx\_frequency\_offset: This 24-bit field gives the signed offset of the RCST transmit centre frequency relative to the Superframe\_centre\_frequency parameter (SCT). The frequency is given in multiples of 100 Hz. This parameter shall be ignored by the RCSTs when it is used in broadcast messages;
- private\_data\_byte: This is an 8 bit field, the value of which is privately defined. It can be used, for example, to indicate system specific NCC information.

#### 8.5.5.10.13 Table Update Descriptor

The Table Update descriptor (see table 46) provides a mechanism for notifying RCSTs when there is an upcoming change to one of the RCS Tables. In normal operation, these tables are expected to change infrequently. This descriptor provides a method for notifying the RCSTs on a particular interactive network when an update to one or more of these tables is imminent. This allows the RCSTs to only scan for table changes at this time, so freeing up processing resources. An RCST that monitors the composition tables continually may ignore this descriptor. Use of this descriptor does not modify the normal update control method implemented by the version\_number and current\_next\_indicator fields of the section header. In addition, the Table Update Descriptor may be utilized to indicate an update to any General SI Table defined in the present document or DVB-SI table, such as NIT, BAT, SDT, EIT, etc...

NOTE: In the case of the SCT, the Superframe start time (parameters Superframe\_start\_time\_base and Superframe\_start\_time\_ext) for each superframe definition, changes in a systematic fashion to reflect the regular period of the superframe. No update will be signalled where this is the only change and where the change has no material effect on superframe timing (i.e. the new start time equals a previous start time plus an integer number of superframe periods for the superframe in question).

**Table 46: Table Update Descriptor**

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
table_update_descriptor() {			
descriptor_tag		8	uimsbf
descriptor_length		8	uimsbf
interactive_network_id		16	uimsbf
for (i=0; i<n; i++) {			
table_id		8	uimsbf
new_version	3	5	uimsbf
}			
}			
NOTE: *Reserved bits are of type bslbf, and shall precede the Information bits on the same line. They shall be ignored by the RCST. The value of these bits is undefined. For an encrypted uni-cast TIM, the bit values shall be varied in a random manner to avoid encryption spoofing.			

Semantics for the table\_update\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- interactive\_network\_id: This 16 bit field indicates which interactive network the tables belong to, and matches the field of the same name in the table section header;
- table\_id: This 8 bit field indicates a change to the associated table. The table\_id values are defined in table 10 in the present document for the General SI Tables and in EN 300 468 [4] for the standard DVB-SI tables;
- new\_version: This 5 bit field defines the new version number of the table after the change, and matches the version\_number field in the table section header.

## 8.5.5.10.14 Contention control descriptor

Table 47: Contention Control Descriptor

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
Contention_Control_descriptor(){			
descriptor_tag		8	uimsbf
Descriptor_length		8	uimsbf
Superframe_ID		8	uimsbf
CSC_response_timeout		32	upcrmsf
CSC_max_losses		8	uimsbf
Max_time_before_retry		32	upcrmsf
}			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. For non-encrypted message types, these bits shall be set to '0'.		

The Contention control descriptor (see table 47) defines the necessary retransmission parameters for the CSC bursts.

Semantics for the contention\_control\_descriptor:

- descriptor\_tag: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- descriptor\_length: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the descriptor\_length field;
- Superframe\_ID: this is an 8-bit field which serves as a label for identification of the relevant superframe to which this descriptor applies;
- CSC\_response\_timeout: this 32-bit field gives the value of the timeout after which the RCST shall consider that there was a collision on its previous CSC burst. It is expressed in terms of PCR count intervals. The 32 bits correspond to a maximum duration of 93,2 seconds;
- CSC\_max\_losses: this 8-bit field specifies the maximum number of unsuccessful CSC attempts before the RCST gives up its logon procedure,
- Max\_time\_before\_retry: this 32-bit field gives the upper bound before retransmission of a CSC burst, expressed in terms of PCR count intervals, after the CSC\_response\_timeout has expired. In other words, upon expiry of the CSC\_response\_timeout and if the CSC\_max\_losses has not been reached, the RCST shall wait a random time between 0 second and this upper bound before transmitting a new CSC burst. The 32 bits correspond to a maximum duration of 93,2 seconds.

## 8.5.5.10.15 Correction Control descriptor

The Correction control descriptor (see table 48) defines the necessary timeouts for the coarse synchronization, fine synchronization and synchronization maintenance procedures.

Table 48: Correction Control Descriptor

Syntax	No. of bits		Information Mnemonic
	Reserved*	Information	
Correction_Control_descriptor(){			
descriptor_tag		8	uimsbf
Descriptor_length		8	uimsbf
ACQ_response_timeout		32	upcrmsf
SYNC_response_timeout		32	upcrmsf
ACQ_max_losses		8	uimsbf
SYNC_max_losses		8	uimsbf
}			
NOTE:	*Reserved bits are of type bslbf, and shall precede the Information bits on the same line. For non-encrypted message types, these bits shall be set to '0'.		

Semantics for the `correction_control_descriptor`:

- `descriptor_tag`: The descriptor tag is an 8 bit field which identifies each descriptor. Its value is given in the Tag value column of table 29;
- `descriptor_length`: The descriptor length is an 8 bit field specifying the number of bytes of the descriptor immediately following the `descriptor_length` field;
- `ACQ_response_timeout`: this 32-bit field gives the value of the timeout after which the RCST shall assume there was a problem on its previous ACQ burst. It corresponds to the "Correction received in time" test of the coarse sync procedure. It is expressed in terms of PCR count intervals. The 32 bits correspond to a maximum duration of 93.2 seconds;
- `SYNC_response_timeout`: this 32-bit field gives the value of the timeout after which the RCST shall assume there was a problem on its previous SYNC burst. It corresponds to the "Correction received in time" test of the fine sync and sync maintenance procedures. It is expressed in terms of PCR count intervals. The 32 bits correspond to a maximum duration of 93,2 seconds;
- `ACQ_max_losses`: this 8-bit field specifies the maximum number of consecutive ACQs with no correction received before expiry of `ACQ_response_timeout`. When this number is reached, the RCST shall return to the "Off / Stand-by" state. It corresponds to the "max losses exceeded" test of the coarse sync procedure (see subclause 7.4);
- `SYNC_max_losses`: this 8-bit field specifies the maximum number of consecutive SYNCs with no correction received before expiry of `SYNC_response_timeout`. When this number is reached, the RCST shall return to the "Off / Stand-by" state. It corresponds to the "max losses exceeded" test of the fine sync procedure and of the sync maintenance procedure (see subclauses 7.5 and 7.6).

#### 8.5.5.11 Accessing of the forward link signalling

Service Information for Forward Link Signalling (FLS) shall be transmitted in one or more Forward Link Signalling services. Being a DVB data service, the FLS service shall be referenced like any other DVB service in the Program Association Table (PAT), Program Map Table (PMT) and optionally Service Description Table (SDT). If contained in the SDT, the `service_type` shall be set to 0x0F for 'RCS FLS'. The PMT of the FLS service shall contain the PIDs for the NCR, the General SI tables and RCST specific messages as defined in the present document.

Additional information about the Satellite Interactive Network shall be conveyed in a RCS Map Table (RMT). This table having the same syntax as the NIT -but being transmitted on a different PID- shall contain one or more linkage descriptors pointing to one or more FLS services. Each linkage descriptor shall contain an `interactive_network_id` and a list of `population_ids`, which are being used by the RCST to select its appropriate FLS service. For this purpose the RCST has a default `interactive_network_id` and `population_id` for the installation process which may be changed at a later stage and be memorized for following network logons.

An interactive network is defined by a group of terminals serviced by a single NCC. In the simplest case, each NCC handles a single interactive network, however a large NCC logically splits its capacity into several `interactive_networks`. A population is a subset of a particular interactive network.

For the purposes of this specification, the forward link acquisition procedure is concerned with acquiring the forward link that carries the Forward Link Signalling service. This forward link acquisition procedure comprises the following steps:

- physical link synchronization;
- forward link location;
- NCR synchronization;
- DVB-RCS specific table loading.

*Physical link synchronization* is the process of tuning to the required transponder, recovering the transmitted signal and synchronizing to the DVB transport stream formatting on that signal. This shall be done each time an RCST is required to move to a different transponder.



*Forward link location* is the procedure for identifying which transponder carries the Forward Link Signalling service, and for identifying the PIDs. This procedure uses three standard DVB table types as follows:

- Network Information Table (NIT);
- Program Allocation Table (PAT);
- Program Map Table (PMT).

The detailed procedure for forward link location shall be as defined below in this section. It should be noted that this procedure may involve tuning to a maximum of three transponders in sequence, requiring a physical link synchronization for each transponder.

An outline of the forward link location procedure flow is shown in figure 35. This procedure requires that the following parameters be stored at the RCST as power-up configuration data:

- location details for forward link start-up Transport Stream (TS). This can be any Transport Stream in the network. The parameters correspond to those defined in the *satellite\_delivery\_system\_descriptor* of EN 300 468 [4].
- *population\_ID* value.

The RCST shall tune to the start-up Transport Stream and scan the copy of the NIT to locate the Transport Stream carrying the RMT. This shall be identified by locating the linkage descriptor containing the *linkage\_type* code (0x07) for the RCS Map service. It shall use the *TS\_id* parameter from that linkage descriptor to locate the *satellite\_delivery\_system\_descriptor* for that Transport Stream (in the second loop of the NIT), and hence the tuning details for the Transport Stream, which carries the RMT.

The RCST shall then re-tune (if necessary) to the Transport Stream carrier that carries the RMT and shall load the RMT using the assigned PID. It shall scan the RCS Map Table for all linkage\_descriptors containing the RCS FLS service *linkage\_type* code 0x81, to find the descriptor containing its *population\_id* (in the private data extension to that descriptor type). This descriptor will also contain the *Interactive\_network\_id* to be used by that RCST, together with the *TS\_id* for the Transport Stream to be used by the RCST for logon and the *service\_id* to be used by the PAT later.

The RCST shall then locate the entry for that *TS\_id* in the second loop of the RCS Map Table and extract the satellite forward link and satellite return link descriptors. Where there are multiple instances of these, then it shall extract the signalling satellite forward link and the primary satellite return link as a minimum, and use these for the balance of the procedure. These descriptors contain initial parameters required by the subsequent Return Link Acquisition process, namely the *satellite\_id* for both forward and return links, the *gateway\_id* for the return link, the *NCC\_id* for forward link, plus the *superframe\_id* and transmit centre frequency offset for the return link. The RCST shall then re-tune (if necessary) to the Transport Stream carrier carrying its forward link signalling, using the parameters defined in the signalling satellite forward link descriptor, and locate the PAT entry with *prg\_nbr* parameter equal to the *service\_id*. This defines the PID for the associated PMT.

From the PMT, it shall extract the PID for the PCR Insertion packet (*PCR\_PID*), and the PIDs for the RCS specific signalling tables. These PIDs are identified by locating those stream definitions containing the RCST content descriptor (defined in subclause 8.5.5.10.10), and associating the *table\_id*'s for RCS tables with the *elementary\_PID* value defined for that stream.

The result of this forward link location procedure shall be the following:

- identification of the NCC and satellite for the Forward Link Signalling service;
- identification of the Gateway and satellite for the return link;
- definition of the PIDs for the DVB-RCS forward signalling messages.;
- definition of the superframe to be used for the return link logon procedure.

*NCR synchronization* is the procedure for locking the RCST timing to the NCR count at the NCC. This shall be done by receiving the PCR Insertion Transport Stream Packet (subclause 8.5.5.6) at the transmitted interval and deriving a synchronous PCR count and PCR clock locally in the RCST that are locked to the PCR count in that packet. Upon successful acquisition of the receive NCR, the RCST shall continue to monitor it. As a matter of safety if the reception of the forward link is interrupted, the RCST shall cease transmission if conditions defined in subclause 7.7.3 are fulfilled.

*DVB-RCS specific table loading.* The RCST shall load the following tables, using the PIDs identified in the forward link location procedure:

- SCT, FCT, and TCT to identify the available superframe and frame structures;
- SPT to obtain the satellite ephemeris data for the forward link signalling satellite and return link satellite;
- TIM (broadcast version) to identify possible link anomalies.

In addition, the RCST shall monitor for unicast TIM messages addressed to the RCST MAC address for possible wake-up messages.

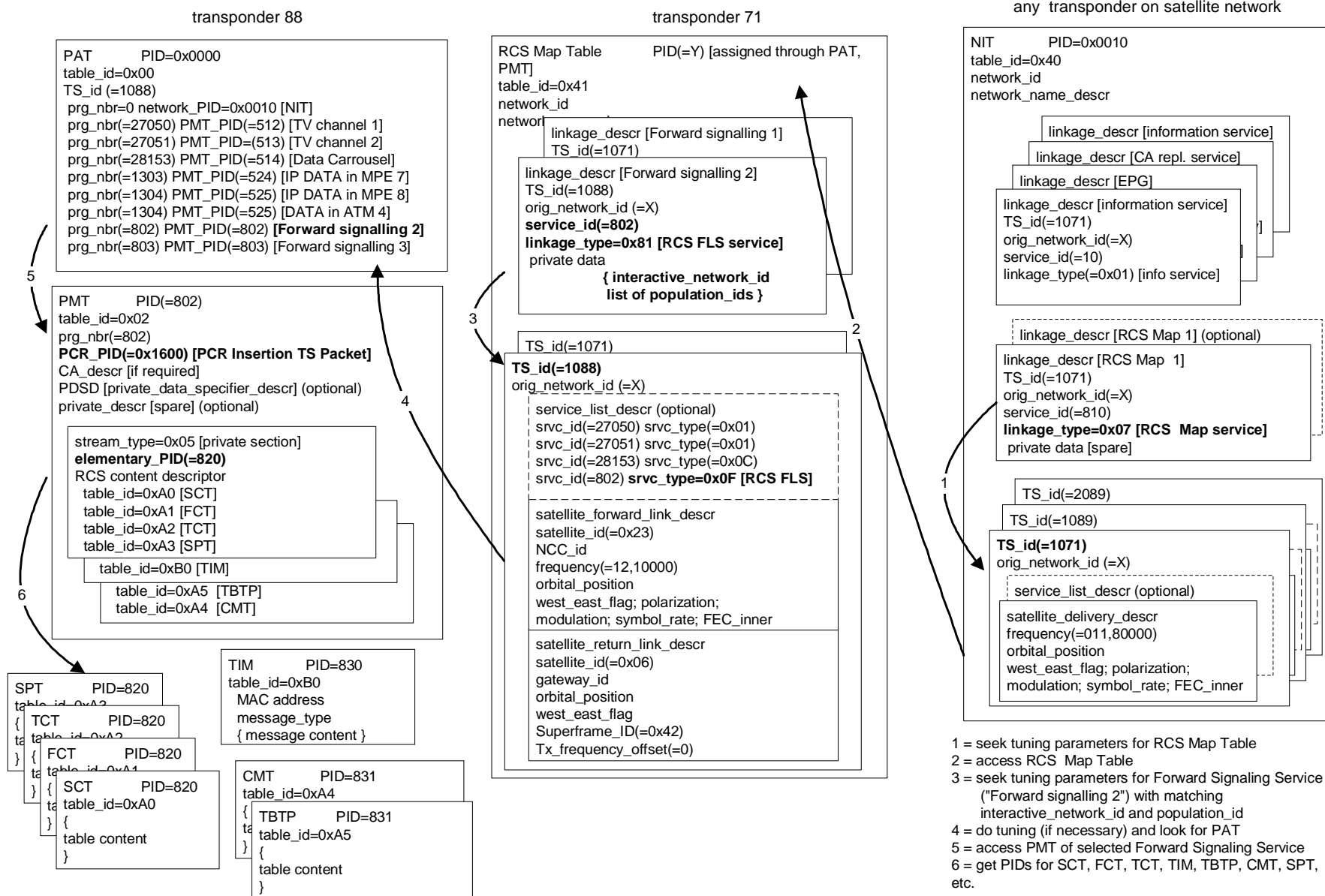


Figure 35: Example of forward Link Location

## 9 Security, Identity, Encryption

Security is intended to protect the user identity including its exact location, the signalling traffic to and from the user, the data traffic to and from the user and the operator/user against use of the network without appropriate authority and subscription. Three levels of security can be applied to the different layers:

- DVB common scrambling in the forward link (could be required by the service provider);
- satellite interactive network individual user scrambling in the forward and return link;
- IP or higher layer security mechanisms (could be used by the service provider, the content provider).

Although the user/service provider could use its own security systems above the data link layer, it may be desirable to provide a security system at the data link layer so that the system is inherently secure on the satellite section without recourse to additional measures. Also, since the satellite interactive network forward link is based on the DVB/MPEG-TS Standard, the DVB common scrambling mechanism could be applied, but is not necessary (it would just add an additional protection to the entire control stream for non-subscribers). This concept is shown in figure 36.

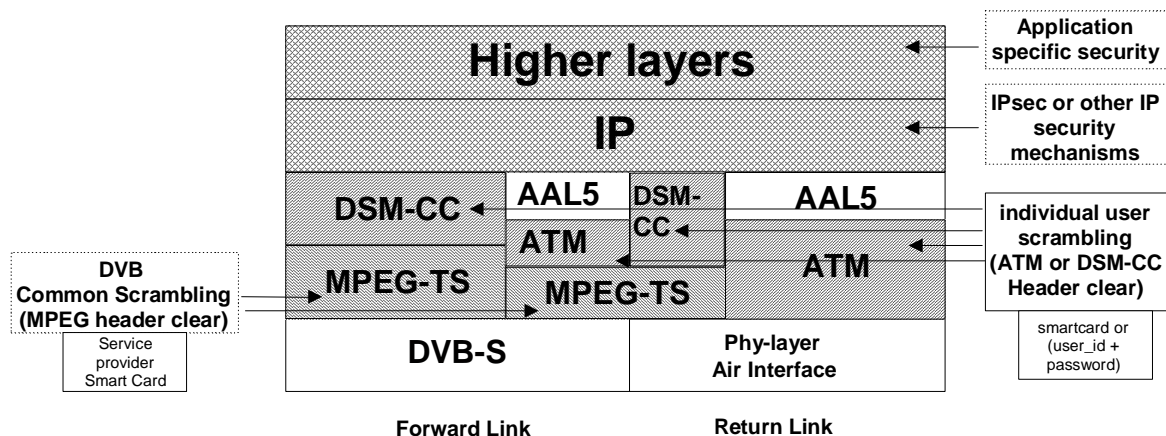


Figure 36: Security layers for satellite interactive network (example)

In the following it is assumed there can be more than one user per RCST and that such users will have security in their own right. Security is thus defined at a level higher than the individual RCST. On a user basis, an authentication algorithm may either check for user name and password on the client device or may use a Smart Card within the RCST. All data and control to and from each user may be scrambled on an individual user basis. Each user may have a control word for the return and the forward link that does not allow anybody other than the NCC/Gateway or the user himself to descramble the data, except for lawful interceptors such as country authorities. An optional security mechanism derived from the one used in [14] is also considered.

### 9.1 Authentication

Authentication may be implemented by request for a user name or password on the client device. In the case of a PC used as the client device, the RCST does not need to carry any special implementation. However, if the RCST contains a proxy client, then the proxy may be able to authenticate itself to the NCC. This means that an authentication server may be implemented at the NCC, which manages the authentication of each user.

Authentication could also be replaced by a Smart Card on the RCST, also used for the link layer individual control word encryption.

### 9.2 Forward Link

DVB Common Scrambling could be required in the Forward Link.

Individual scrambling may be implemented at the section level, but the MAC address of the user may remain in the clear, since the RCST uses the MAC address to filter messages.

## 9.3 Return Link

The client device may handle IPsec, so the router at the Gateway may be able to handle IPsec.

Individual layer 2 scrambling may also be implemented.

## 9.4 Security (optional)

The optional security mechanism described here is derived from the one used in [14]. The RCSTs indicates in the CSC bursts (see subclause 6.2.3) if it implements this option. The NCC indicates in the logon initialize descriptor (see subclause 8.5.5.10.4) if this mechanism is to be used. This security mechanism consists of two separate sub-systems:

- a set of MAC messages used for authentication and key-agreement between NCC and RCST. These messages are used for key negotiation during a session setup as well as for on-the-fly update of keys (see subclause 9.4.7);
- on-the-fly encryption and decryption of payload data streams passed between NCC and RCST.

When a session is being setup, before payload data is transferred, one of three request/response MAC message-pairs is used to generate a session key specific to the payload stream associated with the session.

The session key is a shared secret between the NCC and the RCST: even if every MAC message is intercepted, the cryptographic properties of the protocol ensure that an eavesdropper cannot determine the session key value.

This is achieved by using a public-key protocol, which requires no up-front shared secret, or a simpler protocol based on a long-term shared secret between NCC and RCST called a cookie. The cookie is 160 bits long. It can be used for authenticating the RCST to the NCC at logon.

Each RCST will store its own cookie in non-volatile storage, whereas the NCC will maintain a data-base of the cookie values of the RCSTs on its network. Cookie values will be updated occasionally as dictated by security policy, but they are less vulnerable than session keys: a successful brute-force attack on a session key reveals nothing about the cookie value, nor any other session key.

The MAC messages also implement a defence against clones: an RCST that is a physical copy of an existing RCST and attempts to operate on the network under the cloned identity (when the cloned RCST itself is not registered on the network). The anti-cloning measure is a simple non-volatile 8-bit counter that is incremented synchronously at the NCC and RCST over time: if a clone RCST engages in traffic with the NCC, this will be detected the next time the cloned RCST connects because the counter value will be out of synchronization.

If the clone attempts to operate concurrently with the cloned unit, there will be an immediate break-down of functionality for both units, due to confusion within the MAC protocol. This amounts to a denial-of-service attack, and the NCC should be prepared for this kind of protocol failure.

### 9.4.1 Cryptographic primitives

The key exchange protocols and data stream encryption is based on a set of well-established primitive cryptographic functions. The functions and their associated key sizes can be changed in the future, in case crypt-analytic or brute-force attacks become a realistic threat.

The specific set of functions and key sizes are negotiated between NCC and RCST at sign-on time. The functions supported at the present time are Diffie-Hellman, HMAC-SHA1, and DES. Check current cryptographic literature for any updates regarding their security and use.

The following sections give a brief overview of the cryptographic primitives, and details on how they are used in the protocol. Later sections describe the exact field layout of the MAC messages.

The protocol parameters are described in terms of byte strings. Integer quantities are represented as base-256 byte strings. Big-endian byte-ordering is used, that is, the most significant byte comes first. If necessary to reach a fixed length, the string is padded with zeros at the most significant end.

### 9.4.1.1 Public key exchange

A public key exchange primitive is used to allow the NCC and RCST to agree on a secret, although communicating in public. The Diffie-Hellman scheme is based on unsigned integer arithmetic and works as follows.

The NCC chooses two public values, a large prime number  $m$ , and a (small) number  $g$  which is a generator modulo  $m$  (that is,  $g^a \bmod m$  will generate all number from 0 to  $m-1$  for varying  $a$ ). The NCC also chooses a secret number  $x < m$ , and sends the following three values to the RCST:  $m$ ,  $g$ ,  $X = g^x \bmod m$ .

The RCST chooses a secret value  $y < m$ , and responds to the NCC with the value  $Y = g^y \bmod m$ .

The RCST now calculates  $s = X^y \bmod m = (g^x)^y \bmod m = g^{(x*y)} \bmod m$ , whereas the NCC calculates  $Y^x \bmod m = (g^y)^x \bmod m = g^{(y*x)} = s$ , so the NCC and RCST now agree on the value  $s$ .

The value of  $s$  is a secret shared between NCC and RCST. To determine its value from the publicly communicated values  $m$ ,  $g$ ,  $X$ , and  $Y$ , an eavesdropper shall determine  $x$  or  $y$  by solving an equation of the form  $Z = g^z \bmod m$  for unknown  $z$ . This is known as the discreet logarithm problem and is computationally infeasible with current algorithms for sufficiently large values of  $m$ .

The parameter size supported are 512 bits for the prime number  $m$ , and hence also for the remaining values since all arithmetic is modulo  $m$ .

In the applicable MAC messages, the unsigned integer quantities  $m$ ,  $g$ ,  $X$ , and  $Y$  are encoded into fixed-size fields (64, 96, or 128 bytes) using big-endian byte-ordering.

### 9.4.1.2 Hashing

The protocol makes use of a keyed hash function that computes secure checksums which can only be verified with the possession of a secret key. The function has the one-way property, meaning that it is computationally infeasible to find an input value that maps to a given output value.

The hash function is also used to generate derived secret material based on a master secret. Because of the one-way property, the master secret is protected even if the derived secret is discovered.

In generic terms, the keyed hash function takes two byte strings as input, the key and a data string, and produces another string of bytes, the digest:

- $\text{digest} = H(\text{key}, \text{data})$ .

The  $H$  function shall accept key and data parameters of any size, whereas the protocol is designed to accept digests of any size.

The specification currently supports the HMAC-SHA1 function defined in RFC 2104 [15]. It produces a 20-byte digest.

### 9.4.1.3 Encryption

Payload data is encrypted and decrypted using a symmetric-key block cipher, which is used in Cipher Block Chaining (CBC) mode with special handling of any final odd-size block.

In generic terms, the encryption and decryption functions take two byte strings as input, the key and a data block, and produce as output another data block of the same length:

- $\text{ciphertext} = E(\text{key}, \text{plaintext})$ ;
- $D(\text{key}, \text{ciphertext}) = \text{plaintext}$ .

The key length and block length is given by the chosen cipher, and the payload stream processing logic will apply it as appropriate to data units of various sizes.

The specification currently supports the DES algorithm, which has a block size of 8 bytes, and various options for key length based on an 8-byte raw key block (see subclause 9.4.5).

#### 9.4.1.4 Pseudo-random numbers

The protocols used for generating secret values depend on the availability of a pseudo-random, that is, practically unpredictable, endless string of bytes. This will typically be produced with a Pseudo-Random Number Generator, PRNG, algorithm.

The random bytes are used to generate the secret Diffie-Hellman values,  $x$  and  $y$ , and for nonce values used during key exchange. The unpredictable nature of the random input ensures that different secret values are produced each time, and also prevents replay of old intercepted messages.

This specification does not require any particular algorithm, only that the NCC and RCST each choose one that is well-established and cryptographically analysed.

The hardest aspect of using a PRNG is to initialize it with an unpredictable seed value. The seed should contain multiple high-granularity device-dependent time-samples, as well as any other available pseudo-random material, like file allocation tables, etc. These random source values are then hashed together to squeeze out the entropy for the seed value.

#### 9.4.2 Main Key Exchange, MKE

Main Key Exchange (MKE) uses Diffie-Hellman to develop a shared secret between the NCC and RCST, which is independent of the cookie value. Furthermore, it uses the cookie value to authenticate the RCST to the NCC. It optionally uses the newly developed shared secret to update the cookie value. Finally, it derives a shared secret key used for the security context that is used to process payload stream data.

The exchange is initiated by the NCC sending a message containing the Diffie-Hellman values,  $m$ ,  $g$ ,  $X$ , and a random nonce string,  $nonce1$ . The RCST responds with a message containing its Diffie-Hellman value,  $Y$ , a random nonce string,  $nonce2$ , and an authentication string,  $auth$ .

The NCC and RCST each use the same formula to calculate the authentication string:

- $auth = H(\text{cookie}, nonce1 \sim nonce2)$ ;

which is communicated by the RCST and checked by the NCC. This proves the identity of the RCST, since it requires knowledge of the cookie to calculate the correct value of  $auth$ .

The RCST and NCC each use the Diffie-Hellman values (see subclause 9.4.1.1) to arrive at the same secret value,  $s$ :

- $s = g^{(x*y)} \text{ mod } m$ .

This unsigned integer value is encoded as a byte string, of length specified by the Diffie-Hellman parameter size, using big-endian byte ordering. It is then used to calculate a temporary shared secret string,  $temp$ :

- $temp = H(\text{encode}(s), nonce2 \sim nonce1)$ .

If the cookie is to be updated, the new value is computed in sections for  $n = 1, 2, \dots$ :

- $newcookie(n) = H(temp \sim (\text{unsigned char})1 \sim (\text{unsigned char})n, "")$ .

These string values are computed and concatenated until the total length matches or exceeds the length of the cookie. The cookie is then obtained by taking the first 20 bytes out of the concatenated sections, starting from the beginning.

The session key used for payload stream encryption is likewise computed in sections:

- $key(n) = H(temp \sim (\text{unsigned char})2 \sim (\text{unsigned char})n, "")$ .

Where, again, a sufficient number of sections are calculated to produce enough bytes to cover the length of the key. The session key is obtained "in the same manner as the cookie" by taking the required number of bytes out of the concatenated sections, starting from the beginning.

### 9.4.3 Quick Key Exchange, QKE

Quick Key Exchange uses the existing cookie value to authenticate the RCST to the NCC, and then derive a shared secret key used for the security context that is used to process payload stream data.

The exchange is initiated by the NCC sending a message containing a random nonce string, *nonce1*. The RCST responds with a message containing a random nonce string, *nonce2*, and an authentication value, *auth*.

The value of *auth* is calculated in the same way as for Main Key Exchange, and can be used to verify the identity of the RCST (see subclause 9.4.2).

The RCST and NCC then each calculate a temporary shared secret string, *temp*:

- $temp = H(\text{cookie} \sim (\text{unsigned char})^3, \text{nonce2} \sim \text{nonce1})$ .

This value is used to produce the payload encryption key in the same way as for Main Key Exchange (see subclause 9.4.2).

### 9.4.4 Explicit Key Exchange, EKE

Explicit Key Exchange is used by the NCC to deliver a pre-determined session key to the RCST. The session key is encrypted under a temporary key derived from the cookie value, and is used for the security context that is used to process payload stream data.

The delivery is performed by the NCC sending a message containing a random nonce string, *nonce1*, and a byte string value, *encryptedkey*, which has the same length as a key used for payload encryption. The RCST responds with a message containing a random nonce string, *nonce2*, and an authentication value, *auth*.

The value of *auth* is calculated in the same way as for Main Key Exchange, and can be used to verify the identity of the RCST (see subclause 9.4.2).

Both the NCC and RCST calculate a temporary shared secret string, *temp*:

- $temp = H(\text{cookie} \sim (\text{unsigned char})^4, \text{nonce1})$ .

which is used to produce sections of a temporary key, in the same way as for Main Key Exchange (see subclause 9.4.2). The NCC uses these temporary key string sections to XOR with the session key to obtain the *encryptedkey* value, and the RCST performs a second XOR operation to decrypt the session key value.

For normal DES, 8 bytes of raw key data are delivered, which are used to derive the actual key with the appropriate number of effective bits, as described below (see subclause 9.4.2).

### 9.4.5 Key derivation

The actual key value used for processing payload data is derived from the key sections developed during key exchange. For DES, 8 bytes of raw key data is required, so a single 20-byte section, *key(1)*, computed by HMAC-SHA1 is sufficient.

In each byte, the least significant bit is not used (it can be used as an odd-parity bit of the remaining 7 bits), bringing the effective key size down to 56 bits.

Furthermore, when used in 40-bit mode, the two most significant bits of each byte in the key are zeroed.

### 9.4.6 Data stream processing

Security can be applied to various payload data streams selectively. The elementary unit is called a security context, which contains two session keys used for encrypting and decrypting a stream of payload data. Only one of the keys is used to process any particular payload unit. Each key can be used for processing both upstream and downstream payload data.



Having two keys allows negotiation of a new key to take place while payload data is processed using the old one, and then do an immediate switch-over once the new key is agreed upon, without interrupting payload traffic. The NCC initiates the key exchanges, and can start using a session key for downstream traffic encryption once the key exchange is complete. For upstream traffic encryption, the RCST should use whichever key was used by the NCC in the most recent encrypted payload unit.

#### 9.4.6.1 Payload streams

A payload stream is identified by either of:

- a 24-bit (UNI) ATM virtual circuit VPI/VCI;
- a 48-bit MAC-address: this is used for DVB Multiprotocol Encapsulation downstream payload data.

When a payload stream is secured, the RCST and the NCC will have matching security contexts, which are used to encrypt/decrypt both upstream and downstream traffic. For unsecured payload streams there is no security context, and payload data is not encrypted.

To support encrypted multi-cast traffic, the same security context will be created for each member using EKE (see subclause 9.4.4), so that each RCST can decrypt the common payload data stream.

#### 9.4.6.2 Data encryption

Within a payload data stream, data is carried in individual units at the various protocol layers. Encryption is applied at the lowest layer possible, consistent with the payload stream:

- ATM-based payload streams: the unit of encryption is a single ATM cell. The 48-byte cell payload is encrypted using the security context implied by the session. Encryption is transparent to higher-level protocol layers, which see only unencrypted cell payloads;
- DVB Multiprotocol Encapsulation payload streams: the unit of encryption is a single DVB Multiprotocol Encapsulation section. The datagram\_data\_bytes (between the MAC-address and the CRC/checksum) are encrypted using the security context implied by the session. The DVB Multiprotocol Encapsulation payload to be encrypted will be adjusted to have a length of  $n * 8$  bytes ( $n$  is an integer) by adding an appropriate amount (0 .. 7 bytes) of stuffing bytes before the CRC/checksum according to Annex B of EN 300 468 [4] after processing the entire section. The CRC/checksum is calculated on the encrypted datagram bytes, while higher-level protocol layers see only unencrypted datagrams.

#### 9.4.6.3 Encryption flags

There are flags in the header of each encryption unit specifying which of the two sessions keys of the security context is used.

The receiver will use the security context of the payload stream to see if decryption shall be done.

- ATM cells: the least significant two bits of the Generic Flow Control, GFC, field of the cell header are used:
  - 00: not encrypted;
  - 01: reserved;
  - 10: encrypted using session key 0;
  - 11: encrypted using session key 1.

The most significant two bits of the GFC field are reserved for future use, and shall be set to 00.

- DVB Multiprotocol Encapsulation sections, according to EN 301 192 [5], the 2-bit payload\_scrambling\_control field in the section header is used:
  - 00: not encrypted;
  - 01: reserved;

- 10: encrypted using session key 0;
- 11: encrypted using session key 1.

The 2-bit address\_scrambling\_control field in the section header is 00 all the time (the address is not scrambled).

#### 9.4.6.4 Chaining and initialization vector

Within encryption units, the block encryption algorithm is used in Cipher Block Chaining mode, CBC: the first plain-text block is XOR'ed with an Initialization Vector (IV), and subsequent blocks are XOR'ed with the previous cipher-text block, before the block is encrypted. Decryption is opposite: each cipher-text block is first decrypted and then XOR'ed with the previous chaining value.

The value of the IV for a given encryption unit is zero.

### 9.4.7 Security Establishment

Security issues are handled in the following situations:

- when an RCST registers on the network it will first go through the normal logon procedure as described in clause 7. It can then do an initial handshake with the NCC to establish the level of security support, in particular the cryptographic algorithms and key sizes to be used subsequently. During logon, the RCST indicates its encryption/security capability to the NCC on the CSC burst (see subclause 6.2.3) and the NCC communicates on the logon initialize descriptor (see subclause 8.5.5.10.4) whether security shall be used;

the handshake consists of **<MAC>Security Sign-On** and **<MAC>Security Sign-On Response** messages (see subclauses 9.4.9.1 and 9.4.9.2);

- The MAC security messages transmitted over the return path shall be transported using DULM (see subclause 6.6.2).

The MAC security messages transmitted over the forward path will be sent on a dedicated and well-known PID. "Security enhanced" RCSTs will thus have to MAC-filter this PID.

These remarks do apply to all the security related MAC messages.

A failure during this stage of the protocol causes the NCC to revert to non-secure interaction with the RCST.

- one session key is agreed for a logon session (i.e. it is valid between RCST logon and logoff), and the cookie and/or clone counter values may be updated as part of the exchange;

The key exchange consists of **<MAC>Main/Quick/Explicit Key Exchange** and **<MAC>Main/Quick/Explicit Key Exchange Response** messages (see subclauses 9.4.9.3 to 9.4.9.8) which are exchanged after the initial "security sign-on" handshake.

A failure during this stage of the protocol causes the RCST to be logged off.

- after an RCST is logged on, each session key of the security context of the payload stream can be updated on-the-fly, that is, without new logon, and without interrupting payload data traffic. The cookie and/or clone counter values cannot be updated as part of the exchange.

A new session key is negotiated using the same MAC messages used during logon.

A failure during this stage of the protocol causes the RCST to be logged off.

While a session key of the security context is being updated, payload stream data traffic should be encrypted using the other session key or not at all. Once the key exchange is complete, the NCC can start using it for subsequent downstream traffic, thereby directing the RCST to use it for upstream traffic.

All three variants of key exchange messages can authenticate the RCST based on the existing cookie value. They also perform the clone detection counter check, and optionally increment the clone counter. Only MKE can update the cookie.

## 9.4.8 Persistent state variables

To facilitate authentication, key exchange, and clone detection, the RCST has a set of state variables whose values are retained across registrations and power cycles:

**Table 49: Persistent RCST variables**

Name	Function	Size
Cookie	authentication cookie	160 bits
Cookie_SN	cookie sequence number	1 bit
Clone_Counter	clone detection counter	8 bits
Clone_Counter_SN	clone counter sequence number	1 bit

The sequence numbers are used to ensure that the NCC and RCST can stay synchronized even in case the RCST drops off the net in the middle of a protocol exchange.

### 9.4.8.1 Guaranteed delivery

The NCC will ensure that a protocol exchange is complete before proceeding. If it doesn't receive a response MAC message within a given time-interval, it will re-transmit the original message unchanged. The RCST will do likewise in situations where it requires a response. If the number of re-transmissions exceeds three, the protocol fails.

Due to race conditions, superfluous re-transmissions may be generated by both NCC and RCST. They shall discard such messages after the first message has in fact been received.

If the RCST is not ready to respond within the specified time-out, it can send <MAC>Wait messages (see subclause 9.4.9.9) to extend the time it has available to generate a proper response. Upon receiving the wait message, the NCC will restart its timer and reset the retry count.

The protocol time-out values can be set by the <MAC> **Default Configuration Message** as defined in subclause 5.5.3.2 of [14], otherwise the following default values apply.

**Table 50: Protocol time-out values**

Code	Protocol stage	Default Value
0xD	Security Sign-On	700
0xE	Main Key Exchange	1200
0xF	Quick Key Exchange Explicit Key Exchange	900

The Unit for the timeouts is msec.

## 9.4.9 Security MAC Messages

### 9.4.9.1 <MAC>Security Sign-On

As part of the registration process when an RCST performs the logon, the NCC and RCST will negotiate the specific set of cryptographic algorithms and parameters used in the key exchange protocols and for payload encryption.

The selections are global, and apply to all subsequent security exchanges for as long as the RCST is registered on the network.

The selections affect the layout of the subsequent key exchange messages, since they have fields that vary in size according to the choice of algorithms and parameters.

The NCC indicates which algorithms and parameters it supports by setting the appropriate bits in the <MAC>**Security Sign-On** message. There are four classes of algorithms, and the NCC will set one or more bits in each of the four fields to indicate which specific choices it supports.

Table 51: Security Sign-On message structure

Security_Sign-On () {	Bits	Bytes	Bit Number / Description	Parameter bytes
<b>Public_Key_Alg</b> PKA_Reserved PKA_DH_512	7 1	<b>1</b>	<b>Public key algorithm choices:</b> 7..1: Reserved, shall be 0 0:(yes/no) Diffie-Hellman, 512 bits	P <sub>pka</sub> : 64
<b>Hash_Alg</b> HA_Reserved HA_HMACSHA1	7 1	<b>1</b>	<b>Hash algorithm choices:</b> 7..1: Reserved, shall be 0 0:(yes/no) HMAC-SHA1	P <sub>ha</sub> : 20
<b>Encryption_Alg</b> EA_Reserved EA_DES_56 EA_DES_40	6 1 1	<b>1</b>	<b>Encryption algorithm choices:</b> 7..2: Reserved, shall be 0 1:(yes/no) DES, 56 bit key 0:(yes/no) DES, 40 bit key	P <sub>ea</sub> : 8
<b>Nonce_Size</b> NS_Reserved NS_64 NS_64	7 1	<b>1</b>	<b>Nonce size choices:</b> 7..1: Reserved, shall be 0 0: (yes/no) 8 random bytes	P <sub>ns</sub> : 8
<b>Reserved</b>	<b>32</b>	<b>4</b>	<b>Reserved for future use, shall be 0</b>	
}				

If the security option is supported, the minimum subset to support is PKA\_DH\_512, HA\_HMACSHA1, EA\_DES\_40, and NS\_64. EA\_DES\_56 is optional.

#### 9.4.9.2 <MAC>Security Sign-On Response

In its security sign-on response, the RCST indicates which specific algorithms and parameters to use. It does so by choosing one of the suggestions offered by the NCC within each of the four classes.

The fields of the response message have the same definition as the message from the NCC, except that exactly one bit will be set in each field.

If the RCST is unable to support any of the suggested algorithms for any class, it shall return an all-zero field value, and the NCC will revert to non-secure communication or re-issue the <MAC>Security Sign-On message with different choices.

Table 52: Security Sign-On Response message structure

Security_Sign-On_Response() {	Bits	Bytes	Bit Number / Description	Parameter bytes
<b>Public_Key_Alg</b> PKA_Reserved PKA_DH_512	7 1	<b>1</b>	<b>Public key algorithm choices:</b> 7..1: Reserved, shall be 0 0:(yes/no) Diffie-Hellman, 512 bits	P <sub>pka</sub> : 64
<b>Hash_Alg</b> HA_Reserved HA_HMACSHA1	7 1	<b>1</b>	<b>Hash algorithm choices:</b> 7..1: Reserved, shall be 0 0:(yes/no) HMAC-SHA1	P <sub>ha</sub> : 20
<b>Encryption_Alg</b> EA_Reserved EA_DES_56 EA_DES_40	6 1 1	<b>1</b>	<b>Encryption algorithm choices:</b> 7..2: Reserved, shall be 0 1:(yes/no) DES, 56 bit key 0:(yes/no) DES, 40 bit key	P <sub>ea</sub> : 8
<b>Nonce_Size</b> NS_Reserved NS_64	7 1	<b>1</b>	<b>Nonce size choices:</b> 7..1: Reserved, shall be 0 0: (yes/no) 8 random bytes	P <sub>ns</sub> : 8
<b>Reserved</b>	<b>32</b>	<b>4</b>	<b>Reserved for future use, shall be 0</b>	
}				

### 9.4.9.3 <MAC>Main Key Exchange

The Main Key Exchange message is used to start a cookie-independent key exchange with the RCST, and also instructs the RCST whether to update its cookie value and clone counter value. The `connection_id` is not used and shall be set to all '0'.

**Table 53: Main Key Exchange message structure**

Main_Key_Exchange () {	Bits	Bytes	Bit Number / Description
<b>Connection_ID</b>	<b>32</b>	<b>4</b>	<b>MAC connection identifier</b>
<b>Flags</b>		<b>1</b>	
Reserved	4		7..4: shall be 0
FL_Initializing	1		3:(yes/no) first ever key exchange
FL_Update_Cookie	1		2:(yes/no) make new cookie value
FL_Update_Counter	1		1:(yes/no) increment clone counter
FL_Session_Key	1		0: select session key 0 or 1
<b>Reserved</b>	<b>8</b>	<b>1</b>	<b>Reserved for future use, shall be 0</b>
<b>Nonce</b>		<b>P<sub>ns</sub></b>	<b>Random string nonce1</b>
<b>DH_Modulus</b>		<b>P<sub>pka</sub></b>	<b>Diffie-Hellman modulus m</b>
<b>DH_Generator</b>		<b>P<sub>pka</sub></b>	<b>Diffie-Hellman generator g</b>
<b>DH_Public_X</b>		<b>P<sub>pka</sub></b>	<b>Diffie-Hellman public value X</b>
}			

The `FL_Session_Key` bit specifies which session key of the security context to update.

If the `FL_Update_Counter` bit is set, it instructs the RCST to increment its clone detection counter.

If the `FL_Update_Cookie` bit is set, it instructs the RCST to generate a new cookie value to be used for future authentications and key exchanges, and to reset the clone detection counter to zero.

If the `FL_Initializing` bit is set, it tells the RCST that the `Authenticator` field in the response will be ignored.

The sizes of the multi-byte fields are determined by the parameters of the algorithms selected during security sign-on (see subclause 9.4.9.1).

The NCC will use its own private Diffie-Hellman value, `x`, together with the fields of the response message from the RCST to derive the new session key value, as well as any new value for the cookie (see subclause 9.4.2).

### 9.4.9.4 <MAC>Main Key Exchange Response

The Main Key Exchange Response message authenticates the RCST and completes the cookie-independent key exchange with the NCC. It also contains the current value of the clone detection counter. The `connection_id` is not used and shall be set to all '0'.

**Table 54: Main Key Exchange Response message structure**

Main_Key_Exchange_Response () {	Bits	Bytes	Bit Number / Description
<b>Connection_ID</b>	<b>32</b>	<b>4</b>	<b>MAC connection identifier</b>
<b>Flags</b>		<b>1</b>	
Reserved FL_Cookie_SN	6		7..2: shall be 0
FL_cookie_SN	1		1: cookie sequence number
FL_Counter_SN	1		0: clone counter sequence number
<b>Clone_Counter</b>	<b>8</b>	<b>1</b>	<b>Current clone counter value</b>
<b>Nonce</b>		<b>P<sub>ns</sub></b>	<b>Random string nonce2</b>
<b>Authenticator</b>		<b>P<sub>ha</sub></b>	<b>Authentication value auth</b>
<b>DH_Public_Y</b>		<b>P<sub>pka</sub></b>	<b>Diffie-Hellman public value Y</b>
}			

The `FL_Counter_SN` bit is the current sequence number of the clone detection counter. The `Clone_Counter` field is the current value of the counter. A clone collision has been detected if the NCC finds a mismatch from the expected value.

The `FL_Cookie_SN` bit is the sequence number of the cookie used for authentication.

If the `FL_Update_Cookie` bit was set by the NCC, the RCST will generate a new cookie value and complement the cookie sequence number bit. It will also reset the clone counter value to zero and clear the clone counter sequence number bit.

If the `FL_Update_Counter` bit was set by the NCC, the RCST will increment the value of the clone counter (modulo 256) and complement the clone counter sequence number bit.

Any updates to the cookie, clone counter, or their associated sequence number bits do not take effect, and shall not be committed to non-volatile storage, until the following **<MAC>Connect Confirm** message is received by the RCST.

The RCST uses its private Diffie-Hellman value,  $y$ , together with the message fields to derive the new session key value, as well as any new value for the cookie (see subclause 9.4.2).

#### 9.4.9.5 <MAC>Quick Key Exchange

The Quick Key Exchange message is used to start a cookie-dependent key exchange with the RCST, and also instructs the RCST whether to update its clone counter value. The `connection_id` is not used and shall be set to all '0'.

**Table 55: Quick Key Exchange message structure**

Quick_Key_Exchange () {	Bits	Bytes	Bit Number / Description
<b>Connection_ID</b>	<b>32</b>	<b>4</b>	<b>MAC connection identifier</b>
<b>Flags</b>	<b>8</b>	<b>1</b>	
Reserved	6		7..2: shall be 0
FL_Update_Counter	1		1:(yes/no) increment clone counter
FL_Session_Key	1		0: select session key 0 or 1
<b>Reserved</b>	<b>8</b>	<b>1</b>	<b>Reserved for future use, shall be 0</b>
<b>Nonce</b>		<b>P<sub>ns</sub></b>	<b>Random string nonce1</b>
}			

The `FL_Session_Key` bit specifies which session key of the security context to update.

If the `FL_Update_Counter` bit is set, it instructs the RCST to increment its clone detection counter.

The NCC will use its knowledge of the cookie value together with the fields of the response message from the RCST to derive the session key value (see subclause 9.4.3).

#### 9.4.9.6 <MAC>Quick Key Exchange Response

The Quick Key Exchange Response message authenticates the RCST and completes the cookie-dependent key exchange with the NCC. It also contains the current value of the clone detection counter. The `connection_id` is not used and shall be set to all '0'.

**Table 56: Quick Key Exchange Response message structure**

Quick_Key_Exchange_Response () {	Bits	Bytes	Bit Number / Description
<b>Connection_ID</b>	<b>32</b>	<b>4</b>	<b>MAC connection identifier</b>
<b>Flags</b>		<b>1</b>	
Reserved	6		7..2: shall be 0
FL_Cookie_SN	1		1: cookie sequence number
FL_Counter_SN	1		0: clone counter sequence number
<b>Clone_Counter</b>	<b>8</b>	<b>1</b>	<b>Current clone counter value</b>
<b>Nonce</b>		<b>P<sub>ns</sub></b>	<b>Random string nonce2</b>
<b>Authenticator</b>		<b>P<sub>ha</sub></b>	<b>Authentication value auth</b>
}			

The `FL_Cookie_SN` bit is the sequence number of the cookie used for authentication.

The `FL_Counter_SN` bit is the current sequence number of the clone detection counter. The `Clone_Counter` field is the current value of the counter. A clone collision has been detected if the NCC finds a mismatch from the expected value.

If the `FL_Update_Counter` bit was set by the NCC, the RCST will increment the value of the clone counter (modulo 256) and complement the clone counter sequence number bit.

The RCST uses the cookie value together with the message fields to derive the session key value (see subclause 9.4.3).

#### 9.4.9.7 <MAC>Explicit Key Exchange

The Explicit Key Exchange message is used to securely deliver an existing session key value to the RCST, and also instructs the RCST whether to update its clone counter value. The `connection_id` is not used and shall be set to all '0'.

**Table 57: Explicit Key Exchange message structure**

Explicit_Key_Exchange () {	Bits	Bytes	Bit Number / Description
<b>Connection_ID</b>	<b>32</b>	<b>4</b>	<b>MAC connection identifier</b>
<b>Flags</b>		<b>1</b>	
Reserved	6		7..2: shall be 0
FL_Update_Counter	1		1:(yes/no) increment clone counter
FL_Session_Key	1		0: select session key 0 or 1
<b>Reserved</b>	<b>8</b>	<b>1</b>	<b>Reserved for future use, shall be 0</b>
<b>Nonce</b>		<b>P<sub>ns</sub></b>	<b>Random string nonce1</b>
<b>Encryptedkey</b>		<b>P<sub>ea</sub></b>	<b>Encrypted session key</b>
}			

The `FL_Session_Key` bit specifies which session key of the security context to update.

If the `FL_Update_Counter` bit is set, it instructs the RCST to increment its clone detection counter.

The NCC has used its knowledge of the cookie value to encrypt the session key value (see subclause 9.4.4).

#### 9.4.9.8 <MAC>Explicit Key Exchange Response

The Explicit Key Exchange Response message authenticates the RCST and acknowledges receipt of the delivered key. It also contains the current value of the clone detection counter. The `connection_id` is not used and shall be set to all '0'.

**Table 58: Explicit Key Exchange Response message structure**

Explicit_Key_Exchange_Response () {	Bits	Bytes	Bit Number / Description
<b>Connection_ID</b>	<b>32</b>	<b>4</b>	<b>MAC connection identifier</b>
<b>Flags</b>		<b>1</b>	
Reserved	6		7..2: shall be 0
FL_Cookie_SN	1		1: cookie sequence number
FL_Counter_SN	1		0: clone counter sequence number
<b>Clone_Counter</b>	<b>8</b>	<b>1</b>	<b>Current clone counter value</b>
<b>Nonce</b>		<b>P<sub>ns</sub></b>	<b>Random string nonce2</b>
<b>Authenticator</b>		<b>P<sub>ha</sub></b>	<b>Authentication value auth</b>
}			

The `FL_Cookie_SN` bit is the sequence number of the cookie used for authentication and session key decryption. If the NCC determines that it has used the wrong cookie for session key encryption it will re-issue the <MAC>Explicit Key Exchange using the old cookie value.

The `FL_Counter_SN` bit is the current sequence number of the clone detection counter. The `Clone_Counter` field is the current value of the counter. A clone collision has been detected if the NCC finds a mismatch from the expected value.

If the `FL_Update_Counter` bit was set by the NCC, the RCST will increment the value of the clone counter (modulo 256) and complement the clone counter sequence number bit.

The RCST uses the cookie value together with the message fields to decrypt the session key value (see subclause 9.4.4).

### 9.4.9.9 <MAC>Wait

The Wait message is used by the RCST to extend the time the NCC waits for a reply to a given message. Upon receiving it, the NCC will reset its time-out value and retry count (see subclause 9.4.8.1).

**Table 59: Wait message structure**

<b>Wait () {</b>	<b>Bits</b>	<b>Bytes</b>	<b>Bit Number / Description</b>
<b>Connection_ID</b>	<b>32</b>	<b>4</b>	<b>MAC connection identifier</b>
<b>Message_Type</b>	<b>8</b>	<b>1</b>	<b>Type of message from NCC</b>
<b>Reserved</b>	<b>8</b>	<b>1</b>	<b>Reserved for future use, shall be 0</b>
<b>}</b>			

The **Message\_Type** field is the message type value of the message received from the NCC being processed. The **connection\_id** is not used and shall be set to all '0'. The RCST indicates that it is currently unable to send a reply to the message.



## Annex A (informative): Compliance table

**Table A.1: RCST compliance table**

PROFILE NAME	Baseline	ATM (option)	MPEG2 (option)
Access scheme			
Fixed MF-TDMA	•	•	•
Dynamic MF-TDMA	o	o	o
Traffic Burst Format			
ATM	•	•	•
MPEG2			•
Connectivity			
IP	•	•	•
Native ATM		•	o
Channel Coding			
RS	•	•	•
Convolutional	•	•	•
Turbo	•	•	•
CRC	•	•	•
Capacity Requests & management information			
Prefix	•	•	•
Data Unit Labelling	•	•	•
Mini-Slots	•	•	•
Contention Mini-Slot	•	•	•
Security mechanism	o	o	o
RCST forward link receivers			
1 DVB-S	•	•	•
2+ DVB-S	o	o	o
•:	Minimum Compliance Requirement for RCST.		
o:	Optional Compliance Point (statement by manufacturer).		

RCST symbol rate range goes from 128 to 4 000 ksymbol/s.

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## Bibliography

The following material, though not specifically referenced in the body of the present document (or not publicly available), gives supporting information.

- ITU-T Recommendation I.363-3: "B-ISDN ATM Adaptation Layer specification; Part 3: Type 3/4 AAL".

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## History

<b>Document history</b>		
V1.2.1	July 2000	One-step Approval Procedure OAP 20001124: 2000-07-26 to 2000-11-24
V1.2.2	December 2000	Publication